

booleans: THEORY

BEGIN

boolean: TYPE+

bool: TYPE+ = boolean

FALSE, TRUE: bool

NOT, \neg : [bool \rightarrow bool]

AND, $\&$, \wedge , OR, \vee , IMPLIES, \Rightarrow , \Rightarrow , WHEN, IFF, \Leftrightarrow , \Leftrightarrow : [bool, bool \rightarrow bool]

END booleans

```
equalities [T : TYPE] : THEORY
BEGIN

  = : [T, T → boolean]

END equalities
```

```
notequal[T: TYPE]: THEORY
BEGIN

  x, y: VAR T

  ≠(x, y): boolean = NOT (x = y);

  ≠: [T, T → bool] = ≠

END notequal
```

```
if_def [T: TYPE]: THEORY
  BEGIN

    IF: [boolean, T, T → T]

  END if_def
```

```

boolean_props: THEORY
BEGIN

  A, B: VAR bool

  bool_exclusive: POSTULATE NOT (FALSE = TRUE)

  bool_inclusive: POSTULATE A = FALSE OR A = TRUE

  not_def: POSTULATE (NOT A) = IF A THEN FALSE ELSE TRUE ENDIF

  and_def: POSTULATE (A AND B) = IF A THEN B ELSE FALSE ENDIF

  syand_def: POSTULATE & = AND

  or_def: POSTULATE (A OR B) = IF A THEN TRUE ELSE B ENDIF

  implies_def: POSTULATE (A IMPLIES B) = IF A THEN B ELSE TRUE ENDIF

  syimplies_def: POSTULATE  $\Rightarrow$  = IMPLIES

  when_def: POSTULATE (A WHEN B) = (B IMPLIES A)

  iff_def: POSTULATE (A IFF B) = ((A AND B) OR (NOT A AND NOT B))

  syiff_def: POSTULATE  $\Leftrightarrow$  = IFF

  excluded_middle: LEMMA A OR NOT A

END boolean_props

```

```
xor_def: THEORY
BEGIN

  A, B: VAR bool

  xor(A, B): bool = (A ≠ B)

  xor_def: LEMMA (A XOR B) = IF A THEN NOT B ELSE B ENDIF

END xor_def
```

```

quantifier_props[t: TYPE]: THEORY
BEGIN

  x: VAR t

  p, q: VAR [t → bool]

  not_exists: LEMMA (∃ x: p(x)) = NOT (∀ x: NOT p(x))

  exists_not: LEMMA (∃ x: NOT p(x)) = NOT (∀ x: p(x))

  exists_or: LEMMA
    (∃ x: p(x) OR q(x)) =
      ((∃ x: p(x)) OR (∃ x: q(x)))

  exists_implies: LEMMA
    (∃ x: p(x) IMPLIES q(x)) =
      ((∃ x: NOT p(x)) OR (∃ x: q(x)))

  exists_and: LEMMA
    (∃ x: p(x) AND q(x)) IMPLIES
      ((∃ x: p(x)) AND (∃ x: q(x)))

  not_forall: LEMMA (∀ x: p(x)) = NOT (∃ x: NOT p(x))

  forall_not: LEMMA (∀ x: NOT p(x)) = NOT (∃ x: p(x))

  forall_and: LEMMA
    (∀ x: p(x) AND q(x)) =
      ((∀ x: p(x)) AND (∀ x: q(x)))

  forall_or: LEMMA
    ((∀ x: p(x)) OR (∀ x: q(x))) IMPLIES
      (∀ x: p(x) OR q(x))

END quantifier_props

```

```
defined_types[t: TYPE]: THEORY
BEGIN

  pred: TYPE = [t → bool]

  PRED: TYPE = [t → bool]

  predicate: TYPE = [t → bool]

  PREDICATE: TYPE = [t → bool]

  setof: TYPE = [t → bool]

  SETOF: TYPE = [t → bool]

END defined_types
```



```

exists1[T: TYPE]: THEORY
BEGIN

  x, y: VAR T

  p, q: VAR pred[T]

  unique?(p): bool =  $\forall x, y: p(x) \text{ AND } p(y) \text{ IMPLIES } x = y$ 

  exists1(p): bool =  $(\exists x: p(x)) \text{ AND } \text{unique?}(p)$ 

  unique_lem: LEMMA
     $(\forall x: p(x) \text{ IMPLIES } q(x)) \text{ IMPLIES } (\text{unique?}(q) \text{ IMPLIES } \text{unique?}(p))$ 

  exists1_lem: LEMMA  $(\text{exists1! } x: p(x)) \text{ IMPLIES } (\exists x: p(x))$ 

END exists1

```

```

equality_props [T: TYPE]: THEORY
BEGIN

  x, y, z: VAR T

  b: VAR bool

  IF_true: POSTULATE IF TRUE THEN x ELSE y ENDIF = x

  IF_false: POSTULATE IF FALSE THEN x ELSE y ENDIF = y

  IF_same: LEMMA IF b THEN x ELSE x ENDIF = x

  reflexivity_of_equals: POSTULATE x = x

  transitivity_of_equals: POSTULATE x = y AND y = z IMPLIES x = z

  symmetry_of_equals: POSTULATE x = y IMPLIES y = x

END equality_props

```

```

if_props[s, t: TYPE]: THEORY
BEGIN

  a, b, c: VAR bool

  x, y: VAR s

  f: VAR [s → t]

  lift_if1: LEMMA
    f(IF a THEN x ELSE y ENDIF) =
      IF a THEN f(x) ELSE f(y) ENDIF

  lift_if2: LEMMA
    IF (IF a THEN b ELSE c ENDIF) THEN x ELSE y ENDIF =
      IF a
        THEN (IF b THEN x ELSE y ENDIF)
      ELSE (IF c THEN x ELSE y ENDIF)
      ENDIF

END if_props

```

```

functions[D, R: TYPE]: THEORY
BEGIN

  f, g: VAR [D → R]

  x, x1, x2: VAR D

  y: VAR R

  Drel: VAR PRED[[D]]

  Rrel: VAR PRED[[R]]

  extensionality_postulate: POSTULATE
    (∀ (x: D): f(x) = g(x) IFF f = g

  extensionality: LEMMA (∀ (x: D): f(x) = g(x) IMPLIES f = g

  congruence: POSTULATE f = g AND x1 = x2 IMPLIES f(x1) = g(x2)

  η: LEMMA (λ (x: D): f(x)) = f

  injective?(f): bool =
    (∀ x1, x2: (f(x1) = f(x2) ⇒ (x1 = x2)))

  surjective?(f): bool = (∀ y: (∃ x: f(x) = y))

  bijective?(f): bool = injective?(f) & surjective?(f)

  bij_is_inj: JUDGEMENT (bijective?) SUBTYPE_OF (injective?)

  bij_is_surj: JUDGEMENT (bijective?) SUBTYPE_OF (surjective?)

  domain(f): TYPE = D

  range(f): TYPE = R

  graph(f)(x, y): bool = (f(x) = y)

  preserves(f, Drel, Rrel): bool =
    ∀ x1, x2: Drel(x1, x2) IMPLIES Rrel(f(x1), f(x2))

```

preserves(Drel, Rrel)(f): bool = preserves(f, Drel, Rrel)

inverts(f, Drel, Rrel): bool =

$\forall x_1, x_2: \text{Drel}(x_1, x_2) \text{ IMPLIES } \text{Rrel}(f(x_2), f(x_1))$

inverts(Drel, Rrel)(f): bool = inverts(f, Drel, Rrel)

END functions

```
functions_alt[D, R: TYPE, Drel: PRED[[D]], Rrel: PRED[[R]]]: THEORY
BEGIN

  f: VAR [D → R]

  preserves: [[D → R] → bool] = preserves(Drel, Rrel)

  inverts: [[D → R] → bool] = inverts(Drel, Rrel)

END functions_alt
```

```
transpose [ $T_1$ ,  $T_2$ ,  $T_3$ : TYPE]: THEORY
BEGIN

   $f$ : VAR [ $T_1 \rightarrow [T_2 \rightarrow T_3]$ ]

   $x$ : VAR  $T_1$ 

   $y$ : VAR  $T_2$ 

  transpose( $f$ )( $y$ )( $x$ ):  $T_3 = f(x)(y)$ 

END transpose
```

```

restrict[T: TYPE, S: TYPE FROM T, R: TYPE]: THEORY
BEGIN

  f: VAR [T → R]

  s: VAR S

  restrict(f)(s): R = f(s)

  CONVERSION restrict

  injective_restrict: LEMMA
    injective?(f) IMPLIES injective?(restrict(f))

  restrict_of_inj_is_inj: JUDGEMENT restrict(f: (injective?[T, R])) HAS_TYPE
    (injective?[S, R])

END restrict

```



```
restrict_props[T: TYPE, R: TYPE]: THEORY
BEGIN

  f: VAR [T → R]

  restrict_full: LEMMA restrict[T, T, R](f) = f

END restrict_props
```

```
extend[T: TYPE, S: TYPE FROM T, R: TYPE, d: R]: THEORY
BEGIN

  f: VAR [S → R]

  t: VAR T

  extend(f)(t): R = IF S_pred(t) THEN f(t) ELSE d ENDIF

  restrict_extend: LEMMA restrict[T, S, R](extend(f)) = f

END extend
```

```
extend_bool[T: TYPE, S: TYPE FROM T]: THEORY
BEGIN

  CONVERSION extend[T, S, bool, FALSE]

END extend_bool
```

```
extend_props [T: TYPE, R: TYPE, d: R]: THEORY
BEGIN

  f: VAR [T → R]

  extend_full: LEMMA extend [T, T, R, d] (f) = f

END extend_props
```

```
extend_func_props[T : TYPE, S : TYPE FROM T, R : TYPE, d : R] : THEORY
BEGIN

  surjective_extend : JUDGEMENT extend[T, S, R, d](f : (surjective?[S, R])) HAS_TYPE
    (surjective?[T, R])

END extend_func_props
```

```
K_conversion[ $T_1$ ,  $T_2$ : TYPE]: THEORY
BEGIN

  K_conversion( $x$ :  $T_1$ )( $y$ :  $T_2$ ):  $T_1 = x$ 

END K_conversion
```

```

K_props[T1, T2: TYPE, S: TYPE FROM T1]: THEORY
BEGIN

  K_preserves: JUDGEMENT K_conversion[T1, T2](x: S)(y: T2) HAS_TYPE
              S

  K_preserves1: JUDGEMENT K_conversion[T1, T2](x: S) HAS_TYPE
                [T2 → S]

END K_props

```

```
identity[T: TYPE]: THEORY
BEGIN

  x: VAR T

  I: (bijective?[T, T]) = ( $\lambda x: x$ )

  id: (bijective?[T, T]) = ( $\lambda x: x$ )

  identity: (bijective?[T, T]) = ( $\lambda x: x$ )

END identity
```



```
identity_props[T: TYPE, S: TYPE FROM T]: THEORY
BEGIN

  x: VAR S

  I_preserves: JUDGEMENT I[T](x) HAS_TYPE S

  id_preserves: JUDGEMENT id[T](x) HAS_TYPE S

  identity_preserves: JUDGEMENT identity[T](x) HAS_TYPE S

END identity_props
```

```

relations[T: TYPE]: THEORY
BEGIN

  R: VAR PRED[[T]]

  x, y, z: VAR T

  eq: pred[[T]] = (λ x, y: x = y)

  reflexive?(R): bool = ∀ x: R(x, x)

  irreflexive?(R): bool = ∀ x: NOT R(x, x)

  symmetric?(R): bool = ∀ x, y: R(x, y) IMPLIES R(y, x)

  antisymmetric?(R): bool =
    ∀ x, y: R(x, y) & R(y, x) ⇒ x = y

  connected?(R): bool =
    ∀ x, y: x ≠ y IMPLIES R(x, y) OR R(y, x)

  transitive?(R): bool =
    ∀ x, y, z: R(x, y) & R(y, z) ⇒ R(x, z)

  equivalence?(R): bool =
    reflexive?(R) AND symmetric?(R) AND transitive?(R)

  equivalence: TYPE = (equivalence?)

  equiv_is_reflexive: JUDGEMENT (equivalence?) SUBTYPE_OF (reflexive?)

  equiv_is_symmetric: JUDGEMENT (equivalence?) SUBTYPE_OF (symmetric?)

  equiv_is_transitive: JUDGEMENT (equivalence?) SUBTYPE_OF
    (transitive?)

  RC(R)(x, y): bool = (x = y) OR R(x, y)

  TC(R)(x, y): INDUCTIVE bool =
    R(x, y) OR (∃ z: TC(R)(x, z) AND TC(R)(z, y))

  RTC(R)(x, y): bool = (x = y) OR TC(R)(x, y)

```

END relations

```

orders[T : TYPE] : THEORY
BEGIN

  x, y : VAR T

  ≤, < : VAR pred[[T]]

  p : VAR pred[T]

  preorder?(≤) : bool = reflexive?(≤) & transitive?(≤)

  preorder_is_reflexive : JUDGEMENT (preorder?) SUBTYPE_OF
    (reflexive?[T])

  preorder_is_transitive : JUDGEMENT (preorder?) SUBTYPE_OF
    (transitive?[T])

  equiv_is_preorder : JUDGEMENT (equivalence?[T]) SUBTYPE_OF
    (preorder?)

  partial_order?(≤) : bool = preorder?(≤) & antisymmetric?(≤)

  po_is_preorder : JUDGEMENT (partial_order?) SUBTYPE_OF (preorder?)

  po_is_antisymmetric : JUDGEMENT (partial_order?) SUBTYPE_OF
    (antisymmetric?[T])

  strict_order?(<) : bool = irreflexive?(<) & transitive?(<)

  strict_is_irreflexive : JUDGEMENT (strict_order?) SUBTYPE_OF
    (irreflexive?[T])

  strict_order_is_antisymmetric : JUDGEMENT (strict_order?) SUBTYPE_OF
    (antisymmetric?[T])

  strict_is_transitive : JUDGEMENT (strict_order?) SUBTYPE_OF
    (transitive?[T])

  dichotomous?(≤) : bool = (∀ x, y : (x ≤ y OR y ≤ x))

  total_order?(≤) : bool = partial_order?(≤) & dichotomous?(≤)

```

total_is_po: JUDGEMENT (total_order?) SUBTYPE_OF (partial_order?)

total_is_dichotomous: JUDGEMENT (total_order?) SUBTYPE_OF
(dichotomous?)

linear_order?(\leq): bool = total_order?(\leq)

linear_is_total: JUDGEMENT (linear_order?) SUBTYPE_OF (total_order?)

total_is_linear: JUDGEMENT (total_order?) SUBTYPE_OF (linear_order?)

trichotomous?($<$): bool =
($\forall x, y: x < y$ OR $y < x$ OR $x = y$)

strict_total_order?($<$): bool =
strict_order?($<$) & trichotomous?($<$)

strict_total_is_strict: JUDGEMENT (strict_total_order?) SUBTYPE_OF
(strict_order?)

strict_total_is_trichotomous: JUDGEMENT (strict_total_order?) SUBTYPE_OF
(trichotomous?)

well_founded?($<$): bool =
($\forall p:$
($\exists y: p(y)$) IMPLIES
($\exists (y: (p)): (\forall (x: (p)): (\text{NOT } x < y))$))

strict_well_founded?($<$): bool =
strict_order?($<$) & well_founded?($<$)

strict_well_founded_is_strict: JUDGEMENT (strict_well_founded?) SUBTYPE_OF
(strict_order?)

strict_well_founded_is_well_founded: JUDGEMENT (strict_well_founded?) SUBTYPE_OF
(well_founded?)

well_ordered?($<$): bool =
strict_total_order?($<$) & well_founded?($<$)

well_ordered_is_strict_total: JUDGEMENT (well_ordered?) SUBTYPE_OF
(strict_total_order?)

well_ordered_is_well_founded: JUDGEMENT (well_ordered?) SUBTYPE_OF
(well_founded?)

nonempty_pred: TYPE = {p: pred[T] | $\exists (x: T): p(x)$ }

pe: VAR pred[T]

upper_bound?(<)(x, pe): bool = $\forall (y: (pe)): y < x$

upper_bound?(<)(pe)(x): bool = upper_bound?(<)(x, pe)

lower_bound?(<)(x, pe): bool = $\forall (y: (pe)): x < y$

lower_bound?(<)(pe)(x): bool = lower_bound?(<)(x, pe)

least_upper_bound?(<)(x, pe): bool =
upper_bound?(<)(x, pe) AND
 $\forall y: \text{upper_bound?}(<)(y, \text{pe}) \text{ IMPLIES } (x < y \text{ OR } x = y)$

least_upper_bound?(<)(pe)(x): bool =
least_upper_bound?(<)(x, pe)

greatest_lower_bound?(<)(x, pe): bool =
lower_bound?(<)(x, pe) AND
 $\forall y: \text{lower_bound?}(<)(y, \text{pe}) \text{ IMPLIES } (y < x \text{ OR } x = y)$

greatest_lower_bound?(<)(pe)(x): bool =
greatest_lower_bound?(<)(x, pe)

END orders

```

orders_alt[T: TYPE, <: pred[[T]], pe: nonempty_pred[T]]: THEORY
BEGIN

  x: VAR T

  upper_bound?: [T → bool] = upper_bound?(<)(pe)

  least_upper_bound?: [T → bool] = least_upper_bound?(<)(pe)

  lower_bound?: [T → bool] = lower_bound?(<)(pe)

  greatest_lower_bound?: [T → bool] =
    greatest_lower_bound?(<)(pe)

  least_upper_bound_is_upper_bound: JUDGEMENT (least_upper_bound?) SUBTYPE_OF
    (upper_bound?)

  greatest_lower_bound_is_lower_bound: JUDGEMENT (greatest_lower_bound?) SUBTYPE_OF
    (lower_bound?)

END orders_alt

```

```

restrict_order_props[T : TYPE, S : TYPE FROM T] : THEORY
BEGIN

  reflexive_restrict : JUDGEMENT restrict(R : (reflexive?[T])) HAS_TYPE
    (reflexive?[S])

  irreflexive_restrict : JUDGEMENT restrict(R : (irreflexive?[T])) HAS_TYPE
    (irreflexive?[S])

  symmetric_restrict : JUDGEMENT restrict(R : (symmetric?[T])) HAS_TYPE
    (symmetric?[S])

  antisymmetric_restrict : JUDGEMENT restrict(R : (antisymmetric?[T])) HAS_TYPE
    (antisymmetric?[S])

  connected_restrict : JUDGEMENT restrict(R : (connected?[T])) HAS_TYPE
    (connected?[S])

  transitive_restrict : JUDGEMENT restrict(R : (transitive?[T])) HAS_TYPE
    (transitive?[S])

  equivalence_restrict : JUDGEMENT restrict(R : (equivalence?[T])) HAS_TYPE
    (equivalence?[S])

  preorder_restrict : JUDGEMENT restrict(R : (preorder?[T])) HAS_TYPE
    (preorder?[S])

  partial_order_restrict : JUDGEMENT restrict(R : (partial_order?[T])) HAS_TYPE
    (partial_order?[S])

  strict_order_restrict : JUDGEMENT restrict(R : (strict_order?[T])) HAS_TYPE
    (strict_order?[S])

  dichotomous_restrict : JUDGEMENT restrict(R : (dichotomous?[T])) HAS_TYPE
    (dichotomous?[S])

  total_order_restrict : JUDGEMENT restrict(R : (total_order?[T])) HAS_TYPE
    (total_order?[S])

  trichotomous_restrict : JUDGEMENT restrict(R : (trichotomous?[T])) HAS_TYPE
    (trichotomous?[S])

```


strict_total_order_restrict: JUDGEMENT restrict(R : (strict_total_order? $[T]$)) HAS_TYPE
(strict_total_order? $[S]$)

well_founded_restrict: JUDGEMENT restrict(R : (well_founded? $[T]$)) HAS_TYPE
(well_founded? $[S]$)

well_ordered_restrict: JUDGEMENT restrict(R : (well_ordered? $[T]$)) HAS_TYPE
(well_ordered? $[S]$)

END restrict_order_props

```

extend_order_props [T: TYPE, S: TYPE FROM T]: THEORY
BEGIN

  irreflexive_extend: JUDGEMENT extend [[T], [S], bool, FALSE] (R: (irreflexive?[S]))
    HAS_TYPE (irreflexive?[T])

  symmetric_extend: JUDGEMENT extend [[T], [S], bool, FALSE] (R: (symmetric?[S]))
    HAS_TYPE (symmetric?[T])

  antisymmetric_extend: JUDGEMENT extend
    [[T], [S], bool,
     FALSE] (R: (antisymmetric?[S]))
    HAS_TYPE (antisymmetric?[T])

  transitive_extend: JUDGEMENT extend [[T], [S], bool, FALSE] (R: (transitive?[S]))
    HAS_TYPE (transitive?[T])

  strict_order_extend: JUDGEMENT extend
    [[T], [S], bool, FALSE] (R: (strict_order?[S]))
    HAS_TYPE (strict_order?[T])

END extend_order_props

```

```

wf_induction[T: TYPE, <: (well_founded?[T])]: THEORY
BEGIN

  wf_induction: LEMMA
    (∀ (p: pred[T]):
      (∀ (x: T): (∀ (y: T): y < x IMPLIES p(y)) IMPLIES p(x)) IMPLIES
      (∀ (x: T): p(x)))

END wf_induction

```

```
measure_induction[T: TYPE, M: TYPE, m: [T → M], <: (well_founded?[M])]: THEORY
```

```
BEGIN
```

```
measure_induction: LEMMA
```

```
( $\forall$  (p: pred[T]):
```

```
( $\forall$  (x: T): ( $\forall$  (y: T): m(y) < m(x) IMPLIES p(y)) IMPLIES p(x)) IMPLIES
```

```
( $\forall$  (x: T): p(x))
```

```
END measure_induction
```

```
epsilon[T: TYPE+]: THEORY
BEGIN

  p: VAR pred[T]

  x: VAR T

  epsilon(p): T

  epsilon_ax: AXIOM (exists x: p(x) => p(epsilon(p)))

END epsilon
```

```

sets[T: TYPE]: THEORY
BEGIN

  set: TYPE = setof[T]

  x, y: VAR T

  a, b, c: VAR set

  p: VAR PRED[T]

  (x ∈ a): bool = a(x)

  empty?(a): bool = (∀ x: NOT (x ∈ a))

  ∅: set = {x | FALSE}

  nonempty?(a): bool = NOT empty?(a)

  nonempty_set: TYPE = (nonempty?)

  full?(a): bool = (∀ x: (x ∈ a))

  fullset: set = {x | TRUE}

  nontrivial?(a): bool = a ≠ ∅ & a ≠ fullset

  (a ⊆ b): bool = (∀ x: (x ∈ a) ⇒ (x ∈ b))

  (a ⊂ b): bool = (a ⊆ b) & a ≠ b

  (a ∪ b): set = {x | (x ∈ a) OR (x ∈ b)}

  (a ∩ b): set = {x | (x ∈ a) AND (x ∈ b)}

  disjoint?(a, b): bool = empty?((a ∩ b))

  ā: set = {x | NOT (x ∈ a)}

  (a \ b): set = {x | (x ∈ a) AND NOT (x ∈ b)}

  symmetric_difference(a, b): set = ((a \ b) ∪ (b \ a))

```

every(p)(a): bool = $\forall (x : (a)) : p(x)$

every(p, a): bool = $\forall (x : (a)) : p(x)$

some(p)(a): bool = $\exists (x : (a)) : p(x)$

some(p, a): bool = $\exists (x : (a)) : p(x)$

singleton?(a): bool =
 $(\exists (x : (a)) : (\forall (y : (a)) : x = y))$

singleton(x): (singleton?) = $\{y \mid y = x\}$

$(a \cup \{x\})$: (nonempty?) = $\{y \mid x = y \text{ OR } (y \in a)\}$

$(a \setminus \{x\})$: set = $\{y \mid x \neq y \text{ AND } (y \in a)\}$

choose(p : (nonempty?)): (p)

choose_is_epsilon: AXIOM

$\forall (p : (nonempty?)) : \text{choose}(p) = \varepsilon(p)$

the(p : (singleton?)): (p)

the_lem: LEMMA $\forall (p : (singleton?)) : \text{the}(p) = \varepsilon(p)$

the_prop: LEMMA $\forall (p : (singleton?)) : p(\text{the}(p))$

singleton_elt(a : (singleton?)): $T = \text{the! } x : (x \in a)$

CONVERSION+ singleton_elt

is_singleton: LEMMA

$\forall a$:

(nonempty?(a) AND $\forall x, y : a(x)$ AND $a(y)$ IMPLIES $(x = y)$) IMPLIES
singleton?(a)

singleton_elt_lem: LEMMA

singleton?(a) AND $a(x)$ IMPLIES singleton_elt(a) = x

```

singleton_elt_def: LEMMA
  singleton?(a) IMPLIES singleton_elt(a) = choose(a)

singleton_singleton: LEMMA
  singleton?(a) IMPLIES ( $\exists x: a = \text{singleton}(x)$ )

singleton_rew: LEMMA singleton_elt(singleton(x)) = x

AUTO_REWRITE+ singleton_rew

rest(a): set =
  IF empty?(a) THEN a ELSE ( $a \setminus \{\text{choose}(a)\}$ ) ENDIF

setofsets: TYPE = setof[setof[T]]

A, B: VAR setofsets

powerset(a): setofsets = {b | ( $b \subseteq a$ )}

 $\bigcup A$ : set = {x |  $\exists (a: (A)): a(x)$ }

 $\bigcap A$ : set = {x |  $\forall (a: (A)): a(x)$ }

nonempty_singleton: JUDGEMENT (singleton?) SUBTYPE_OF (nonempty?)

nonempty_union1: JUDGEMENT union(a: (nonempty?), b: set) HAS_TYPE
  (nonempty?)

nonempty_union2: JUDGEMENT union(a: set, b: (nonempty?)) HAS_TYPE
  (nonempty?)

END sets

```



```

sets_lemmas[T : TYPE]: THEORY
BEGIN

  x, y: VAR T

  a, b, c: VAR set[T]

  A, B: VAR setofsets[T]

  extensionality: LEMMA
    ( $\forall x: (x \in a) \text{ IFF } (x \in b)$ ) IMPLIES  $(a = b)$ 

  emptyset_is_empty?: LEMMA empty?(a) IFF  $a = \emptyset$ 

  empty_no_members: LEMMA NOT  $(x \in \emptyset)$ 

  emptyset_min: LEMMA  $(a \subseteq \emptyset)$  IMPLIES  $a = \emptyset$ 

  nonempty_member: LEMMA nonempty?(a) IFF  $\exists x: (x \in a)$ 

  fullset_member: LEMMA  $(x \in \text{fullset})$ 

  fullset_max: LEMMA  $(\text{fullset} \subseteq a)$  IMPLIES  $a = \text{fullset}$ 

  fullset_is_full?: LEMMA full?(a) IFF  $a = \text{fullset}$ 

  nonempty_exists: LEMMA nonempty?(a) IFF  $\exists (x: (a))$ : TRUE

  subset_emptyset: LEMMA  $(\emptyset \subseteq a)$ 

  subset_fullset: LEMMA  $(a \subseteq \text{fullset})$ 

  subset_reflexive: LEMMA  $(a \subseteq a)$ 

  subset_antisymmetric: LEMMA  $(a \subseteq b)$  AND  $(b \subseteq a)$  IMPLIES  $a = b$ 

  subset_transitive: LEMMA
     $(a \subseteq b)$  AND  $(b \subseteq c)$  IMPLIES  $(a \subseteq c)$ 

  subset_partial_order: LEMMA partial_order?(subset?[T])

  subset_is_partial_order: JUDGEMENT subset?[T] HAS_TYPE

```

(partial_order?[set[T]])

strict_subset_irreflexive: LEMMA NOT ($a \subset a$)

strict_subset_transitive: LEMMA
($a \subset b$) AND ($b \subset c$) IMPLIES ($a \subset c$)

strict_subset_strict_order: LEMMA strict_order?(strict_subset?[T])

strict_subset_is_strict_order: JUDGEMENT strict_subset?[T] HAS_TYPE
(strict_order?[set[T]])

union_idempotent: LEMMA ($a \cup a$) = a

union_commutative: LEMMA ($a \cup b$) = ($b \cup a$)

union_associative: LEMMA (($a \cup b$) $\cup c$) = ($a \cup (b \cup c)$)

union_empty: LEMMA ($a \cup \emptyset$) = a

union_full: LEMMA ($a \cup \text{fullset}$) = fullset

union_subset1: LEMMA ($a \subseteq (a \cup b)$)

union_subset2: LEMMA ($a \subseteq b$) IMPLIES ($a \cup b$) = b

subset_union: LEMMA
(($a \cup b$) $\subseteq c$) = (($a \subseteq c$) AND ($b \subseteq c$))

union_upper_bound: LEMMA
($a \subseteq c$) AND ($b \subseteq c$) IMPLIES (($a \cup b$) $\subseteq c$)

union_difference: LEMMA ($a \cup b$) = ($a \cup (b \setminus a)$)

union_diff_subset: LEMMA ($a \subseteq b$) IMPLIES ($a \cup (b \setminus a)$) = b

intersection_idempotent: LEMMA ($a \cap a$) = a

intersection_commutative: LEMMA ($a \cap b$) = ($b \cap a$)

intersection_associative: LEMMA
(($a \cap b$) $\cap c$) = ($a \cap (b \cap c)$)

intersection_empty: LEMMA $(a \cap \emptyset) = \emptyset$

intersection_full: LEMMA $(a \cap \text{fullset}) = a$

intersection_subset1: LEMMA $((a \cap b) \subseteq a)$

intersection_subset2: LEMMA $(a \subseteq b) \text{ IMPLIES } (a \cap b) = a$

intersection_lower_bound: LEMMA
 $(c \subseteq a) \text{ AND } (c \subseteq b) \text{ IMPLIES } (c \subseteq (a \cap b))$

distribute_intersection_union: LEMMA
 $(a \cap (b \cup c)) = ((a \cap b) \cup (a \cap c))$

distribute_union_intersection: LEMMA
 $(a \cup (b \cap c)) = ((a \cup b) \cap (a \cup c))$

complement_emptyset: LEMMA $\overline{\emptyset[T]} = \text{fullset}$

complement_fullset: LEMMA $\overline{\text{fullset}[T]} = \emptyset$

complement_complement: LEMMA $\overline{\overline{a}} = a$

complement_equal: LEMMA $\overline{\overline{a}} = \overline{\overline{b}} \text{ IFF } a = b$

subset_complement: LEMMA $(\overline{\overline{a}} \subseteq \overline{\overline{b}}) \text{ IFF } (b \subseteq a)$

demorgan1: LEMMA $\overline{(a \cup b)} = (\overline{a} \cap \overline{b})$

demorgan2: LEMMA $\overline{(a \cap b)} = (\overline{a} \cup \overline{b})$

difference_emptyset1: LEMMA $(a \setminus \emptyset) = a$

difference_emptyset2: LEMMA $(\emptyset \setminus a) = \emptyset$

difference_fullset1: LEMMA $(a \setminus \text{fullset}) = \emptyset$

difference_fullset2: LEMMA $(\text{fullset} \setminus a) = \overline{a}$

difference_intersection: LEMMA $(a \setminus b) = (a \cap \overline{b})$

difference_difference1: LEMMA

$$((a \setminus b) \setminus c) = (a \setminus (b \cup c))$$

difference_difference2: LEMMA

$$(a \setminus (b \setminus c)) = ((a \setminus b) \cup (a \cap c))$$

difference_subset: LEMMA $((a \setminus b) \subseteq a)$

difference_subset2: LEMMA $(a \subseteq b)$ IMPLIES $(a \setminus b) = \emptyset$

difference_disjoint: LEMMA $\text{disjoint?}(a, (b \setminus a))$

difference_disjoint2: LEMMA $\text{disjoint?}(a, b)$ IMPLIES $(a \setminus b) = a$

diff_union_inter: LEMMA

$$((a \cup b) \setminus a) = (b \setminus (a \cap b))$$

nonempty_add: LEMMA NOT $\text{empty?}((a \cup \{x\}))$

member_add: LEMMA $(x \in a)$ IMPLIES $(a \cup \{x\}) = a$

member_add_reduce: LEMMA

$$(x \in (a \cup \{y\})) = (x = y \text{ OR } (x \in a))$$

member_remove: LEMMA NOT $(x \in a)$ IMPLIES $(a \setminus \{x\}) = a$

add_remove_member: LEMMA

$$(x \in a) \text{ IMPLIES } ((a \setminus \{x\}) \cup \{x\}) = a$$

remove_add_member: LEMMA

$$\text{NOT } (x \in a) \text{ IMPLIES } ((a \cup \{x\}) \setminus \{x\}) = a$$

subset_add: LEMMA $(a \subseteq (a \cup \{x\}))$

add_as_union: LEMMA $(a \cup \{x\}) = (a \cup \text{singleton}(x))$

singleton_as_add: LEMMA $\text{singleton}(x) = (\emptyset \cup \{x\})$

subset_remove: LEMMA $((a \setminus \{x\}) \subseteq a)$

remove_as_difference: LEMMA

$$(a \setminus \{x\}) = (a \setminus \text{singleton}(x))$$

remove_member_singleton: LEMMA $(\text{singleton}(x) \setminus \{x\}) = \emptyset$

choose_rest: LEMMA

NOT $\text{empty?}(a)$ IMPLIES $(\text{rest}(a) \cup \{\text{choose}(a)\}) = a$

choose_member: LEMMA NOT $\text{empty?}(a)$ IMPLIES $(\text{choose}(a) \in a)$

choose_not_member: LEMMA

NOT $\text{empty?}(a)$ IMPLIES NOT $(\text{choose}(a) \in \text{rest}(a))$

rest_not_equal: LEMMA NOT $\text{empty?}(a)$ IMPLIES $\text{rest}(a) \neq a$

rest_member: LEMMA $(x \in \text{rest}(a))$ IMPLIES $(x \in a)$

rest_subset: LEMMA $(\text{rest}(a) \subseteq a)$

choose_add: LEMMA

$\text{choose}((a \cup \{x\})) = x$ OR
 $(\text{choose}((a \cup \{x\})) \in a)$

choose_rest_or: LEMMA

$(x \in a)$ IMPLIES $(x \in \text{rest}(a))$ OR $x = \text{choose}(a)$

choose_singleton: LEMMA $\text{choose}(\text{singleton}(x)) = x$

rest_singleton: LEMMA $\text{rest}(\text{singleton}(x)) = \emptyset[T]$

singleton_subset: LEMMA $(x \in a)$ IFF $(\text{singleton}(x) \subseteq a)$

rest_empty_lem: LEMMA

NOT $\text{empty?}(a)$ AND $\text{empty?}(\text{rest}(a))$ IMPLIES
 $a = \text{extend}[T, (a), \text{bool}, \text{FALSE}](\text{singleton}(\text{choose}(a)))$

singleton_disjoint: LEMMA

NOT $(x \in a)$ IMPLIES $\text{disjoint?}(\text{singleton}(x), a)$

disjoint_remove_left: LEMMA

$\text{disjoint?}(a, b)$ IMPLIES $\text{disjoint?}((a \setminus \{x\}), b)$

disjoint_remove_right: LEMMA

$\text{disjoint?}(a, b)$ IMPLIES $\text{disjoint?}(a, (b \setminus \{x\}))$

union_disj_remove_left: LEMMA

disjoint?(a , b) AND $a(x)$ IMPLIES
 $((a \setminus \{x\}) \cup b) = ((a \cup b) \setminus \{x\})$

union_disj_remove_right: LEMMA

disjoint?(a , b) AND $b(x)$ IMPLIES
 $(a \cup (b \setminus \{x\})) = ((a \cup b) \setminus \{x\})$

subset_powerset: LEMMA ($a \subseteq b$) IFF ($a \in \text{powerset}(b)$)

empty_powerset: LEMMA empty?(a) IFF singleton?(powerset(a))

powerset_emptyset: LEMMA ($\emptyset \in \text{powerset}(a)$)

nonempty_powerset: JUDGEMENT powerset(a) HAS_TYPE
(nonempty?[set[T]])

powerset_union: LEMMA $\bigcup \text{powerset}(a) = a$

powerset_intersection: LEMMA empty?($\bigcap \text{powerset}(a)$)

powerset_subset: LEMMA

($a \subseteq b$) IFF ($\text{powerset}(a) \subseteq \text{powerset}(b)$)

Union_empty: LEMMA

empty?($\bigcup A$) IFF empty?(A) OR every(empty?)(A)

Union_full: LEMMA

full?($\bigcup A$) IFF ($\forall x: \exists (a: (A)): (x \in a)$)

Union_member: LEMMA

($x \in \bigcup A$) IFF ($\exists (a: (A)): (x \in a)$)

Union_subset: LEMMA $\forall (a: (A)): (a \subseteq \bigcup A)$

Union_surjective: JUDGEMENT Union HAS_TYPE

(surjective?[setofsets[T], set[T]])

Union_emptyset_rew: LEMMA $\bigcup \emptyset[\text{set}[T]] = \emptyset$

Union_union_rew: LEMMA

nonempty?(A) IMPLIES
 $\bigcup A = (\text{choose}(A) \cup \bigcup \text{rest}(A))$

Intersection_empty: LEMMA
empty?($\bigcap A$) IFF
 $(\forall x: \exists (a: (A)): \text{NOT } (x \in a))$

Intersection_full: LEMMA full?($\bigcap A$) IFF every(full?)(A)

Intersection_member: LEMMA
 $(x \in \bigcap A)$ IFF $(\forall (a: (A)): (x \in a))$

Intersection_empty_full: COROLLARY full?($\bigcap \emptyset[\text{set}[T]]$)

Intersection_surjective: JUDGEMENT Intersection HAS_TYPE
(surjective?[setofsets[T], set[T]])

Intersection_intersection_rew: LEMMA
nonempty?(A) IMPLIES
 $\bigcap A = (\text{choose}(A) \cap \bigcap \text{rest}(A))$

Complement(A): setofsets[T] =
 $\{a \mid \exists (b: (A)): a = \bar{b}\}$

Complement_empty: LEMMA empty?(Complement(A)) IFF empty?(A)

Complement_full: LEMMA full?(Complement(A)) IFF full?(A)

Complement_Complement: LEMMA Complement(Complement(A)) = A

subset_Complement: LEMMA
(Complement(A) \subseteq Complement(B)) IFF $(A \subseteq B)$

Complement_bijective: JUDGEMENT Complement HAS_TYPE
(bijective?[setofsets[T], setofsets[T]])

Demorgan1: LEMMA $\overline{\bigcup A} = \bigcap \text{Complement}(A)$

Demorgan2: LEMMA $\overline{\bigcap A} = \bigcup \text{Complement}(A)$

END sets_lemmas

```

function_inverse_def [D: TYPE, R: TYPE]: THEORY
BEGIN

  d: VAR D

  r: VAR R

  f: VAR [D → R]

  g: VAR [R → D]

  left_inverse?(g, f): bool = ∀ d: g(f(d)) = d

  right_inverse?(g, f): bool = ∀ r: f(g(r)) = r

  inverse?(g, f): bool =
    ∀ r: (∃ d: f(d) = r) ⇒ f(g(r)) = r

  left_inverse?(f)(g): MACRO bool = left_inverse?(g, f)

  right_inverse?(f)(g): MACRO bool = right_inverse?(g, f)

  inverse?(f)(g): MACRO bool = inverse?(g, f)

  left_inverse_is_inverse: LEMMA
    ∀ f, (g: (λ (g): left_inverse?(g, f))):
      inverse?(g, f)

  left_inj_surj: LEMMA
    ∀ f, (g: (λ (g): left_inverse?(g, f))):
      injective?(f) AND surjective?(g)

  inj_left_alt: LEMMA
    ∀ (f: (injective?[D, R])), (g: (λ (g): inverse?(g, f))):
      left_inverse?(g, f)

  surj_inv_alt: COROLLARY
    ∀ (f: (injective?[D, R])), (g: (λ (g): inverse?(g, f))):
      surjective?(g)

  injective_inverse_alt: LEMMA
    ∀ (f: (injective?[D, R])), (g: (λ (g): inverse?(g, f))):

```


$$r = f(d) \Rightarrow g(r) = d$$

comp_inverse_left_inj_alt: LEMMA

$$\forall (f: (\text{injective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))): \\ g(f(d)) = d$$

noninjective_inverse_exists: LEMMA

$$\forall f: \text{injective?}(f) \text{ OR } (\exists g: \text{inverse?}(g, f))$$

right_inverse_is_inverse: LEMMA

$$\forall f, (g: (\lambda (g): \text{right_inverse?}(g, f))): \\ \text{inverse?}(g, f)$$

right_surj_inj: LEMMA

$$\forall f, (g: (\lambda (g): \text{right_inverse?}(g, f))): \\ \text{surjective?}(f) \text{ AND } \text{injective?}(g)$$

surj_right_alt: LEMMA

$$\forall (f: (\text{surjective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))): \\ \text{right_inverse?}(g, f)$$

inj_inv_alt: COROLLARY

$$\forall (f: (\text{surjective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))): \\ \text{injective?}(g)$$

surjective_inverse_alt: LEMMA

$$\forall (f: (\text{surjective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))): \\ g(r) = d \Rightarrow r = f(d)$$

comp_inverse_right_surj_alt: LEMMA

$$\forall (f: (\text{surjective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))): \\ f(g(r)) = r$$

surjective_inverse_exists: LEMMA

$$\forall (f: (\text{surjective?}[D, R])): \exists g: \text{inverse?}(g, f)$$

left_right_bij: COROLLARY

$$\forall f, g: \\ \text{right_inverse?}(g, f) \text{ AND } \text{left_inverse?}(g, f) \Rightarrow \\ \text{bijective?}(f) \text{ AND } \text{bijective?}(g)$$

bij_left_right: COROLLARY

$\forall (f: (\text{bijective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))):$
 $\text{right_inverse?}(g, f) \text{ AND } \text{left_inverse?}(g, f)$

bij_inv_is_bij_alt: COROLLARY

$\forall (f: (\text{bijective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))):$
 $\text{bijective?}(g)$

bijective_inverse_alt: COROLLARY

$\forall (f: (\text{bijective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))):$
 $g(r) = d \text{ IFF } r = f(d)$

comp_inverse_right_alt: COROLLARY

$\forall (f: (\text{bijective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))):$
 $f(g(r)) = r$

comp_inverse_left_alt: COROLLARY

$\forall (f: (\text{bijective?}[D, R])), (g: (\lambda (g): \text{inverse?}(g, f))):$
 $g(f(d)) = d$

bijective_inverse_exists: LEMMA

$\forall (f: (\text{bijective?}[D, R])):$
 $\text{exists1}(\lambda (g): \text{inverse?}(g, f))$

exists_inv1: LEMMA

$(\exists g: \text{TRUE}) \text{ IFF}$
 $((\exists (d: D): \text{TRUE}) \text{ OR } (\forall (r: R): \text{FALSE}))$

exists_inv2: LEMMA

$(\exists (f: (\text{surjective?}[D, R])): \text{TRUE}) \Rightarrow$
 $((\exists (d: D): \text{TRUE}) \text{ OR } (\forall (r: R): \text{FALSE}))$

exists_inv3: LEMMA

$(\exists f: \text{NOT injective?}(f)) \Rightarrow$
 $((\exists (d: D): \text{TRUE}) \text{ OR } (\forall (r: R): \text{FALSE}))$

END function_inverse_def

```

function_inverse[D: TYPE+, R: TYPE]: THEORY
BEGIN

  x: VAR D

  y: VAR R

  f: VAR [D → R]

  g: VAR [R → D]

  inverse(f)(y): D = (ε! x: f(x) = y)

  unique_bijective_inverse: JUDGEMENT inverse(f: (bijective?[D, R]))(y) HAS_TYPE
    {x: D | f(x) = y}

  bijective_inverse_is_bijective: JUDGEMENT inverse(f: (bijective?[D, R])) HAS_TYPE
    (bijective?[R, D])

  surjective_inverse: LEMMA
    ∀ (f: (surjective?[D, R])):
      inverse(f)(y) = x IMPLIES y = f(x)

  inverse_surjective: LEMMA
    ∀ (f: (surjective?[D, R])): f(inverse(f)(y)) = y

  injective_inverse: LEMMA
    ∀ (f: (injective?[D, R])):
      y = f(x) IMPLIES inverse(f)(y) = x

  inverse_injective: LEMMA
    ∀ (f: (injective?[D, R])): inverse(f)(f(x)) = x

  bijective_inverse: LEMMA
    ∀ (f: (bijective?[D, R])):
      inverse(f)(y) = x IFF y = f(x)

  bij_inv_is_bij: LEMMA bijective?(f) IMPLIES bijective?(inverse(f))

  surj_right: LEMMA
    surjective?(f) IFF right_inverse?(inverse(f), f)

```

```

inj_left: LEMMA injective?(f) IFF left_inverse?(inverse(f), f)

inj_inv: LEMMA surjective?(f) IMPLIES injective?(inverse(f))

surj_inv: LEMMA injective?(f) IMPLIES surjective?(inverse(f))

inv_inj_is_surj: JUDGEMENT inverse(f: (injective?[D, R])) HAS_TYPE
  (surjective?[R, D])

inv_surj_is_inj: JUDGEMENT inverse(f: (surjective?[D, R])) HAS_TYPE
  (injective?[R, D])

comp_inverse_right_surj: LEMMA
   $\forall (f: (\text{surjective?}[D, R])): f(\text{inverse}(f)(y)) = y$ 

comp_inverse_left_inj: LEMMA
   $\forall (f: (\text{injective?}[D, R])): \text{inverse}(f)(f(x)) = x$ 

comp_inverse_right: LEMMA
   $\forall (f: (\text{bijective?}[D, R])): f(\text{inverse}(f)(y)) = y$ 

comp_inverse_left: LEMMA
   $\forall (f: (\text{bijective?}[D, R])): \text{inverse}(f)(f(x)) = x$ 

END function_inverse

```

```

function_inverse_alt[D: TYPE, R: TYPE]: THEORY
BEGIN

  ASSUMING
    inverse_types: ASSUMPTION
      ( $\exists (d: D): \text{TRUE}$ ) OR ( $\forall (r: R): \text{FALSE}$ )
  ENDASSUMING

  d: VAR D

  r: VAR R

  f: VAR [D  $\rightarrow$  R]

  g: VAR [R  $\rightarrow$  D]

  inverses(f): TYPE+ = ( $\lambda (g: [R \rightarrow D]): \text{inverse?}(g, f)$ )

  inverse_alt(f): inverses(f) =
    choose({g: inverses(f) | TRUE})

  bijective_inverse_is_inverse_alt: COROLLARY
     $\forall (f: (\text{bijective?}[D, R])), (g: \text{inverses}(f)):$ 
      g = inverse_alt(f)

  unique_bijective_inverse_alt: JUDGEMENT inverse_alt(f: (bijective?[D, R]))(r)
    HAS_TYPE {d | f(d) = r}

  bijective_inverse_alt_is_bijective: JUDGEMENT inverse_alt(f: (bijective?[D, R]))
    HAS_TYPE (bijective?[R, D])

  inv_inj_is_surj_alt: JUDGEMENT inverse_alt(f: (injective?[D, R])) HAS_TYPE
    (surjective?[R, D])

  inv_surj_is_inj_alt: JUDGEMENT inverse_alt(f: (surjective?[D, R])) HAS_TYPE
    (injective?[R, D])

END function_inverse_alt

```

```

function_image[D, R: TYPE]: THEORY
BEGIN

  f: VAR [D → R]

  x: VAR D

  y: VAR R

  X, X1, X2: VAR set[D]

  Y, Y1, Y2: VAR set[R]

  fun_exists: LEMMA
    (∃ y: TRUE) OR (NOT ∃ x: TRUE) IMPLIES (∃ f: TRUE)

  image(f, X): set[R] =
    {y: R | (∃ (x: (X)): y = f(x))}

  image(f)(X): set[R] = image(f, X)

  inverse_image(f, Y): set[D] = {x: D | (f(x) ∈ Y)}

  inverse_image(f)(Y): set[D] = inverse_image(f, Y)

  image_inverse_image: LEMMA
    (image(f, inverse_image(f, Y)) ⊆ Y)

  inverse_image_image: LEMMA
    (X ⊆ inverse_image(f, image(f, X)))

  image_subset: LEMMA
    (X1 ⊆ X2) IMPLIES (image(f, X1) ⊆ image(f, X2))

  inverse_image_subset: LEMMA
    (Y1 ⊆ Y2) IMPLIES
    (inverse_image(f, Y1) ⊆ inverse_image(f, Y2))

  image_union: LEMMA
    image(f, (X1 ∪ X2)) = (image(f, X1) ∪ image(f, X2))

  image_intersection: LEMMA

```

$$(\text{image}(f, (X_1 \cap X_2)) \subseteq (\text{image}(f, X_1) \cap \text{image}(f, X_2)))$$

inverse_image_union: LEMMA

$$\text{inverse_image}(f, (Y_1 \cup Y_2)) = \\ (\text{inverse_image}(f, Y_1) \cup \text{inverse_image}(f, Y_2))$$

inverse_image_intersection: LEMMA

$$\text{inverse_image}(f, (Y_1 \cap Y_2)) = \\ (\text{inverse_image}(f, Y_1) \cap \text{inverse_image}(f, Y_2))$$

inverse_image_complement: LEMMA

$$\text{inverse_image}(f, \overline{Y}) = \overline{\text{inverse_image}(f, Y)}$$

END function_image

```

function_props [T1, T2, T3: TYPE]: THEORY
BEGIN

  x: VAR T1

  f1: VAR [T1 → T2]

  f2: VAR [T2 → T3]

  X: VAR set[T1]

  R1: VAR PRED [[T1]]

  R2: VAR PRED [[T2]]

  R3: VAR PRED [[T3]]

  f2 ∘ f1(x): T3 = f2(f1(x))

  composition_injective: JUDGEMENT O(f2: (injective?[T2, T3]), f1: (injective?[T1, T2]))
    HAS_TYPE (injective?[T1, T3])

  composition_surjective: JUDGEMENT O(f2: (surjective?[T2, T3]),
    f1: (surjective?[T1, T2]))
    HAS_TYPE (surjective?[T1, T3])

  composition_bijective: JUDGEMENT O(f2: (bijective?[T2, T3]), f1: (bijective?[T1, T2]))
    HAS_TYPE (bijective?[T1, T3])

  image_composition: LEMMA
    image(f2, image(f1, X)) = image(f2 ∘ f1, X)

  preserves_composition: LEMMA
    preserves(f1, R1, R2) AND preserves(f2, R2, R3) IMPLIES
    preserves(f2 ∘ f1, R1, R3)

  inverts_composition1: LEMMA
    preserves(f1, R1, R2) AND inverts(f2, R2, R3) IMPLIES
    inverts(f2 ∘ f1, R1, R3)

  inverts_composition2: LEMMA
    inverts(f1, R1, R2) AND preserves(f2, R2, R3) IMPLIES

```


$\text{inverts}(f_2 \circ f_1, R_1, R_3)$

END function_props

```

function_props_alt [T1, T2, T3: TYPE, R1: PRED[[T1]], R2: PRED[[T2]],
                  R3: PRED[[T3]]]: THEORY
BEGIN

composition_preserves: JUDGEMENT O(f2: (preserves[T2, T3, R2, R3]),
                                   f1: (preserves[T1, T2, R1, R2]))
HAS_TYPE (preserves[T1, T3, R1, R3])

composition_inverts1: JUDGEMENT O(f2: (preserves[T2, T3, R2, R3]),
                                   f1: (inverts[T1, T2, R1, R2]))
HAS_TYPE (inverts[T1, T3, R1, R3])

composition_inverts2: JUDGEMENT O(f2: (inverts[T2, T3, R2, R3]),
                                   f1: (preserves[T1, T2, R1, R2]))
HAS_TYPE (inverts[T1, T3, R1, R3])

END function_props_alt

```

```
function_props2[T1, T2, T3, T4: TYPE]: THEORY
BEGIN

  f1: VAR [T1 → T2]

  f2: VAR [T2 → T3]

  f3: VAR [T3 → T4]

  ASSOC: LEMMA (f3 ∘ f2) ∘ f1 = f3 ∘ (f2 ∘ f1)

END function_props2
```

```

relation_defs[T1, T2: TYPE]: THEORY
BEGIN

  R, S: VAR pred[[T2]]

  x: VAR T1

  y: VAR T2

  X: VAR set[T1]

  Y: VAR set[T2]

  domain?(R)(x: T1): bool = ∃ (y: T2): R(x, y)

  range?(R)(y: T2): bool = ∃ (x: T1): R(x, y)

  domain(R): TYPE = {x: T1 | ∃ (y: T2): R(x, y)}

  range(R): TYPE = {y: T2 | ∃ (x: T1): R(x, y)}

  rinverse(R)(y, x): bool = R(x, y)

  rcomplement(R)(x, y): bool = NOT R(x, y)

  runion(R, S)(x, y): bool = R(x, y) OR S(x, y)

  rintersction(R, S)(x, y): bool = R(x, y) AND S(x, y)

  image(R, X): set[T2] =
    {y: T2 | ∃ (x: (X)): R(x, y)}

  image(R)(X): set[T2] = image(R, X)

  preimage(R, Y): set[T1] =
    {x: T1 | ∃ (y: (Y)): R(x, y)}

  preimage(R)(Y): set[T1] = preimage(R, Y)

  postcondition(R, X): set[T2] =
    {y: T2 | ∃ (x: (X)): R(x, y)}

```

```

postcondition(R)(X): set[T2] = postcondition(R, X)

precondition(R, Y): set[T1] =
  {x: T1 | ∃ (y: T2 | R(x, y)): Y(y)}

precondition(R)(Y): set[T1] = precondition(R, Y)

converse(R): pred[[T1]] =
  (λ (y: T2), (x: T1): R(x, y))

isomorphism?(R): bool =
  (∃ (f: (bijective?[T1, T2])): R = graph(f))

total?(R): bool = ∃ (x: T1): ∃ (y: T2): R(x, y)

onto?(R): bool = ∃ (y: T2): ∃ (x: T1): R(x, y)

END relation_defs

```

```

relation_props[T1, T2, T3: TYPE]: THEORY
BEGIN

  R1: VAR pred[[T2]]

  R2: VAR pred[[T3]]

  x: VAR T1

  y: VAR T2

  z: VAR T3

  R1 ∘ R2(x, z): bool = ∃ y: R1(x, y) AND R2(y, z)

  total_composition: LEMMA
    total?(R1) & total?(R2) ⇒ total?(R1 ∘ R2)

  onto_composition: LEMMA
    onto?(R1) & onto?(R2) ⇒ onto?(R1 ∘ R2)

  composition_total: JUDGEMENT O(R1: (total?[T1, T2]), R2: (total?[T2, T3])) HAS_TYPE
    (total?[T1, T3])

  composition_onto: JUDGEMENT O(R1: (onto?[T1, T2]), R2: (onto?[T2, T3])) HAS_TYPE
    (onto?[T1, T3])

END relation_props

```

```
relation_props2[T1, T2, T3, T4: TYPE]: THEORY
BEGIN

  R1: VAR pred[[T2]]

  R2: VAR pred[[T3]]

  R3: VAR pred[[T4]]

  ASSOC: LEMMA (R1 ∘ R2) ∘ R3 = R1 ∘ (R2 ∘ R3)

END relation_props2
```

```

relation_converse_props [T : TYPE] : THEORY
BEGIN

reflexive_converse : JUDGEMENT converse(R : (reflexive?[T])) HAS_TYPE
  (reflexive?[T])

irreflexive_converse : JUDGEMENT converse(R : (irreflexive?[T])) HAS_TYPE
  (irreflexive?[T])

symmetric_converse : JUDGEMENT converse(R : (symmetric?[T])) HAS_TYPE
  (symmetric?[T])

antisymmetric_converse : JUDGEMENT converse(R : (antisymmetric?[T])) HAS_TYPE
  (antisymmetric?[T])

connected_converse : JUDGEMENT converse(R : (connected?[T])) HAS_TYPE
  (connected?[T])

transitive_converse : JUDGEMENT converse(R : (transitive?[T])) HAS_TYPE
  (transitive?[T])

equivalence_converse : JUDGEMENT converse(R : (equivalence?[T])) HAS_TYPE
  (equivalence?[T])

preorder_converse : JUDGEMENT converse(R : (preorder?[T])) HAS_TYPE
  (preorder?[T])

partial_order_converse : JUDGEMENT converse(R : (partial_order?[T])) HAS_TYPE
  (partial_order?[T])

strict_order_converse : JUDGEMENT converse(R : (strict_order?[T])) HAS_TYPE
  (strict_order?[T])

dichotomous_converse : JUDGEMENT converse(R : (dichotomous?[T])) HAS_TYPE
  (dichotomous?[T])

total_order_converse : JUDGEMENT converse(R : (total_order?[T])) HAS_TYPE
  (total_order?[T])

trichotomous_converse : JUDGEMENT converse(R : (trichotomous?[T])) HAS_TYPE
  (trichotomous?[T])

```



```
strict_total_order_converse: JUDGEMENT converse( $R$ : (strict_total_order?[ $T$ ])) HAS_TYPE
  (strict_total_order?[ $T$ ])

END relation_converse_props
```

```

indexed_sets[index, T: TYPE]: THEORY
BEGIN

  i: VAR index

  x: VAR T

  A, B: VAR [index → set[T]]

  S: VAR set[T]

  ⋃ A: set[T] = {x | ∃ i: A(i)(x)}

  ⋂ A: set[T] = {x | ∀ i: A(i)(x)}

  IUnion_Union: LEMMA ⋃ A = ⋃ image(A)(fullset[index])

  IIntersection_Intersection: LEMMA
    ⋂ A = ⋂ image(A)(fullset[index])

  IUnion_union: LEMMA
    ⋃ λ i: (A(i) ∪ B(i)) = (⋃ A ∪ ⋃ B)

  IIntersection_intersection: LEMMA
    ⋂ λ i: (A(i) ∩ B(i)) = (⋂ A ∩ ⋂ B)

  IUnion_intersection: LEMMA
    ⋃ λ i: (A(i) ∩ S) = (⋃ A ∩ S)

  IIntersection_union: LEMMA
    ⋂ λ i: (A(i) ∪ S) = (⋂ A ∪ S)

END indexed_sets

```

```

operator_defs[T: TYPE]: THEORY
BEGIN

  O, ×: VAR [T, T → T]

  -: VAR [T → T]

  x, y, z: VAR T

  commutative?(O): bool =
    (∀ x, y: x ∘ y = y ∘ x)

  associative?(O): bool =
    (∀ x, y, z:
      (x ∘ y) ∘ z = x ∘ (y ∘ z))

  left_identity?(O)(y): bool = (∀ x: y ∘ x = x)

  right_identity?(O)(y): bool = (∀ x: x ∘ y = x)

  identity?(O)(y): bool =
    (∀ x: x ∘ y = x AND y ∘ x = x)

  has_identity?(O): bool = (∃ y: identity?(O)(y))

  zero?(O)(y): bool =
    (∀ x: x ∘ y = y AND y ∘ x = y)

  has_zero?(O): bool = (∃ y: zero?(O)(y))

  inverses?(O)(-)(y): bool =
    (∀ x: x ∘ -x = y AND (-x) ∘ x = y)

  has_inverses?(O): bool = (∃ -, y: inverses?(O)(-)(y))

  distributive?(×, O): bool =
    (∀ x, y, z:
      x × (y ∘ z) =
      (x × y) ∘ (x × z))

END operator_defs

```

```
numbers: THEORY
BEGIN

    number: TYPE+

END numbers
```

```

number_fields: THEORY
BEGIN

  number_field: TYPE+ FROM number

  numfield: TYPE+ = number_field

  number_field?(n: number): bool = number_field_pred(n)

  nonzero_number: TYPE+ = {r: number_field | r ≠ 0} CONTAINING 1

  nznum: TYPE+ = nonzero_number

  +, -, ×: [numfield, numfield → numfield]

  /: [numfield, nznum → numfield]

  -: [numfield → numfield]

  x, y, z: VAR numfield

  n0x: VAR nznum

  commutative_add: POSTULATE  $x + y = y + x$ 

  associative_add: POSTULATE
     $x + (y + z) = (x + y) + z$ 

  identity_add: POSTULATE  $x + 0 = x$ 

  inverse_add: AXIOM  $x + -x = 0$ 

  minus_add: AXIOM  $x - y = x + -y$ 

  commutative_mult: AXIOM  $x \times y = y \times x$ 

  associative_mult: AXIOM
     $x \times (y \times z) = (x \times y) \times z$ 

  identity_mult: AXIOM  $1 \times x = x$ 

  inverse_mult: AXIOM  $n0x \times (1/n0x) = 1$ 

```

div_def: AXIOM $y/n0x = y \times (1/n0x)$

distributive: POSTULATE

$$x \times (y + z) = (x \times y) + (x \times z)$$

END number_fields

reals: THEORY

BEGIN

\mathbb{R} : TYPE+ FROM number_field

real?(n : number): bool = number_field_pred(n) AND real_pred(n)

$\mathbb{R}_{\neq 0}$: TYPE+ = { r : \mathbb{R} | $r \neq 0$ } CONTAINING 1

$\mathbb{R}_{\neq 0}$: TYPE+ = $\mathbb{R}_{\neq 0}$

x, y : VAR \mathbb{R}

n0z: VAR $\mathbb{R}_{\neq 0}$

closed_plus: AXIOM real_pred($x + y$)

closed_minus: AXIOM real_pred($x - y$)

closed_times: AXIOM real_pred($x \times y$)

closed_divides: AXIOM real_pred($x/n0z$)

closed_neg: AXIOM real_pred($-x$)

real_plus_real_is_real: JUDGEMENT $+(x, y)$ HAS_TYPE \mathbb{R}

real_minus_real_is_real: JUDGEMENT $-(x, y)$ HAS_TYPE \mathbb{R}

real_times_real_is_real: JUDGEMENT $\times(x, y)$ HAS_TYPE \mathbb{R}

real_div_nzreal_is_real: JUDGEMENT $/ (x, n0z)$ HAS_TYPE \mathbb{R}

minus_real_is_real: JUDGEMENT $-(x)$ HAS_TYPE \mathbb{R}

$\langle(x, y)$: bool

$\leq(x, y)$: bool = $x < y$ OR $x = y$;

$\rangle(x, y)$: bool = $y < x$;

$\geq(x, y)$: bool = $y \leq x$

reals_totally_ordered: POSTULATE strict_total_order?(<)

END reals


```

real_axioms: THEORY
BEGIN

  x, y, z: VAR  $\mathbb{R}$ 

  n0x: VAR  $\mathbb{R}_{\neq 0}$ 

  posreal_add_closed: POSTULATE
    x > 0 AND y > 0 IMPLIES x + y > 0

  posreal_mult_closed: AXIOM x > 0 AND y > 0 IMPLIES x × y > 0

  posreal_neg: POSTULATE x > 0 IMPLIES NOT -x > 0

  trichotomy: POSTULATE x > 0 OR x = 0 OR 0 > x

END real_axioms

```

bounded_real_defs: THEORY

BEGIN

x, y : VAR \mathbb{R}

S : VAR (nonempty? $[\mathbb{R}]$)

upper_bound? (x, S) : bool = $\forall (s: (S)): s \leq x$

upper_bound? $(S)(x)$: bool = upper_bound? (x, S)

lower_bound? (x, S) : bool = $\forall (s: (S)): x \leq s$

lower_bound? $(S)(x)$: bool = lower_bound? (x, S)

least_upper_bound? (x, S) : bool =
upper_bound? (x, S) AND
 $\forall y$: upper_bound? (y, S) IMPLIES $(x \leq y)$

least_upper_bound? $(S)(x)$: bool = least_upper_bound? (x, S)

greatest_lower_bound? (x, S) : bool =
lower_bound? (x, S) AND
 $\forall y$: lower_bound? (y, S) IMPLIES $(y \leq x)$

greatest_lower_bound? $(S)(x)$: bool =
greatest_lower_bound? (x, S)

real_complete: AXIOM

$\forall S$:
 $(\exists y$: upper_bound? $(y, S))$ IMPLIES
 $(\exists y$: least_upper_bound? $(y, S))$

real_lower_complete: LEMMA

$\forall S$:
 $(\exists y$: lower_bound? $(y, S))$ IMPLIES
 $(\exists x$: greatest_lower_bound? $(x, S))$

bounded_above? (S) : bool = $(\exists x$: upper_bound? $(x, S))$

bounded_below? (S) : bool = $(\exists x$: lower_bound? $(x, S))$

```

bounded?(S): bool = bounded_above?(S) AND bounded_below?(S)

bounded_set: TYPE = (bounded?)

SA: VAR (bounded_above?)

lub_exists: LEMMA ( $\exists x$ : least_upper_bound?(x, SA))

lub(SA): {x | least_upper_bound?(x, SA)}

lub_lem: LEMMA lub(SA) = x IFF least_upper_bound?(x, SA)

SB: VAR (bounded_below?)

glb_exists: LEMMA ( $\exists x$ : greatest_lower_bound?(x, SB))

glb(SB): {x | greatest_lower_bound?(x, SB)}

glb_lem: LEMMA glb(SB) = x IFF greatest_lower_bound?(x, SB)

END bounded_real_defs

```

```

bounded_real_defs_alt[S: (nonempty?[ $\mathbb{R}$ ])]: THEORY
BEGIN

  x: VAR  $\mathbb{R}$ 

  upper_bound?: [ $\mathbb{R} \rightarrow \text{bool}$ ] = upper_bound?(S)

  lower_bound?: [ $\mathbb{R} \rightarrow \text{bool}$ ] = lower_bound?(S)

  least_upper_bound?: [ $\mathbb{R} \rightarrow \text{bool}$ ] = least_upper_bound?(S)

  greatest_lower_bound?: [ $\mathbb{R} \rightarrow \text{bool}$ ] = greatest_lower_bound?(S)

  lub_is_upper_bound: JUDGEMENT (least_upper_bound?) SUBTYPE_OF
    (upper_bound?)

  glb_is_lower_bound: JUDGEMENT (greatest_lower_bound?) SUBTYPE_OF
    (lower_bound?)

END bounded_real_defs_alt

```

```

real_types: THEORY
BEGIN

  x: VAR ℝ

  ℝ≥0: TYPE+ = {x: ℝ | x ≥ 0} CONTAINING 0

  ℝ≤0: TYPE+ = {x: ℝ | x ≤ 0} CONTAINING 0

  ℝ>0: TYPE+ = {x: ℝ≥0 | x > 0} CONTAINING 1

  ℝ<0: TYPE+ = {x: ℝ≤0 | x < 0} CONTAINING -1

  ℝ≥0: TYPE = ℝ≥0

  ℝ≤0: TYPE = ℝ≤0

  posreal_is_nzreal: JUDGEMENT ℝ>0 SUBTYPE_OF ℝ≠0

  negreal_is_nzreal: JUDGEMENT ℝ<0 SUBTYPE_OF ℝ≠0

  nzx, nzy: VAR ℝ≠0

  px, py: VAR ℝ>0

  nx, ny: VAR ℝ<0

  nnx, nny: VAR ℝ≥0

  npx, npy: VAR ℝ≤0

  nonneg_real_add_closed: LEMMA nnx + nny ≥ 0

  nonpos_real_add_closed: LEMMA npx + npy ≤ 0

  negreal_add_closed: LEMMA nx + ny < 0

  nonneg_real_mult_closed: LEMMA nnx × nny ≥ 0

  nzreal_times_nzreal_is_nzreal: JUDGEMENT ×(nzx, nzy) HAS_TYPE ℝ≠0

  nzreal_div_nzreal_is_nzreal: JUDGEMENT /(nzx, nzy) HAS_TYPE ℝ≠0

```

```

minus_nzreal_is_nzreal: JUDGEMENT  $-(nzx)$  HAS_TYPE  $\mathbb{R}_{\neq 0}$ 

nreal_plus_nreal_is_nreal: JUDGEMENT  $+(nmx, nny)$  HAS_TYPE  $\mathbb{R}_{\geq 0}$ 

nreal_times_nreal_is_nreal: JUDGEMENT  $\times(nmx, nny)$  HAS_TYPE  $\mathbb{R}_{\geq 0}$ 

nreal_div_posreal_is_nreal: JUDGEMENT  $/(nmx, py)$  HAS_TYPE  $\mathbb{R}_{\geq 0}$ 

nreal_div_negreal_is_npreal: JUDGEMENT  $/(nmx, ny)$  HAS_TYPE  $\mathbb{R}_{\leq 0}$ 

npreal_plus_npreal_is_npreal: JUDGEMENT  $+(npx, npy)$  HAS_TYPE  $\mathbb{R}_{\leq 0}$ 

npreal_times_npreal_is_npreal: JUDGEMENT  $\times(npx, npy)$  HAS_TYPE  $\mathbb{R}_{\geq 0}$ 

npreal_div_posreal_is_npreal: JUDGEMENT  $/(npx, py)$  HAS_TYPE  $\mathbb{R}_{\leq 0}$ 

npreal_div_negreal_is_npreal: JUDGEMENT  $/(npx, ny)$  HAS_TYPE  $\mathbb{R}_{\geq 0}$ 

posreal_plus_nreal_is_posreal: JUDGEMENT  $+(px, nny)$  HAS_TYPE  $\mathbb{R}_{> 0}$ 

nreal_plus_posreal_is_posreal: JUDGEMENT  $+(nmx, py)$  HAS_TYPE  $\mathbb{R}_{> 0}$ 

posreal_times_posreal_is_posreal: JUDGEMENT  $\times(px, py)$  HAS_TYPE  $\mathbb{R}_{> 0}$ 

posreal_div_posreal_is_posreal: JUDGEMENT  $/(px, py)$  HAS_TYPE  $\mathbb{R}_{> 0}$ 

negreal_plus_negreal_is_negreal: JUDGEMENT  $+(nx, ny)$  HAS_TYPE  $\mathbb{R}_{< 0}$ 

negreal_times_negreal_is_posreal: JUDGEMENT  $\times(nx, ny)$  HAS_TYPE  $\mathbb{R}_{> 0}$ 

negreal_div_negreal_is_posreal: JUDGEMENT  $/(nx, ny)$  HAS_TYPE  $\mathbb{R}_{> 0}$ 

END real_types

```

```

rationals: THEORY
BEGIN

  ℚ: TYPE+ FROM ℝ

  ℚ: TYPE+ = ℚ

  rational?(n: number): bool =
    number_field_pred(n) AND real_pred(n) AND rational_pred(n)

  ℚ≠0: TYPE+ = {r: ℚ | r ≠ 0} CONTAINING 1

  ℚ≠0: TYPE+ = ℚ≠0

  x, y: VAR ℚ

  n0z: VAR ℚ≠0

  closed_plus: AXIOM rational_pred(x + y)

  closed_minus: AXIOM rational_pred(x - y)

  closed_times: AXIOM rational_pred(x × y)

  closed_divides: AXIOM rational_pred(x/n0z)

  closed_neg: AXIOM rational_pred(-x)

  rat_plus_rat_is_rat: JUDGEMENT +(x, y) HAS_TYPE ℚ

  rat_minus_rat_is_rat: JUDGEMENT -(x, y) HAS_TYPE ℚ

  rat_times_rat_is_rat: JUDGEMENT ×(x, y) HAS_TYPE ℚ

  rat_div_nzrat_is_rat: JUDGEMENT /(x, n0z) HAS_TYPE ℚ

  minus_rat_is_rat: JUDGEMENT -(x) HAS_TYPE ℚ

  ℚ≥0: TYPE+ = {r: ℚ | r ≥ 0} CONTAINING 0

  ℚ≤0: TYPE+ = {r: ℚ | r ≤ 0} CONTAINING 0

```

$\mathbb{Q}_{>0}$: TYPE+ = $\{r: \mathbb{Q}_{\geq 0} \mid r > 0\}$ CONTAINING 1
 $\mathbb{Q}_{<0}$: TYPE+ = $\{r: \mathbb{Q}_{\leq 0} \mid r < 0\}$ CONTAINING -1
 $\mathbb{Q}_{\geq 0}$: TYPE+ = $\mathbb{Q}_{\geq 0}$
 $\mathbb{Q}_{\leq 0}$: TYPE+ = $\mathbb{Q}_{\leq 0}$
nnx, nny: VAR $\mathbb{Q}_{\geq 0}$
npx, npy: VAR $\mathbb{Q}_{\leq 0}$
px, py: VAR $\mathbb{Q}_{>0}$
nx, ny: VAR $\mathbb{Q}_{<0}$
n0x, n0y: VAR $\mathbb{Q}_{\neq 0}$
posrat_is_nzrat: JUDGEMENT $\mathbb{Q}_{>0}$ SUBTYPE_OF $\mathbb{Q}_{\neq 0}$
negrat_is_nzrat: JUDGEMENT $\mathbb{Q}_{<0}$ SUBTYPE_OF $\mathbb{Q}_{\neq 0}$
nzrat_times_nzrat_is_nzrat: JUDGEMENT $\times(n0x, n0y)$ HAS_TYPE $\mathbb{Q}_{\neq 0}$
nzrat_div_nzrat_is_nzrat: JUDGEMENT $/ (n0x, n0y)$ HAS_TYPE $\mathbb{Q}_{\neq 0}$
minus_nzrat_is_nzrat: JUDGEMENT $-(n0x)$ HAS_TYPE $\mathbb{Q}_{\neq 0}$
nnrat_plus_nnrat_is_nnrat: JUDGEMENT $+(nnx, nny)$ HAS_TYPE $\mathbb{Q}_{\geq 0}$
nnrat_times_nnrat_is_nnrat: JUDGEMENT $\times(nnx, nny)$ HAS_TYPE $\mathbb{Q}_{\geq 0}$
nnrat_div_posrat_is_nnrat: JUDGEMENT $/ (nnx, py)$ HAS_TYPE $\mathbb{Q}_{\geq 0}$
nnrat_div_negrat_is_nprat: JUDGEMENT $/ (nnx, ny)$ HAS_TYPE $\mathbb{Q}_{\leq 0}$
nprat_plus_nprat_is_nprat: JUDGEMENT $+(npx, npy)$ HAS_TYPE $\mathbb{Q}_{\leq 0}$
nprat_times_nprat_is_nnrat: JUDGEMENT $\times(npx, npy)$ HAS_TYPE $\mathbb{Q}_{\geq 0}$
nprat_div_posrat_is_nprat: JUDGEMENT $/ (npx, py)$ HAS_TYPE $\mathbb{Q}_{\leq 0}$

$\text{nprat_div_negrat_is_nnrat: JUDGEMENT } /(\text{np}x, \text{ny}) \text{ HAS_TYPE } \mathbb{Q}_{\geq 0}$
 $\text{posrat_plus_nnrat_is_posrat: JUDGEMENT } +(\text{p}x, \text{nn}y) \text{ HAS_TYPE } \mathbb{Q}_{> 0}$
 $\text{nnrat_plus_posrat_is_posrat: JUDGEMENT } +(\text{nn}x, \text{p}y) \text{ HAS_TYPE } \mathbb{Q}_{> 0}$
 $\text{posrat_times_posrat_is_posrat: JUDGEMENT } \times(\text{p}x, \text{p}y) \text{ HAS_TYPE } \mathbb{Q}_{> 0}$
 $\text{posrat_div_posrat_is_posrat: JUDGEMENT } /(\text{p}x, \text{p}y) \text{ HAS_TYPE } \mathbb{Q}_{> 0}$
 $\text{negrat_plus_negrat_is_negrat: JUDGEMENT } +(\text{n}x, \text{n}y) \text{ HAS_TYPE } \mathbb{Q}_{< 0}$
 $\text{negrat_times_negrat_is_posrat: JUDGEMENT } \times(\text{n}x, \text{n}y) \text{ HAS_TYPE } \mathbb{Q}_{> 0}$
 $\text{negrat_div_negrat_is_posrat: JUDGEMENT } /(\text{n}x, \text{n}y) \text{ HAS_TYPE } \mathbb{Q}_{> 0}$

END rationals

```

integers: THEORY
BEGIN

 $\mathbb{Z}$ : TYPE+ FROM  $\mathbb{Q}$ 

 $\mathbb{Z}$ : TYPE+ =  $\mathbb{Z}$ 

integer?(n: number): bool =
  number_field_pred(n) AND
  real_pred(n) AND rational_pred(n) AND integer_pred(n)

nonzero_integer: TYPE+ = {i:  $\mathbb{Z}$  | i  $\neq$  0} CONTAINING 1

 $\mathbb{Z}_{\neq 0}$ : TYPE+ = nonzero_integer

i, j: VAR  $\mathbb{Z}$ 

n0i, n0j: VAR  $\mathbb{Z}_{\neq 0}$ 

closed_plus: AXIOM integer_pred(i + j)

closed_minus: AXIOM integer_pred(i - j)

closed_times: AXIOM integer_pred(i  $\times$  j)

closed_neg: AXIOM integer_pred(-i)

upfrom(i): TYPE+ = {s:  $\mathbb{Z}$  | s  $\geq$  i} CONTAINING i

above(i): TYPE+ = {s:  $\mathbb{Z}$  | s > i} CONTAINING i + 1

int_plus_int_is_int: JUDGEMENT +(i, j) HAS_TYPE  $\mathbb{Z}$ 

int_minus_int_is_int: JUDGEMENT -(i, j) HAS_TYPE  $\mathbb{Z}$ 

int_times_int_is_int: JUDGEMENT  $\times$ (i, j) HAS_TYPE  $\mathbb{Z}$ 

minus_int_is_int: JUDGEMENT -(i) HAS_TYPE  $\mathbb{Z}$ 

minus_nzint_is_nzint: JUDGEMENT -(n0i) HAS_TYPE  $\mathbb{Z}_{\neq 0}$ 

 $\mathbb{N}$ : TYPE+ = {i:  $\mathbb{Z}$  | i  $\geq$  0} CONTAINING 0

```

$\mathbb{Z}_{\neq 0}$: TYPE+ = $\{i: \mathbb{Z} \mid i \leq 0\}$ CONTAINING 0
 $\mathbb{N}_{>0}$: TYPE+ = $\{i: \mathbb{N} \mid i > 0\}$ CONTAINING 1
 $\mathbb{Z}_{<0}$: TYPE+ = $\{i: \mathbb{Z}_{\neq 0} \mid i < 0\}$ CONTAINING -1
 $\mathbb{N}_{>0}$: TYPE+ = $\mathbb{N}_{>0}$
nmi, nnj: VAR \mathbb{N}
npi, npj: VAR $\mathbb{Z}_{\neq 0}$
 π , pj: VAR $\mathbb{N}_{>0}$
ni, nj: VAR $\mathbb{Z}_{<0}$
posint_is_nzint: JUDGEMENT $\mathbb{N}_{>0}$ SUBTYPE_OF $\mathbb{Z}_{\neq 0}$
negint_is_nzint: JUDGEMENT $\mathbb{Z}_{<0}$ SUBTYPE_OF $\mathbb{Z}_{\neq 0}$
nzint_times_nzint_is_nzint: JUDGEMENT $\times(n0i, n0j)$ HAS_TYPE $\mathbb{Z}_{\neq 0}$
nnint_plus_nnint_is_nnint: JUDGEMENT $+(nmi, nnj)$ HAS_TYPE \mathbb{N}
nnint_times_nnint_is_nnint: JUDGEMENT $\times(nmi, nnj)$ HAS_TYPE \mathbb{N}
npint_plus_npint_is_npint: JUDGEMENT $+(npi, npj)$ HAS_TYPE $\mathbb{Z}_{\neq 0}$
npint_times_npint_is_nnint: JUDGEMENT $\times(npi, npj)$ HAS_TYPE \mathbb{N}
posint_plus_nnint_is_posint: JUDGEMENT $+(\pi, nnj)$ HAS_TYPE $\mathbb{N}_{>0}$
nnint_plus_posint_is_posint: JUDGEMENT $+(nmi, pj)$ HAS_TYPE $\mathbb{N}_{>0}$
posint_times_posint_is_posint: JUDGEMENT $\times(\pi, pj)$ HAS_TYPE $\mathbb{N}_{>0}$
negint_plus_negint_is_negint: JUDGEMENT $+(ni, nj)$ HAS_TYPE $\mathbb{Z}_{<0}$
negint_times_negint_is_posint: JUDGEMENT $\times(ni, nj)$ HAS_TYPE $\mathbb{N}_{>0}$
subrange(i, j): TYPE = $\{k: \mathbb{Z} \mid i \leq k \text{ AND } k \leq j\}$

```

even?(i): bool =  $\exists j: i = j \times 2$ 

odd?(i): bool =  $\exists j: i = j \times 2 + 1$ 

even_int: TYPE+ = (even?) CONTAINING 0

odd_int: TYPE+ = (odd?) CONTAINING 1

e1, e2: VAR even_int

o1, o2: VAR odd_int

odd_is_nzint: JUDGEMENT odd_int SUBTYPE_OF  $\mathbb{Z}_{\neq 0}$ 

even_plus_even_is_even: JUDGEMENT +(e1, e2) HAS_TYPE even_int

even_minus_even_is_even: JUDGEMENT -(e1, e2) HAS_TYPE even_int

odd_plus_odd_is_even: JUDGEMENT +(o1, o2) HAS_TYPE even_int

odd_minus_odd_is_even: JUDGEMENT -(o1, o2) HAS_TYPE even_int

odd_plus_even_is_odd: JUDGEMENT +(o1, e2) HAS_TYPE odd_int

odd_minus_even_is_odd: JUDGEMENT -(o1, e2) HAS_TYPE odd_int

even_plus_odd_is_odd: JUDGEMENT +(e1, o2) HAS_TYPE odd_int

even_minus_odd_is_odd: JUDGEMENT -(e1, o2) HAS_TYPE odd_int

even_times_int_is_even: JUDGEMENT  $\times(e1, i)$  HAS_TYPE even_int

int_times_even_is_even: JUDGEMENT  $\times(i, e2)$  HAS_TYPE even_int

odd_times_odd_is_odd: JUDGEMENT  $\times(o1, o2)$  HAS_TYPE odd_int

minus_even_is_even: JUDGEMENT -(e1) HAS_TYPE even_int

minus_odd_is_odd: JUDGEMENT -(o1) HAS_TYPE odd_int

END integers

```

naturalnumbers: THEORY

BEGIN

\mathbb{N} : TYPE = \mathbb{N}

\mathbb{N} : TYPE+ = \mathbb{N}

i, j, k : VAR \mathbb{N}

n : VAR $\mathbb{N}_{>0}$

upfrom_nat_is_nat: JUDGEMENT upfrom(i) SUBTYPE_OF \mathbb{N}

upfrom_posnat_is_posnat: JUDGEMENT upfrom(n) SUBTYPE_OF $\mathbb{N}_{>0}$

above_nat_is_posnat: JUDGEMENT above(i) SUBTYPE_OF $\mathbb{N}_{>0}$

subrange_nat_is_nat: JUDGEMENT subrange(i, j) SUBTYPE_OF \mathbb{N}

subrange_posnat_is_posnat: JUDGEMENT subrange(n, j) SUBTYPE_OF $\mathbb{N}_{>0}$

upto(i): TYPE+ = $\{s: \mathbb{N} \mid s \leq i\}$ CONTAINING i

below(i): TYPE = $\{s: \mathbb{N} \mid s < i\}$

succ(i): $\mathbb{N} = i + 1$

pred(i): $\mathbb{N} =$ IF $i > 0$ THEN $i - 1$ ELSE 0 ENDIF;

$\sim(i, j)$: $\mathbb{N} =$ IF $i > j$ THEN $i - j$ ELSE 0 ENDIF

wf_nat: AXIOM well_founded?($\lambda i, j: i < j$)

p : VAR pred[\mathbb{N}]

nat_induction: LEMMA

($p(0)$ AND ($\forall j: p(j)$ IMPLIES $p(j + 1)$)) IMPLIES
($\forall i: p(i)$)

NAT_induction: LEMMA

($\forall j: (\forall k: k < j$ IMPLIES $p(k)$) IMPLIES $p(j)$) IMPLIES
($\forall i: p(i)$)

```

even_nat: TYPE+ = {i | even?(i)} CONTAINING 0
even_posnat: TYPE+ = {n | even?(n)} CONTAINING 2
odd_posnat: TYPE+ = {n | odd?(n)} CONTAINING 1
even_negint: TYPE+ = {n:  $\mathbb{Z}_{<0}$  | even?(n)} CONTAINING -2
odd_negint: TYPE+ = {n:  $\mathbb{Z}_{<0}$  | odd?(n)} CONTAINING -1
x: VAR  $\mathbb{Z}$ 
even_or_odd: THEOREM even?(x) IFF NOT odd?(x)
odd_iff_not_even: LEMMA odd?(x) IFF NOT even?(x)
even_iff_not_odd: LEMMA even?(x) IFF NOT odd?(x)
odd_or_even_int: LEMMA odd?(x) OR even?(x)
odd_iff_even_succ: LEMMA odd?(x) IFF even?(x + 1)
even_iff_odd_succ: LEMMA even?(x) IFF odd?(x + 1)
even_plus1: LEMMA even?(x) IFF NOT even?(x + 1)
odd_plus1: LEMMA odd?(x) IFF NOT odd?(x + 1)
even_div2: LEMMA even?(x) IMPLIES integer_pred(x/2)
odd_div2: LEMMA odd?(x) IMPLIES integer_pred((x - 1)/2)
END naturalnumbers

```

```

min_nat[T: TYPE FROM ℕ]: THEORY
BEGIN

  S: VAR (nonempty?[T])

  a, x: VAR T

  min(S): {a | S(a) AND (∀ x: S(x) IMPLIES a ≤ x)}

  minimum?(a, S): bool = S(a) AND (∀ x: S(x) IMPLIES a ≤ x)

  min_def: LEMMA
    ∀ (S: (nonempty?[T])): min(S) = a IFF minimum?(a, S)

END min_nat

```

real_defs: THEORY

BEGIN

m, n : VAR \mathbb{R}

$\text{sgn}(m)$: $\mathbb{Z} =$ IF $m \geq 0$ THEN 1 ELSE -1 ENDIF

$|m|$: $\{n: \mathbb{R}_{\geq 0} \mid n \geq m \text{ AND } n \geq -m\} =$
IF $m < 0$ THEN $-m$ ELSE m ENDIF

nonzero_abs_is_pos: JUDGEMENT $\text{abs}(x: \mathbb{R}_{\neq 0})$ HAS_TYPE
 $\{y: \mathbb{R}_{> 0} \mid y \geq x\}$

rat_abs_is_nonneg: JUDGEMENT $\text{abs}(q: \mathbb{Q})$ HAS_TYPE $\{r: \mathbb{Q}_{\geq 0} \mid r \geq q\}$

nzrat_abs_is_pos: JUDGEMENT $\text{abs}(q: \mathbb{Q}_{\neq 0})$ HAS_TYPE $\{r: \mathbb{Q}_{> 0} \mid r \geq q\}$

int_abs_is_nonneg: JUDGEMENT $\text{abs}(i: \mathbb{Z})$ HAS_TYPE $\{j: \mathbb{N} \mid j \geq i\}$

nzint_abs_is_pos: JUDGEMENT $\text{abs}(i: \mathbb{Z}_{\neq 0})$ HAS_TYPE $\{j: \mathbb{N}_{> 0} \mid j \geq i\}$

$\text{max}(m, n)$: $\{p: \mathbb{R} \mid p \geq m \text{ AND } p \geq n\} =$
IF $m < n$ THEN n ELSE m ENDIF

$\text{min}(m, n)$: $\{p: \mathbb{R} \mid p \leq m \text{ AND } p \leq n\} =$
IF $m > n$ THEN n ELSE m ENDIF

nzreal_max: JUDGEMENT $\text{max}(x, y: \mathbb{R}_{\neq 0})$ HAS_TYPE
 $\{z: \mathbb{R}_{\neq 0} \mid z \geq x \text{ AND } z \geq y\}$

nzreal_min: JUDGEMENT $\text{min}(x, y: \mathbb{R}_{\neq 0})$ HAS_TYPE
 $\{z: \mathbb{R}_{\neq 0} \mid z \leq x \text{ AND } z \leq y\}$

nonneg_real_max: JUDGEMENT $\text{max}(x, y: \mathbb{R}_{\geq 0})$ HAS_TYPE
 $\{z: \mathbb{R}_{\geq 0} \mid z \geq x \text{ AND } z \geq y\}$

nonneg_real_min: JUDGEMENT $\text{min}(x, y: \mathbb{R}_{\geq 0})$ HAS_TYPE
 $\{z: \mathbb{R}_{\geq 0} \mid z \leq x \text{ AND } z \leq y\}$

posreal_max: JUDGEMENT $\text{max}(x, y: \mathbb{R}_{> 0})$ HAS_TYPE
 $\{z: \mathbb{R}_{> 0} \mid z \geq x \text{ AND } z \geq y\}$

$\text{posreal_min: JUDGEMENT } \min(x, y: \mathbb{R}_{>0}) \text{ HAS_TYPE}$
 $\{z: \mathbb{R}_{>0} \mid z \leq x \text{ AND } z \leq y\}$

$\text{rat_max: JUDGEMENT } \max(q, r: \mathbb{Q}) \text{ HAS_TYPE}$
 $\{s: \mathbb{Q} \mid s \geq q \text{ AND } s \geq r\}$

$\text{rat_min: JUDGEMENT } \min(q, r: \mathbb{Q}) \text{ HAS_TYPE}$
 $\{s: \mathbb{Q} \mid s \leq q \text{ AND } s \leq r\}$

$\text{nzrat_max: JUDGEMENT } \max(q, r: \mathbb{Q}_{\neq 0}) \text{ HAS_TYPE}$
 $\{s: \mathbb{Q}_{\neq 0} \mid s \geq q \text{ AND } s \geq r\}$

$\text{nzrat_min: JUDGEMENT } \min(q, r: \mathbb{Q}_{\neq 0}) \text{ HAS_TYPE}$
 $\{s: \mathbb{Q}_{\neq 0} \mid s \leq q \text{ AND } s \leq r\}$

$\text{nonneg_rat_max: JUDGEMENT } \max(q, r: \mathbb{Q}_{\geq 0}) \text{ HAS_TYPE}$
 $\{s: \mathbb{Q}_{\geq 0} \mid s \geq q \text{ AND } s \geq r\}$

$\text{nonneg_rat_min: JUDGEMENT } \min(q, r: \mathbb{Q}_{\geq 0}) \text{ HAS_TYPE}$
 $\{s: \mathbb{Q}_{\geq 0} \mid s \leq q \text{ AND } s \leq r\}$

$\text{posrat_max: JUDGEMENT } \max(q, r: \mathbb{Q}_{>0}) \text{ HAS_TYPE}$
 $\{s: \mathbb{Q}_{>0} \mid s \geq q \text{ AND } s \geq r\}$

$\text{posrat_min: JUDGEMENT } \min(q, r: \mathbb{Q}_{>0}) \text{ HAS_TYPE}$
 $\{s: \mathbb{Q}_{>0} \mid s \leq q \text{ AND } s \leq r\}$

$\text{int_max: JUDGEMENT } \max(i, j: \mathbb{Z}) \text{ HAS_TYPE}$
 $\{k: \mathbb{Z} \mid i \leq k \text{ AND } j \leq k\}$

$\text{int_min: JUDGEMENT } \min(i, j: \mathbb{Z}) \text{ HAS_TYPE}$
 $\{k: \mathbb{Z} \mid k \leq i \text{ AND } k \leq j\}$

$\text{nzint_max: JUDGEMENT } \max(i, j: \mathbb{Z}_{\neq 0}) \text{ HAS_TYPE}$
 $\{k: \mathbb{Z}_{\neq 0} \mid i \leq k \text{ AND } j \leq k\}$

$\text{nzint_min: JUDGEMENT } \min(i, j: \mathbb{Z}_{\neq 0}) \text{ HAS_TYPE}$
 $\{k: \mathbb{Z}_{\neq 0} \mid k \leq i \text{ AND } k \leq j\}$

$\text{nat_max: JUDGEMENT } \max(i, j: \mathbb{N}) \text{ HAS_TYPE}$
 $\{k: \mathbb{N} \mid i \leq k \text{ AND } j \leq k\}$

nat_min: JUDGEMENT $\min(i, j: \mathbb{N})$ HAS_TYPE
 $\{k: \mathbb{N} \mid k \leq i \text{ AND } k \leq j\}$

posint_max: JUDGEMENT $\max(i, j: \mathbb{N}_{>0})$ HAS_TYPE
 $\{k: \mathbb{N}_{>0} \mid i \leq k \text{ AND } j \leq k\}$

posint_min: JUDGEMENT $\min(i, j: \mathbb{N}_{>0})$ HAS_TYPE
 $\{k: \mathbb{N}_{>0} \mid k \leq i \text{ AND } k \leq j\}$

$a, b, c: \text{VAR } \mathbb{R}$

min_le: LEMMA $\min(a, b) \leq c$ IFF $(a \leq c \text{ OR } b \leq c)$

min_lt: LEMMA $\min(a, b) < c$ IFF $(a < c \text{ OR } b < c)$

min_ge: LEMMA $\min(a, b) \geq c$ IFF $(a \geq c \text{ AND } b \geq c)$

min_gt: LEMMA $\min(a, b) > c$ IFF $(a > c \text{ AND } b > c)$

le_min: LEMMA $a \leq \min(b, c)$ IFF $(a \leq b \text{ AND } a \leq c)$

lt_min: LEMMA $a < \min(b, c)$ IFF $(a < b \text{ AND } a < c)$

ge_min: LEMMA $a \geq \min(b, c)$ IFF $(a \geq b \text{ OR } a \geq c)$

gt_min: LEMMA $a > \min(b, c)$ IFF $(a > b \text{ OR } a > c)$

max_le: LEMMA $\max(a, b) \leq c$ IFF $(a \leq c \text{ AND } b \leq c)$

max_lt: LEMMA $\max(a, b) < c$ IFF $(a < c \text{ AND } b < c)$

max_ge: LEMMA $\max(a, b) \geq c$ IFF $(a \geq c \text{ OR } b \geq c)$

max_gt: LEMMA $\max(a, b) > c$ IFF $(a > c \text{ OR } b > c)$

le_max: LEMMA $a \leq \max(b, c)$ IFF $(a \leq b \text{ OR } a \leq c)$

lt_max: LEMMA $a < \max(b, c)$ IFF $(a < b \text{ OR } a < c)$

ge_max: LEMMA $a \geq \max(b, c)$ IFF $(a \geq b \text{ AND } a \geq c)$

gt_max: LEMMA $a > \max(b, c)$ IFF $(a > b \text{ AND } a > c)$

END real_defs

real_props: THEORY

BEGIN

w, x, y, z : VAR \mathbb{R}

$n0w, n0x, n0y, n0z$: VAR $\mathbb{R}_{\neq 0}$

nnw, nnx, nny, nnz : VAR $\mathbb{R}_{\geq 0}$

pw, px, py, pz : VAR $\mathbb{R}_{> 0}$

npw, npx, npy, npz : VAR $\mathbb{R}_{\leq 0}$

nw, nx, ny, nz : VAR $\mathbb{R}_{< 0}$

inv_ne_0: LEMMA $1/n0x \neq 0$

both_sides_plus1: LEMMA $(x + z = y + z)$ IFF $x = y$

both_sides_plus2: LEMMA $(z + x = z + y)$ IFF $x = y$

both_sides_minus1: LEMMA $(x - z = y - z)$ IFF $x = y$

both_sides_minus2: LEMMA $(z - x = z - y)$ IFF $x = y$

both_sides_times1: LEMMA

$(x \times n0z = y \times n0z)$ IFF $x = y$

both_sides_times2: LEMMA

$(n0z \times x = n0z \times y)$ IFF $x = y$

both_sides_div1: LEMMA $(x/n0z = y/n0z)$ IFF $x = y$

both_sides_div2: LEMMA

$(n0z/n0x = n0z/n0y)$ IFF $n0x = n0y$

times_plus: LEMMA

$(x + y) \times (z + w) =$
 $x \times z + x \times w + y \times z + y \times w$

times_div1: LEMMA

$x \times (y/n0z) = (x \times y)/n0z$

times_div2: LEMMA

$$(x/n0z) \times y = (x \times y)/n0z$$

div_times: LEMMA

$$(x/n0x) \times (y/n0y) = \\ (x \times y)/(n0x \times n0y)$$

div_eq_zero: LEMMA $x/n0z = 0$ IFF $x = 0$

div_simp: LEMMA $n0x/n0x = 1$

div_cancel1: LEMMA $n0z \times (x/n0z) = x$

div_cancel2: LEMMA $(x/n0z) \times n0z = x$

div_cancel3: LEMMA $x/n0z = y$ IFF $x = y \times n0z$

div_cancel4: LEMMA $x = y/n0z$ IFF $x \times n0z = y$

cross_mult: LEMMA

$$(x/n0x = y/n0y) \text{ IFF} \\ (x \times n0y = y \times n0x)$$

add_div: LEMMA

$$(x/n0x) + (y/n0y) = \\ (x \times n0y + y \times n0x)/(n0x \times n0y)$$

minus_div1: LEMMA

$$(x/n0x) - (y/n0y) = \\ (x \times n0y - y \times n0x)/(n0x \times n0y)$$

minus_div2: LEMMA

$$(x/n0x - y/n0x) = (x - y)/n0x$$

div_distributes: LEMMA

$$(x/n0z) + (y/n0z) = (x + y)/n0z$$

div_distributes_minus: LEMMA

$$(x/n0z) - (y/n0z) = (x - y)/n0z$$

div_div1: LEMMA

$$(x/(n0y/n0z)) = ((x \times n0z)/n0y)$$

div_div2: LEMMA

$$((x/n0y)/n0z) = (x/(n0y \times n0z))$$

idem_add_is_zero: LEMMA $x + x = x$ IMPLIES $x = 0$

zero_times1: LEMMA $0 \times x = 0$

zero_times2: LEMMA $x \times 0 = 0$

zero_times3: LEMMA $x \times y = 0$ IFF $x = 0$ OR $y = 0$

neg_times_neg: LEMMA $(-x) \times (-y) = x \times y$

zero_is_neg_zero: LEMMA $-0 = 0$

strict_lt: LEMMA strict_total_order?($<$)

trich_lt: LEMMA $x < y$ OR $x = y$ OR $y < x$

tri_unique_lt1: LEMMA $x < y$ IMPLIES ($x \neq y$ AND NOT ($y < x$))

tri_unique_lt2: LEMMA

$$x = y \text{ IMPLIES } (\text{NOT } (x < y) \text{ AND NOT } (y < x))$$

zero_not_lt_zero: LEMMA NOT $0 < 0$

neg_lt: LEMMA $0 < -x$ IFF $x < 0$

pos_times_lt: LEMMA

$$0 < x \times y \text{ IFF } (0 < x \text{ AND } 0 < y) \text{ OR } (x < 0 \text{ AND } y < 0)$$

neg_times_lt: LEMMA

$$x \times y < 0 \text{ IFF } (0 < x \text{ AND } y < 0) \text{ OR } (x < 0 \text{ AND } 0 < y)$$

quotient_pos_lt: FORMULA $0 < 1/n0x$ IFF $0 < n0x$

quotient_neg_lt: FORMULA $1/n0x < 0$ IFF $n0x < 0$

pos_div_lt: LEMMA
 $0 < x/n0y$ IFF
 $(0 < x \text{ AND } 0 < n0y) \text{ OR } (x < 0 \text{ AND } n0y < 0)$

neg_div_lt: LEMMA
 $x/n0y < 0$ IFF
 $(0 < x \text{ AND } n0y < 0) \text{ OR } (x < 0 \text{ AND } 0 < n0y)$

div_mult_pos_lt1: LEMMA
 $z/py < x$ IFF $z < x \times py$

div_mult_pos_lt2: LEMMA
 $x < z/py$ IFF $x \times py < z$

div_mult_neg_lt1: LEMMA
 $z/ny < x$ IFF $x \times ny < z$

div_mult_neg_lt2: LEMMA
 $x < z/ny$ IFF $z < x \times ny$

both_sides_plus_lt1: LEMMA
 $x + z < y + z$ IFF $x < y$

both_sides_plus_lt2: LEMMA
 $z + x < z + y$ IFF $x < y$

both_sides_minus_lt1: LEMMA
 $x - z < y - z$ IFF $x < y$

both_sides_minus_lt2: LEMMA
 $z - x < z - y$ IFF $y < x$

both_sides_times_pos_lt1: LEMMA
 $x \times pz < y \times pz$ IFF $x < y$

both_sides_times_pos_lt2: LEMMA
 $pz \times x < pz \times y$ IFF $x < y$

both_sides_times_neg_lt1: LEMMA
 $x \times nz < y \times nz$ IFF $y < x$

both_sides_times_neg_lt2: LEMMA

$$nz \times x < nz \times y \text{ IFF } y < x$$

both_sides_div_pos_lt1: LEMMA

$$x/pz < y/pz \text{ IFF } x < y$$

both_sides_div_pos_lt2: LEMMA

$$pz/px < pz/py \text{ IFF } py < px$$

both_sides_div_pos_lt3: LEMMA

$$nz/px < nz/py \text{ IFF } px < py$$

both_sides_div_neg_lt1: LEMMA

$$x/nz < y/nz \text{ IFF } y < x$$

both_sides_div_neg_lt2: LEMMA

$$pz/nx < pz/ny \text{ IFF } ny < nx$$

both_sides_div_neg_lt3: LEMMA

$$nz/nx < nz/ny \text{ IFF } nx < ny$$

lt_plus_lt1: LEMMA

$$x \leq y \text{ AND } z < w \text{ IMPLIES } x + z < y + w$$

lt_plus_lt2: LEMMA

$$x < y \text{ AND } z \leq w \text{ IMPLIES } x + z < y + w$$

lt_minus_lt1: LEMMA

$$x \leq y \text{ AND } w < z \text{ IMPLIES } x - z < y - w$$

lt_minus_lt2: LEMMA

$$x < y \text{ AND } w \leq z \text{ IMPLIES } x - z < y - w$$

lt_times_lt_pos1: LEMMA

$$px \leq y \text{ AND } nnz < w \text{ IMPLIES } px \times nnz < y \times w$$

lt_times_lt_pos2: LEMMA

$$nnx < y \text{ AND } pz \leq w \text{ IMPLIES } nnx \times pz < y \times w$$

lt_div_lt_pos1: LEMMA

$$px \leq y \text{ AND } pz < w \text{ IMPLIES } px/w < y/pz$$

lt_div_lt_pos2: LEMMA

$nx < y$ AND $pz \leq w$ IMPLIES $nx/w < y/pz$

lt_times_lt_neg1: LEMMA

$x \leq ny$ AND $z < npw$ IMPLIES $ny \times npw < x \times z$

lt_times_lt_neg2: LEMMA

$x < npy$ AND $z \leq nw$ IMPLIES $npy \times nw < x \times z$

lt_div_lt_neg1: LEMMA

$x \leq ny$ AND $z < nw$ IMPLIES $ny/z < x/nw$

lt_div_lt_neg2: LEMMA

$x < npy$ AND $z \leq nw$ IMPLIES $npy/z < x/nw$

total_le: LEMMA total_order?(\leq)

dich_le: LEMMA $x \leq y$ OR $y \leq x$

zero_le_zero: LEMMA $0 \leq 0$

neg_le: LEMMA $0 \leq -x$ IFF $x \leq 0$

pos_times_le: LEMMA

$0 \leq x \times y$ IFF
($0 \leq x$ AND $0 \leq y$) OR ($x \leq 0$ AND $y \leq 0$)

neg_times_le: LEMMA

$x \times y \leq 0$ IFF
($0 \leq x$ AND $y \leq 0$) OR ($x \leq 0$ AND $0 \leq y$)

quotient_pos_le: FORMULA $0 \leq 1/n0x$ IFF $0 \leq n0x$

quotient_neg_le: FORMULA $1/n0x \leq 0$ IFF $n0x \leq 0$

pos_div_le: LEMMA

$0 \leq x/n0y$ IFF
($0 \leq x$ AND $0 \leq n0y$) OR ($x \leq 0$ AND $n0y \leq 0$)

neg_div_le: LEMMA

$x/n0y \leq 0$ IFF
($0 \leq x$ AND $n0y \leq 0$) OR ($x \leq 0$ AND $0 \leq n0y$)

div_mult_pos_le1: LEMMA

$$z/py \leq x \text{ IFF } z \leq x \times py$$

div_mult_pos_le2: LEMMA

$$x \leq z/py \text{ IFF } x \times py \leq z$$

div_mult_neg_le1: LEMMA

$$z/ny \leq x \text{ IFF } x \times ny \leq z$$

div_mult_neg_le2: LEMMA

$$x \leq z/ny \text{ IFF } z \leq x \times ny$$

both_sides_plus_le1: LEMMA

$$x + z \leq y + z \text{ IFF } x \leq y$$

both_sides_plus_le2: LEMMA

$$z + x \leq z + y \text{ IFF } x \leq y$$

both_sides_minus_le1: LEMMA

$$x - z \leq y - z \text{ IFF } x \leq y$$

both_sides_minus_le2: LEMMA

$$z - x \leq z - y \text{ IFF } y \leq x$$

both_sides_times_pos_le1: LEMMA

$$x \times pz \leq y \times pz \text{ IFF } x \leq y$$

both_sides_times_pos_le2: LEMMA

$$pz \times x \leq pz \times y \text{ IFF } x \leq y$$

both_sides_times_neg_le1: LEMMA

$$x \times nz \leq y \times nz \text{ IFF } y \leq x$$

both_sides_times_neg_le2: LEMMA

$$nz \times x \leq nz \times y \text{ IFF } y \leq x$$

both_sides_div_pos_le1: LEMMA

$$x/pz \leq y/pz \text{ IFF } x \leq y$$

both_sides_div_pos_le2: LEMMA

$$pz/px \leq pz/py \text{ IFF } py \leq px$$

both_sides_div_pos_le3: LEMMA
 $nz/px \leq nz/py$ IFF $px \leq py$

both_sides_div_neg_le1: LEMMA
 $x/nz \leq y/nz$ IFF $y \leq x$

both_sides_div_neg_le2: LEMMA
 $pz/nx \leq pz/ny$ IFF $ny \leq nx$

both_sides_div_neg_le3: LEMMA
 $nz/nx \leq nz/ny$ IFF $nx \leq ny$

le_plus_le: LEMMA
 $x \leq y$ AND $z \leq w$ IMPLIES $x + z \leq y + w$

le_minus_le: LEMMA
 $x \leq y$ AND $w \leq z$ IMPLIES $x - z \leq y - w$

le_times_le_pos: LEMMA
 $nnx \leq y$ AND $nnz \leq w$ IMPLIES $nnx \times nnz \leq y \times w$

le_div_le_pos: LEMMA
 $nnx \leq y$ AND $pz \leq w$ IMPLIES $nnx/w \leq y/pz$

le_times_le_neg: LEMMA
 $x \leq npy$ AND $z \leq npw$ IMPLIES $npy \times npw \leq x \times z$

le_div_le_neg: LEMMA
 $x \leq npy$ AND $z \leq nw$ IMPLIES $npy/z \leq x/nw$

strict_gt: LEMMA strict_total_order?(>)

trich_gt: LEMMA $x > y$ OR $x = y$ OR $y > x$

tri_unique_gt1: LEMMA $x > y$ IMPLIES $(x \neq y$ AND NOT $(y > x))$

tri_unique_gt2: LEMMA
 $x = y$ IMPLIES (NOT $(x > y)$ AND NOT $(y > x))$

zero_not_gt_zero: LEMMA NOT $0 > 0$

neg_gt: LEMMA $0 > -x$ IFF $x > 0$

pos_times_gt: LEMMA

$$x \times y > 0 \text{ IFF} \\ (0 > x \text{ AND } 0 > y) \text{ OR } (x > 0 \text{ AND } y > 0)$$

neg_times_gt: LEMMA

$$0 > x \times y \text{ IFF} \\ (0 > x \text{ AND } y > 0) \text{ OR } (x > 0 \text{ AND } 0 > y)$$

quotient_pos_gt: FORMULA $1/n0x > 0 \text{ IFF } n0x > 0$

quotient_neg_gt: FORMULA $0 > 1/n0x \text{ IFF } 0 > n0x$

pos_div_gt: LEMMA

$$x/n0y > 0 \text{ IFF} \\ (0 > x \text{ AND } 0 > n0y) \text{ OR } (x > 0 \text{ AND } n0y > 0)$$

neg_div_gt: LEMMA

$$0 > x/n0y \text{ IFF} \\ (0 > x \text{ AND } n0y > 0) \text{ OR } (x > 0 \text{ AND } 0 > n0y)$$

both_sides_plus_gt1: LEMMA

$$x + z > y + z \text{ IFF } x > y$$

both_sides_plus_gt2: LEMMA

$$z + x > z + y \text{ IFF } x > y$$

both_sides_minus_gt1: LEMMA

$$x - z > y - z \text{ IFF } x > y$$

both_sides_minus_gt2: LEMMA

$$z - x > z - y \text{ IFF } y > x$$

both_sides_times_pos_gt1: LEMMA

$$x \times pz > y \times pz \text{ IFF } x > y$$

both_sides_times_pos_gt2: LEMMA

$$pz \times x > pz \times y \text{ IFF } x > y$$

both_sides_times_neg_gt1: LEMMA

$$x \times nz > y \times nz \text{ IFF } y > x$$

both_sides_times_neg_gt2: LEMMA

$$nz \times x > nz \times y \text{ IFF } y > x$$

both_sides_div_pos_gt1: LEMMA

$$x/pz > y/pz \text{ IFF } x > y$$

both_sides_div_pos_gt2: LEMMA

$$pz/px > pz/py \text{ IFF } py > px$$

both_sides_div_pos_gt3: LEMMA

$$nz/px > nz/py \text{ IFF } px > py$$

both_sides_div_neg_gt1: LEMMA

$$x/nz > y/nz \text{ IFF } y > x$$

both_sides_div_neg_gt2: LEMMA

$$pz/nx > pz/ny \text{ IFF } ny > nx$$

both_sides_div_neg_gt3: LEMMA

$$nz/nx > nz/ny \text{ IFF } nx > ny$$

gt_plus_gt1: LEMMA

$$x \geq y \text{ AND } z > w \text{ IMPLIES } x + z > y + w$$

gt_plus_gt2: LEMMA

$$x > y \text{ AND } z \geq w \text{ IMPLIES } x + z > y + w$$

gt_minus_gt1: LEMMA

$$x \geq y \text{ AND } w > z \text{ IMPLIES } x - z > y - w$$

gt_minus_gt2: LEMMA

$$x > y \text{ AND } w \geq z \text{ IMPLIES } x - z > y - w$$

gt_times_gt_pos1: LEMMA

$$x \geq py \text{ AND } z > nnw \text{ IMPLIES } x \times z > py \times nnw$$

gt_times_gt_pos2: LEMMA

$$x > nny \text{ AND } z \geq pw \text{ IMPLIES } x \times z > nny \times pw$$

gt_div_gt_pos1: LEMMA

$$x \geq py \text{ AND } z > pw \text{ IMPLIES } x/pw > py/z$$

gt_div_gt_pos2: LEMMA

$x > nny$ AND $z \geq pw$ IMPLIES $x/pw > nny/z$

gt_times_gt_neg1: LEMMA

$nx \geq y$ AND $npz > w$ IMPLIES $y \times w > nx \times npz$

gt_times_gt_neg2: LEMMA

$npz > y$ AND $nz \geq w$ IMPLIES $y \times w > npz \times nz$

gt_div_gt_neg1: LEMMA

$nx \geq y$ AND $nz > w$ IMPLIES $y/nz > nx/w$

gt_div_gt_neg2: LEMMA

$npz > y$ AND $nz \geq w$ IMPLIES $y/nz > npz/w$

total_ge: LEMMA total_order?(\geq)

dich_ge: LEMMA $x \geq y$ OR $y \geq x$

zero_ge_zero: LEMMA $0 \geq 0$

neg_ge: LEMMA $0 \geq -x$ IFF $x \geq 0$

pos_times_ge: LEMMA

$x \times y \geq 0$ IFF
($0 \geq x$ AND $0 \geq y$) OR ($x \geq 0$ AND $y \geq 0$)

neg_times_ge: LEMMA

$0 \geq x \times y$ IFF
($0 \geq x$ AND $y \geq 0$) OR ($x \geq 0$ AND $0 \geq y$)

quotient_pos_ge: FORMULA $1/n0x \geq 0$ IFF $n0x \geq 0$

quotient_neg_ge: FORMULA $0 \geq 1/n0x$ IFF $0 \geq n0x$

pos_div_ge: LEMMA

$x/n0y \geq 0$ IFF
($0 \geq x$ AND $0 \geq n0y$) OR ($x \geq 0$ AND $n0y \geq 0$)

neg_div_ge: LEMMA

$0 \geq x/n0y$ IFF
($0 \geq x$ AND $n0y \geq 0$) OR ($x \geq 0$ AND $0 \geq n0y$)

div_mult_pos_ge1: LEMMA

$$z/py \geq x \text{ IFF } z \geq x \times py$$

div_mult_pos_ge2: LEMMA

$$x \geq z/py \text{ IFF } x \times py \geq z$$

div_mult_neg_ge1: LEMMA

$$z/ny \geq x \text{ IFF } x \times ny \geq z$$

div_mult_neg_ge2: LEMMA

$$x \geq z/ny \text{ IFF } z \geq x \times ny$$

both_sides_plus_ge1: LEMMA

$$x + z \geq y + z \text{ IFF } x \geq y$$

both_sides_plus_ge2: LEMMA

$$z + x \geq z + y \text{ IFF } x \geq y$$

both_sides_minus_ge1: LEMMA

$$x - z \geq y - z \text{ IFF } x \geq y$$

both_sides_minus_ge2: LEMMA

$$z - x \geq z - y \text{ IFF } y \geq x$$

both_sides_times_pos_ge1: LEMMA

$$x \times pz \geq y \times pz \text{ IFF } x \geq y$$

both_sides_times_pos_ge2: LEMMA

$$pz \times x \geq pz \times y \text{ IFF } x \geq y$$

both_sides_times_neg_ge1: LEMMA

$$x \times nz \geq y \times nz \text{ IFF } y \geq x$$

both_sides_times_neg_ge2: LEMMA

$$nz \times x \geq nz \times y \text{ IFF } y \geq x$$

both_sides_div_pos_ge1: LEMMA

$$x/pz \geq y/pz \text{ IFF } x \geq y$$

both_sides_div_pos_ge2: LEMMA

$$pz/px \geq pz/py \text{ IFF } py \geq px$$

both_sides_div_pos_ge3: LEMMA
 $nz/px \geq nz/py$ IFF $px \geq py$

both_sides_div_neg_ge1: LEMMA
 $x/nz \geq y/nz$ IFF $y \geq x$

both_sides_div_neg_ge2: LEMMA
 $pz/nx \geq pz/ny$ IFF $ny \geq nx$

both_sides_div_neg_ge3: LEMMA
 $nz/nx \geq nz/ny$ IFF $nx \geq ny$

ge_plus_ge: LEMMA
 $x \geq y$ AND $z \geq w$ IMPLIES $x + z \geq y + w$

ge_minus_ge: LEMMA
 $x \geq y$ AND $w \geq z$ IMPLIES $x - z \geq y - w$

ge_times_ge_pos: LEMMA
 $x \geq nny$ AND $z \geq nnw$ IMPLIES $x \times z \geq nny \times nnw$

ge_div_ge_pos: LEMMA
 $x \geq nny$ AND $z \geq pw$ IMPLIES $x/pw \geq nny/z$

ge_times_ge_neg: LEMMA
 $npz \geq y$ AND $npz \geq w$ IMPLIES $y \times w \geq npz \times npz$

ge_div_ge_neg: LEMMA
 $npz \geq y$ AND $nz \geq w$ IMPLIES $y/nz \geq npz/w$

nonzero_times1: LEMMA $n0x \times y = 0$ IFF $y = 0$

nonzero_times2: LEMMA $x \times n0y = 0$ IFF $x = 0$

nonzero_times3: LEMMA $n0x \times n0y = 0$ IFF FALSE

eq1_gt: FORMULA $x > 1$ AND $x \times y = 1$ IMPLIES $y < 1$

eq1_ge: FORMULA $x \geq 1$ AND $x \times y = 1$ IMPLIES $y \leq 1$

eqm1_gt: FORMULA

$$x > 1 \text{ AND } x \times y = -1 \text{ IMPLIES } y > -1$$

eqm1_ge: FORMULA

$$x \geq 1 \text{ AND } x \times y = -1 \text{ IMPLIES } y \geq -1$$

eqm1_lt: FORMULA

$$x < -1 \text{ AND } x \times y = -1 \text{ IMPLIES } y < 1$$

eqm1_le: FORMULA

$$x \leq -1 \text{ AND } x \times y = -1 \text{ IMPLIES } y \leq 1$$

sqrt_1: LEMMA $x \times x = 1$ IFF $x = 1$ OR $x = -1$

sqrt_1_lt: LEMMA $x \times x < 1$ IMPLIES $|x| < 1$

sqrt_1_le: LEMMA $x \times x \leq 1$ IMPLIES $|x| \leq 1$

idem_mult: LEMMA $x \times x = x$ IFF $x = 0$ OR $x = 1$

i, j : VAR \mathbb{Z}

product_1: LEMMA

$$i \geq 0 \text{ AND } j \geq 0 \text{ AND } i \times j = 1 \text{ IMPLIES } \\ i = 1 \text{ AND } j = 1$$

product_m1: LEMMA

$$i \geq 0 \text{ AND } j \leq 0 \text{ AND } i \times j = -1 \text{ IMPLIES } \\ i = 1 \text{ AND } j = -1$$

triangle: LEMMA $|x + y| \leq |x| + |y|$

abs_mult: LEMMA $|x \times y| = |x| \times |y|$

abs_div: LEMMA $|x/\mathbf{n}0y| = |x|/|\mathbf{n}0y|$

abs_abs: LEMMA $||x|| = |x|$

abs_square: LEMMA $|x \times x| = x \times x$

abs_limits: LEMMA

$$-(|x| + |y|) \leq x + y \text{ AND } \\ x + y \leq |x| + |y|$$

```

axiom_of_archimedes: LEMMA  $\forall (x: \mathbb{R}): \exists (i: \mathbb{Z}): x < i$ 

archimedean: LEMMA
   $\forall (px: \mathbb{R}_{>0}): \exists (n: \mathbb{N}_{>0}): 1/n < px$ 

real_lt_is_strict_total_order: JUDGEMENT < HAS_TYPE
  (strict_total_order? $[\mathbb{R}]$ )

real_le_is_total_order: JUDGEMENT  $\leq$  HAS_TYPE (total_order? $[\mathbb{R}]$ )

real_gt_is_strict_total_order: JUDGEMENT > HAS_TYPE
  (strict_total_order? $[\mathbb{R}]$ )

real_ge_is_total_order: JUDGEMENT  $\geq$  HAS_TYPE (total_order? $[\mathbb{R}]$ )

END real_props

```

```

extra_real_props: THEORY
BEGIN

  w, x, y, z: VAR ℝ

  n0w, n0x, n0y, n0z: VAR ℝ≠0

  nnw, nnx, nny, nnz: VAR ℝ≥0

  pw, px, py, pz: VAR ℝ>0

  npw, npx, npy, npz: VAR ℝ≤0

  nw, nx, ny, nz: VAR ℝ<0

  pos_neg_split: LEMMA ∃ px, nx: n0x = px OR n0x = nx

  div_mult_pos_neg_lt1: LEMMA
    z/n0y < x IFF
    IF n0y > 0
      THEN z < x × n0y
    ELSE x × n0y < z
    ENDIF

  div_mult_pos_neg_lt2: LEMMA
    x < z/n0y IFF
    IF n0y > 0
      THEN x × n0y < z
    ELSE z < x × n0y
    ENDIF

  div_mult_pos_neg_le1: LEMMA
    z/n0y ≤ x IFF
    IF n0y > 0
      THEN z ≤ x × n0y
    ELSE x × n0y ≤ z
    ENDIF

  div_mult_pos_neg_le2: LEMMA
    x ≤ z/n0y IFF
    IF n0y > 0
      THEN x × n0y ≤ z

```

ELSE $z \leq x \times n0y$
ENDIF

div_mult_pos_neg_gt1: LEMMA

$z/n0y > x$ IFF
IF $n0y > 0$
THEN $z > x \times n0y$
ELSE $x \times n0y > z$
ENDIF

div_mult_pos_neg_gt2: LEMMA

$x > z/n0y$ IFF
IF $n0y > 0$
THEN $x \times n0y > z$
ELSE $z > x \times n0y$
ENDIF

div_mult_pos_neg_ge1: LEMMA

$z/n0y \geq x$ IFF
IF $n0y > 0$
THEN $z \geq x \times n0y$
ELSE $x \times n0y \geq z$
ENDIF

div_mult_pos_neg_ge2: LEMMA

$x \geq z/n0y$ IFF
IF $n0y > 0$
THEN $x \times n0y \geq z$
ELSE $z \geq x \times n0y$
ENDIF

both_sides_times_pos_neg_lt1: LEMMA

IF $n0z > 0$
THEN $x \times n0z < y \times n0z$
ELSE $y \times n0z < x \times n0z$
ENDIF
IFF $x < y$

both_sides_times_pos_neg_lt2: LEMMA

IF $n0z > 0$
THEN $n0z \times x < n0z \times y$
ELSE $n0z \times y < n0z \times x$

ENDIF
IFF $x < y$

both_sides_times_pos_neg_le1: LEMMA

IF $n0z > 0$
THEN $x \times n0z \leq y \times n0z$
ELSE $y \times n0z \leq x \times n0z$
ENDIF
IFF $x \leq y$

both_sides_times_pos_neg_le2: LEMMA

IF $n0z > 0$
THEN $n0z \times x \leq n0z \times y$
ELSE $n0z \times y \leq n0z \times x$
ENDIF
IFF $x \leq y$

both_sides_times_pos_neg_gt1: LEMMA

IF $n0z > 0$
THEN $x \times n0z > y \times n0z$
ELSE $y \times n0z > x \times n0z$
ENDIF
IFF $x > y$

both_sides_times_pos_neg_gt2: LEMMA

IF $n0z > 0$
THEN $n0z \times x > n0z \times y$
ELSE $n0z \times y > n0z \times x$
ENDIF
IFF $x > y$

both_sides_times_pos_neg_ge1: LEMMA

IF $n0z > 0$
THEN $x \times n0z \geq y \times n0z$
ELSE $y \times n0z \geq x \times n0z$
ENDIF
IFF $x \geq y$

both_sides_times_pos_neg_ge2: LEMMA

IF $n0z > 0$
THEN $n0z \times x \geq n0z \times y$
ELSE $n0z \times y \geq n0z \times x$

```

ENDIF
  IFF  $x \geq y$ 

both_sides_div_pos_neg_lt1: LEMMA
  IF  $n0z > 0$ 
    THEN  $x/n0z < y/n0z$ 
  ELSE  $y/n0z < x/n0z$ 
  ENDIF
  IFF  $x < y$ 

both_sides_div_pos_neg_lt2: LEMMA
  IF  $n0z > 0$ 
    THEN  $n0z/px < n0z/py$ 
  ELSE  $n0z/py < n0z/px$ 
  ENDIF
  IFF  $py < px$ 

both_sides_div_pos_neg_lt3: LEMMA
  IF  $n0z > 0$ 
    THEN  $n0z/nx < n0z/ny$ 
  ELSE  $n0z/ny < n0z/nx$ 
  ENDIF
  IFF  $ny < nx$ 

both_sides_div_pos_neg_le1: LEMMA
  IF  $n0z > 0$ 
    THEN  $x/n0z \leq y/n0z$ 
  ELSE  $y/n0z \leq x/n0z$ 
  ENDIF
  IFF  $x \leq y$ 

both_sides_div_pos_neg_le2: LEMMA
  IF  $n0z > 0$ 
    THEN  $n0z/px \leq n0z/py$ 
  ELSE  $n0z/py \leq n0z/px$ 
  ENDIF
  IFF  $py \leq px$ 

both_sides_div_pos_neg_le3: LEMMA
  IF  $n0z > 0$ 
    THEN  $n0z/nx \leq n0z/ny$ 
  ELSE  $n0z/ny \leq n0z/nx$ 

```

```

ENDIF
  IFF  $ny \leq nx$ 

both_sides_div_pos_neg_gt1: LEMMA
  IF  $n0z > 0$ 
    THEN  $x/n0z > y/n0z$ 
  ELSE  $y/n0z > x/n0z$ 
ENDIF
  IFF  $x > y$ 

both_sides_div_pos_neg_gt2: LEMMA
  IF  $n0z > 0$ 
    THEN  $n0z/px > n0z/py$ 
  ELSE  $n0z/py > n0z/px$ 
ENDIF
  IFF  $py > px$ 

both_sides_div_pos_neg_gt3: LEMMA
  IF  $n0z > 0$ 
    THEN  $n0z/nx > n0z/ny$ 
  ELSE  $n0z/ny > n0z/nx$ 
ENDIF
  IFF  $ny > nx$ 

both_sides_div_pos_neg_ge1: LEMMA
  IF  $n0z > 0$ 
    THEN  $x/n0z \geq y/n0z$ 
  ELSE  $y/n0z \geq x/n0z$ 
ENDIF
  IFF  $x \geq y$ 

both_sides_div_pos_neg_ge2: LEMMA
  IF  $n0z > 0$ 
    THEN  $n0z/px \geq n0z/py$ 
  ELSE  $n0z/py \geq n0z/px$ 
ENDIF
  IFF  $py \geq px$ 

both_sides_div_pos_neg_ge3: LEMMA
  IF  $n0z > 0$ 
    THEN  $n0z/nx \geq n0z/ny$ 
  ELSE  $n0z/ny \geq n0z/nx$ 

```

ENDIF
 IFF $ny \geq nx$

both_sides_times1_imp: LEMMA
 $x = y$ IMPLIES $x \times w = y \times w$

both_sides_times2_imp: LEMMA
 $x = y$ IMPLIES $w \times x = w \times y$

both_sides_times_pos_le1_imp: LEMMA
 $x \leq y$ IMPLIES $x \times nnw \leq y \times nnw$

both_sides_times_pos_le2_imp: LEMMA
 $x \leq y$ IMPLIES $nnw \times x \leq nnw \times y$

both_sides_times_neg_le1_imp: LEMMA
 $y \leq x$ IMPLIES $x \times npw \leq y \times npw$

both_sides_times_neg_le2_imp: LEMMA
 $y \leq x$ IMPLIES $npw \times x \leq npw \times y$

both_sides_times_pos_ge1_imp: LEMMA
 $x \geq y$ IMPLIES $x \times nnw \geq y \times nnw$

both_sides_times_pos_ge2_imp: LEMMA
 $x \geq y$ IMPLIES $nnw \times x \geq nnw \times y$

both_sides_times_neg_ge1_imp: LEMMA
 $y \geq x$ IMPLIES $x \times npw \geq y \times npw$

both_sides_times_neg_ge2_imp: LEMMA
 $y \geq x$ IMPLIES $npw \times x \geq npw \times y$

both_sides_times_pos_neg_le1_imp: LEMMA
 $x \leq y$ IMPLIES
 IF $w \geq 0$
 THEN $x \times w \leq y \times w$
 ELSE $y \times w \leq x \times w$
 ENDIF

both_sides_times_pos_neg_le2_imp: LEMMA
 $x \leq y$ IMPLIES


```

IF  $w \geq 0$ 
  THEN  $w \times x \leq w \times y$ 
ELSE  $w \times y \leq w \times x$ 
ENDIF

```

both_sides_times_pos_neg_ge1_imp: LEMMA

```

 $x \geq y$  IMPLIES
IF  $w \geq 0$ 
  THEN  $x \times w \geq y \times w$ 
ELSE  $y \times w \geq x \times w$ 
ENDIF

```

both_sides_times_pos_neg_ge2_imp: LEMMA

```

 $x \geq y$  IMPLIES
IF  $w \geq 0$ 
  THEN  $w \times x \geq w \times y$ 
ELSE  $w \times y \geq w \times x$ 
ENDIF

```

zero_times4: LEMMA $0 = x \times y$ IFF $x = 0$ OR $y = 0$

times_div_cancel1: LEMMA $(n0z \times x)/n0z = x$

times_div_cancel2: LEMMA $(x \times n0z)/n0z = x$

div_mult_pos_gt1: LEMMA $z/py > x$ IFF $z > x \times py$

div_mult_pos_gt2: LEMMA $x > z/py$ IFF $x \times py > z$

div_mult_neg_gt1: LEMMA $z/ny > x$ IFF $x \times ny > z$

div_mult_neg_gt2: LEMMA $x > z/ny$ IFF $z > x \times ny$

lt_cut: LEMMA $x < y$ AND $y < z$ IMPLIES $x < z$

le_cut: LEMMA $x \leq y$ AND $y \leq z$ IMPLIES $x \leq z$

gt_cut: LEMMA $x > y$ AND $y > z$ IMPLIES $x > z$

ge_cut: LEMMA $x \geq y$ AND $y \geq z$ IMPLIES $x \geq z$

le_times_le_any1: LEMMA

```

IF  $w \geq 0$  IFF  $x \geq 0$ 
  THEN IF  $y \geq 0$  IFF  $z \geq 0$ 
    THEN  $|w| \leq |y|$  AND  $|x| \leq |z|$ 
    ELSE  $(w = 0$  OR  $x = 0)$  AND  $(y = 0$  OR  $z = 0)$ 
    ENDIF
  ELSE IF  $y \geq 0$  IFF  $z \geq 0$ 
    THEN TRUE
    ELSE  $|w| \geq |y|$  AND  $|x| \geq |z|$ 
    ENDIF
ENDIF
IMPLIES  $w \times x \leq y \times z$ 

```

ge_times_ge_any1 : LEMMA

```

IF  $y \geq 0$  IFF  $z \geq 0$ 
  THEN IF  $w \geq 0$  IFF  $x \geq 0$ 
    THEN  $|w| \geq |y|$  AND  $|x| \geq |z|$ 
    ELSE  $(w = 0$  OR  $x = 0)$  AND  $(y = 0$  OR  $z = 0)$ 
    ENDIF
  ELSE IF  $w \geq 0$  IFF  $x \geq 0$ 
    THEN TRUE
    ELSE  $|w| \leq |y|$  AND  $|x| \leq |z|$ 
    ENDIF
ENDIF
IMPLIES  $w \times x \geq y \times z$ 

```

lt_times_lt_any1 : LEMMA

```

IF  $w = 0$  OR  $x = 0$ 
  THEN  $0 < y$  AND  $0 < z$  OR  $y < 0$  AND  $z < 0$ 
ELSIF  $w > 0$  IFF  $x > 0$ 
  THEN  $(y > 0$  IFF  $z > 0)$  AND
     $(|w| \leq |y|$  AND  $|x| < |z|$  OR
       $|w| < |y|$  AND  $|x| \leq |z|)$ 
ELSIF  $y > 0$  IFF  $z > 0$  THEN TRUE
ELSE  $|w| \geq |y|$  AND  $|x| > |z|$  OR
   $|w| > |y|$  AND  $|x| \geq |z|$ 
ENDIF
IMPLIES  $w \times x < y \times z$ 

```

gt_times_gt_any1 : LEMMA

```

IF  $y = 0$  OR  $z = 0$ 
  THEN  $0 < w$  AND  $0 < x$  OR  $w < 0$  AND  $x < 0$ 
ELSIF  $y > 0$  IFF  $z > 0$ 

```

```

    THEN ( $w > 0$  IFF  $x > 0$ ) AND
           ( $|w| \geq |y|$  AND  $|x| > |z|$  OR
             $|w| > |y|$  AND  $|x| \geq |z|$ )
  ELSIF  $w > 0$  IFF  $x > 0$  THEN TRUE
  ELSE  $|w| \leq |y|$  AND  $|x| < |z|$  OR
        $|w| < |y|$  AND  $|x| \leq |z|$ 
  ENDIF
  IMPLIES  $w \times x > y \times z$ 

```

le_times_le_any2: LEMMA

```

 $w \times x \leq y \times z$  IMPLIES
  IF  $y = 0$  OR  $z = 0$ 
    THEN ( $0 \geq w$  OR  $0 \geq x$ ) AND ( $w \geq 0$  OR  $x \geq 0$ )
  ELSIF  $y > 0$  IFF  $z > 0$ 
    THEN ( $w > 0$  IFF  $x > 0$ ) IMPLIES
           ( $|w| < |y|$  OR  $|x| \leq |z|$ ) AND
           ( $|w| \leq |y|$  OR  $|x| < |z|$ )
  ELSE NOT ( $w > 0$  IFF  $x > 0$ ) AND
           ( $|w| > |y|$  OR  $|x| \geq |z|$ ) AND
           ( $|w| \geq |y|$  OR  $|x| > |z|$ )
  ENDIF

```

ge_times_ge_any2: LEMMA

```

 $w \times x \geq y \times z$  IMPLIES
  IF  $w = 0$  OR  $x = 0$ 
    THEN ( $0 \geq y$  OR  $0 \geq z$ ) AND ( $y \geq 0$  OR  $z \geq 0$ )
  ELSIF  $w > 0$  IFF  $x > 0$ 
    THEN ( $y > 0$  IFF  $z > 0$ ) IMPLIES
           ( $|w| > |y|$  OR  $|x| \geq |z|$ ) AND
           ( $|w| \geq |y|$  OR  $|x| > |z|$ )
  ELSE NOT ( $y > 0$  IFF  $z > 0$ ) AND
           ( $|w| < |y|$  OR  $|x| \leq |z|$ ) AND
           ( $|w| \leq |y|$  OR  $|x| < |z|$ )
  ENDIF

```

lt_times_lt_any2: LEMMA

```

 $w \times x < y \times z$  IMPLIES
  IF  $y \geq 0$  IFF  $z \geq 0$ 
    THEN IF  $w \geq 0$  IFF  $x \geq 0$ 
           THEN  $|w| < |y|$  OR  $|x| < |z|$ 
         ELSE ( $w \neq 0$  AND  $x \neq 0$ ) OR ( $y \neq 0$  AND  $z \neq 0$ )
         ENDIF
  ENDIF

```

```

ELSE NOT ( $w \geq 0$  IFF  $x \geq 0$ ) AND
      ( $|w| > |y|$  OR  $|x| > |z|$ )
ENDIF

```

gt_times_gt_any2: LEMMA

```

 $w \times x > y \times z$  IMPLIES
  IF  $w \geq 0$  IFF  $x \geq 0$ 
    THEN IF  $y \geq 0$  IFF  $z \geq 0$ 
      THEN  $|w| > |y|$  OR  $|x| > |z|$ 
      ELSE ( $w \neq 0$  AND  $x \neq 0$ ) OR ( $y \neq 0$  AND  $z \neq 0$ )
      ENDIF
    ELSE NOT ( $y \geq 0$  IFF  $z \geq 0$ ) AND
      ( $|w| < |y|$  OR  $|x| < |z|$ )
    ENDIF
ENDIF

```

END extra_real_props

```

extra_tegies: THEORY
BEGIN

  x, y: VAR ℝ

  n0z: VAR ℝ≠0

  neg_mult: LEMMA  $-x \times y = -(x \times y)$ 

  mult_neg: LEMMA  $x \times -y = -(x \times y)$ 

  neg_add: LEMMA  $-x + y = y - x$ 

  add_neg: LEMMA  $x + -y = x - y$ 

  neg_div: LEMMA  $-x/n0z = -(x/n0z)$ 

  div_neg: LEMMA  $x/-n0z = -(x/n0z)$ 

  one_times: LEMMA  $1 \times x = x$ 

  neg_one_times: LEMMA  $-1 \times x = -x$ 

  zero_div: LEMMA  $0/n0z = 0$ 

  neg_neg: LEMMA  $--x = x$ 

END extra_tegies

```

```

rational_props: THEORY
BEGIN

  x, y: VAR ℝ

  i: VAR ℤ

  n0j: VAR ℤ≠0

  p: VAR ℕ>0

  r: VAR ℚ

  rational_pred_ax: AXIOM ∃ i, n0j: r = i/n0j

  rational_pred_ax2: LEMMA ∃ i, p: r = i/p

  density_positive: LEMMA
    0 ≤ x AND x < y IMPLIES (∃ r: x < r AND r < y)

  density: LEMMA x < y IMPLIES (∃ r: x < r AND r < y)

END rational_props

```

```

integer_props: THEORY
BEGIN

  m, n: VAR ℕ

  i, j, k: VAR ℤ

  n0j: VAR ℤ≠0

  N: VAR (nonempty?[ℕ])

  I: VAR (nonempty?[ℤ])

  integer_pred_ax: LEMMA ∃ n: i = n OR i = -n

  div_simple: LEMMA
    (∃ k: i = k × n0j) = integer_pred(i/n0j)

  lub_nat: LEMMA
    upper_bound?(m, extend[ℝ, ℕ, bool, FALSE](N)) ⇒
      ∃ (n: (N)):
        least_upper_bound?(n, extend[ℝ, ℕ, bool, FALSE](N))

  lub_int: LEMMA
    upper_bound?(i, extend[ℝ, ℤ, bool, FALSE](I)) ⇒
      ∃ (j: (I)):
        least_upper_bound?(j, extend[ℝ, ℤ, bool, FALSE](I))

  glb_nat: LEMMA
    ∃ (n: (N)):
      greatest_lower_bound?(n, extend[ℝ, ℕ, bool, FALSE](N))

  glb_int: LEMMA
    lower_bound?(i, extend[ℝ, ℤ, bool, FALSE](I)) ⇒
      ∃ (j: (I)):
        greatest_lower_bound?(j, extend[ℝ, ℤ, bool, FALSE](I))

END integer_props

```

```

floor_ceil: THEORY
BEGIN

  x, y: VAR ℝ

  py: VAR ℝ>0

  j: VAR nonzero_integer

  i, k: VAR ℤ

  floor_exists: LEMMA ∃ i: i ≤ x & x < i + 1

  ceiling_exists: LEMMA ∃ i: x ≤ i & i < x + 1

  [x]: {i | i ≤ x & x < i + 1}

  fractional(x): {x | 0 ≤ x & x < 1} = x - [x]

  ⌈x⌉: {i | x ≤ i & i < x + 1}

  floor_def: LEMMA [x] ≤ x & x < [x] + 1

  ceiling_def: LEMMA x ≤ ⌈x⌉ & ⌈x⌉ < x + 1

  floor_ceiling_reflect1: LEMMA [-x] = -⌈x⌉

  floor_ceiling_reflect2: LEMMA ⌈-x⌉ = -[x]

  nonneg_floor_is_nat: JUDGEMENT floor(x: ℝ≥0) HAS_TYPE ℕ

  nonneg_ceiling_is_nat: JUDGEMENT ceiling(x: ℝ≥0) HAS_TYPE ℕ

  floor_int: LEMMA [i] = i

  ceiling_int: LEMMA ⌈i⌉ = i

  floor_plus_int: LEMMA [x + i] = [x] + i

  ceiling_plus_int: LEMMA ⌈x + i⌉ = ⌈x⌉ + i

  floor_ceiling_nonint: LEMMA

```


$$\text{NOT integer?}(x) \Rightarrow \lceil x \rceil - \lfloor x \rfloor = 1$$

$$\text{floor_ceiling_int: LEMMA } \lfloor i \rfloor = \lceil i \rceil$$

floor_neg: LEMMA

$$\begin{aligned} \lfloor x \rfloor &= \\ \text{IF integer?}(x) & \\ \text{THEN } -\lfloor -x \rfloor & \\ \text{ELSE } -\lfloor -x \rfloor - 1 & \\ \text{ENDIF} & \end{aligned}$$

$$\text{real_parts: LEMMA } x = \lfloor x \rfloor + \text{fractional}(x)$$

floor_plus: LEMMA

$$\begin{aligned} \lfloor x + y \rfloor &= \\ \lfloor x \rfloor + \lfloor y \rfloor + \lfloor \text{fractional}(x) + \text{fractional}(y) \rfloor & \end{aligned}$$

ceiling_plus: LEMMA

$$\begin{aligned} \lceil x + y \rceil &= \\ \lfloor x \rfloor + \lfloor y \rfloor + \lceil \text{fractional}(x) + \text{fractional}(y) \rceil & \end{aligned}$$

$$\text{floor_split: LEMMA } i = \lfloor i/2 \rfloor + \lceil i/2 \rceil$$

$$\text{floor_within_1: LEMMA } x - \lfloor x \rfloor < 1$$

$$\text{ceiling_within_1: LEMMA } \lceil x \rceil - x < 1$$

floor_val: LEMMA

$$\begin{aligned} i \geq k \times j \text{ AND } i < (k + 1) \times j \text{ IMPLIES} \\ \lfloor i/j \rfloor = k \end{aligned}$$

floor_small: LEMMA

$$\begin{aligned} |i| < |j| \text{ IMPLIES} \\ \lfloor i/j \rfloor = \\ \text{IF } i/j \geq 0 \text{ THEN } 0 \text{ ELSE } -1 \text{ ENDIF} \end{aligned}$$

$$\text{floor_eq_0: LEMMA } \lfloor x \rfloor = 0 \text{ IMPLIES } x \geq 0 \text{ AND } x < 1$$

fractional_plus: LEMMA

$$\begin{aligned} \text{fractional}(x + y) &= \\ \text{fractional}(\text{fractional}(x) + \text{fractional}(y)) & \end{aligned}$$

floor_div: LEMMA

$$\lfloor x/\text{py} \rfloor = i \text{ IFF}$$

$$\text{py} \times i \leq x \text{ AND } x < \text{py} \times (i + 1)$$

floor_0: LEMMA $\lfloor x \rfloor = 0$ IFF $0 \leq x$ AND $x < 1$

END floor_ceil

exponentiation: THEORY

BEGIN

r : VAR \mathbb{R}

q : VAR \mathbb{Q}

nnq: VAR $\mathbb{Q}_{\geq 0}$

m, n : VAR \mathbb{N}

pm, pn: VAR $\mathbb{N}_{>0}$

i, j : VAR \mathbb{Z}

n0i, n0j: VAR $\mathbb{Z}_{\neq 0}$

x, y : VAR \mathbb{R}

px, py: VAR $\mathbb{R}_{>0}$

n0x, n0y: VAR $\mathbb{R}_{\neq 0}$

gtlx, gtly: VAR $\{r: \mathbb{R}_{>0} \mid r > 1\}$

ltlx, ltly: VAR $\{r: \mathbb{R}_{>0} \mid r < 1\}$

gely, gely: VAR $\{r: \mathbb{R}_{>0} \mid r \geq 1\}$

lelx, lely: VAR $\{r: \mathbb{R}_{>0} \mid r \leq 1\}$

r^n : RECURSIVE $\mathbb{R} = \text{IF } n = 0 \text{ THEN } 1 \text{ ELSE } r \times r^{n-1} \text{ ENDIF}$
MEASURE n ;

expt_pos_aux: LEMMA $px^n > 0$

expt_nonzero_aux: LEMMA $n0x^n \neq 0$;

nnreal_expt: JUDGEMENT $\text{expt}(x: \mathbb{R}_{\geq 0}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{R}_{\geq 0}$

posreal_expt: JUDGEMENT $\text{expt}(x: \mathbb{R}_{>0}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{R}_{>0}$

$\text{nzreal_expt: JUDGEMENT } \text{expt}(x: \mathbb{R}_{\neq 0}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{R}_{\neq 0}$
 $\text{rat_expt: JUDGEMENT } \text{expt}(x: \mathbb{Q}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{Q}$
 $\text{nnrat_expt: JUDGEMENT } \text{expt}(x: \mathbb{Q}_{\geq 0}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{Q}_{\geq 0}$
 $\text{posrat_expt: JUDGEMENT } \text{expt}(x: \mathbb{Q}_{> 0}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{Q}_{> 0}$
 $\text{int_expt: JUDGEMENT } \text{expt}(x: \mathbb{Z}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{Z}$
 $\text{nat_expt: JUDGEMENT } \text{expt}(x: \mathbb{N}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{N}$
 $\text{posnat_expt: JUDGEMENT } \text{expt}(x: \mathbb{N}_{> 0}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{N}_{> 0}$
 $r: \mathbb{R}^i: \{i: \mathbb{Z} \mid r \neq 0 \text{ OR } i \geq 0\}; \mathbb{R} =$
 $\quad \text{IF } i \geq 0 \text{ THEN } r^i \text{ ELSE } 1/r^{-i} \text{ ENDIF}$
 $\text{expt_pos: LEMMA } \text{px}^i > 0$
 $\text{expt_nonzero: LEMMA } \text{n0x}^i \neq 0$
 $\text{nnreal_exp: JUDGEMENT } \wedge(x: \mathbb{R}_{\geq 0}, i: \mathbb{Z} \mid x \neq 0 \text{ OR } i \geq 0) \text{ HAS_TYPE}$
 $\quad \mathbb{R}_{\geq 0}$
 $\text{posreal_exp: JUDGEMENT } \wedge(x: \mathbb{R}_{> 0}, i: \mathbb{Z}) \text{ HAS_TYPE } \mathbb{R}_{> 0}$
 $\text{nzreal_exp: JUDGEMENT } \wedge(x: \mathbb{R}_{\neq 0}, i: \mathbb{Z}) \text{ HAS_TYPE } \mathbb{R}_{\neq 0}$
 $\text{rat_exp: JUDGEMENT } \wedge(x: \mathbb{Q}, i: \mathbb{Z} \mid x \neq 0 \text{ OR } i \geq 0) \text{ HAS_TYPE}$
 $\quad \mathbb{Q}$
 $\text{nnrat_exp: JUDGEMENT } \wedge(x: \mathbb{Q}_{\geq 0}, i: \mathbb{Z} \mid x \neq 0 \text{ OR } i \geq 0) \text{ HAS_TYPE}$
 $\quad \mathbb{Q}_{\geq 0}$
 $\text{posrat_exp: JUDGEMENT } \wedge(x: \mathbb{Q}_{> 0}, i: \mathbb{Z}) \text{ HAS_TYPE } \mathbb{Q}_{> 0}$
 $\text{int_exp: JUDGEMENT } \wedge(x: \mathbb{Z}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{Z}$
 $\text{nat_exp: JUDGEMENT } \wedge(x: \mathbb{N}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{N}$
 $\text{posint_exp: JUDGEMENT } \wedge(x: \mathbb{N}_{> 0}, n: \mathbb{N}) \text{ HAS_TYPE } \mathbb{N}_{> 0}$

expt_x0_aux: LEMMA $x^0 = 1$

expt_x1_aux: LEMMA $x^1 = x$

expt_1n_aux: LEMMA $1^n = 1$

increasing_expt_aux: LEMMA $gt1x^{m+2} > gt1x$

decreasing_expt_aux: LEMMA $lt1x^{m+2} < lt1x$

expt_1_aux: LEMMA $px^{n+1} = 1$ IFF $px = 1$

expt_plus_aux: LEMMA
 $n0x^{m+n} = n0x^m \times n0x^n$

expt_minus_aux: LEMMA
 $m \geq n$ IMPLIES $n0x^{m-n} = n0x^m / n0x^n$

expt_times_aux: LEMMA $n0x^{m \times n} = n0x^{mn}$

expt_divide_aux: LEMMA
 $1/n0x^{m \times n} = 1/n0x^{mn}$

both_sides_expt1_aux: LEMMA
 $px^{m+1} = px^{n+1}$ IFF $m = n$ OR $px = 1$

both_sides_expt2_aux: LEMMA $px^{pm} = py^{pm}$ IFF $px = py$

both_sides_expt_pos_lt_aux: LEMMA
 $px^{m+1} < py^{m+1}$ IFF $px < py$

both_sides_expt_gt1_lt_aux: LEMMA
 $gt1x^{m+1} < gt1x^{n+1}$ IFF $m < n$

both_sides_expt_lt1_lt_aux: LEMMA
 $lt1x^{m+1} < lt1x^{n+1}$ IFF $n < m$

both_sides_expt_pos_le_aux: LEMMA
 $px^{m+1} \leq py^{m+1}$ IFF $px \leq py$

both_sides_expt_gt1_le_aux: LEMMA
 $gt1x^{m+1} \leq gt1x^{n+1}$ IFF $m \leq n$

both_sides_expt_lt1_le_aux: LEMMA
 $lt1x^{m+1} \leq lt1x^{n+1}$ IFF $n \leq m$

both_sides_expt_pos_gt_aux: LEMMA
 $px^{m+1} > py^{m+1}$ IFF $px > py$

both_sides_expt_gt1_gt_aux: LEMMA
 $gt1x^{m+1} > gt1x^{n+1}$ IFF $m > n$

both_sides_expt_lt1_gt_aux: LEMMA
 $lt1x^{m+1} > lt1x^{n+1}$ IFF $n > m$

both_sides_expt_pos_ge_aux: LEMMA
 $px^{m+1} \geq py^{m+1}$ IFF $px \geq py$

both_sides_expt_gt1_ge_aux: LEMMA
 $gt1x^{m+1} \geq gt1x^{n+1}$ IFF $m \geq n$

both_sides_expt_lt1_ge_aux: LEMMA
 $lt1x^{m+1} \geq lt1x^{n+1}$ IFF $n \geq m$

expt_of_mult: LEMMA
 $x \times y^n = x^n \times y^n$

expt_of_div: LEMMA
 $x/n0y^n = x^n/n0y^n$

expt_of_inv: LEMMA $1/n0x^n = 1/n0x^n$

expt_of_abs: LEMMA $|x|^n = |x^n|$

abs_of_expt_inv: LEMMA
 $|1/n0x^n| = 1/|n0x|^n$

expt_x0: LEMMA $x^0 = 1$

expt_x1: LEMMA $x^1 = x$

expt_x2: LEMMA $x^2 = x \times x$

expt_x3: LEMMA $x^3 = x \times x \times x$

expt_x4: LEMMA $x^4 = x \times x \times x \times x$

expt_1i: LEMMA $1^i = 1$

expt_eq_0: LEMMA $x^{pn} = 0$ IFF $x = 0$

expt_plus: LEMMA
 $n0x^{(i+j)} =$
 $n0x^i \times n0x^j$

expt_times: LEMMA
 $n0x^{(i \times j)} = (n0x^i)^j$

expt_inverse: LEMMA
 $n0x^{(-i)} = 1/(n0x^i)$

expt_div: LEMMA
 $n0x^i/n0x^j =$
 $n0x^{(i-j)}$

both_sides_expt1: LEMMA
 $px^{n0i} = px^{n0j}$ IFF $n0i = n0j$ OR $px = 1$

both_sides_expt2: LEMMA $px^{n0i} = py^{n0i}$ IFF $px = py$

b : VAR above(1)

pos_expt_gt: LEMMA $n < b^n$

expt_ge1: LEMMA $b^n \geq 1$

both_sides_expt_pos_lt: LEMMA
 $px^{pm} < py^{pm}$ IFF $px < py$

both_sides_expt_gt1_lt: LEMMA
 $gt1x^i < gt1x^j$ IFF $i < j$

both_sides_expt_lt1_lt: LEMMA
 $lt1x^i < lt1x^j$ IFF $j < i$

both_sides_expt_pos_le: LEMMA

$$px^{pm} \leq py^{pm} \text{ IFF } px \leq py$$

both_sides_expt_gt1_le: LEMMA
 $gt1x^i \leq gt1x^j \text{ IFF } i \leq j$

both_sides_expt_lt1_le: LEMMA
 $lt1x^i \leq lt1x^j \text{ IFF } j \leq i$

both_sides_expt_pos_gt: LEMMA
 $px^{pm} > py^{pm} \text{ IFF } px > py$

both_sides_expt_gt1_gt: LEMMA
 $gt1x^i > gt1x^j \text{ IFF } i > j$

both_sides_expt_lt1_gt: LEMMA
 $lt1x^i > lt1x^j \text{ IFF } j > i$

both_sides_expt_pos_ge: LEMMA
 $px^{pm} \geq py^{pm} \text{ IFF } px \geq py$

both_sides_expt_gt1_ge: LEMMA
 $gt1x^i \geq gt1x^j \text{ IFF } i \geq j$

both_sides_expt_lt1_ge: LEMMA
 $lt1x^i \geq lt1x^j \text{ IFF } j \geq i$

expt_gt1_pos: LEMMA $gt1x^{pm} \geq gt1x$

expt_gt1_neg: LEMMA $gt1x^{(-pm)} < 1$

expt_gt1_nonpos: LEMMA $gt1x^{(-m)} \leq 1$

mult_expt: LEMMA
 $(n0x \times n0y)^i = n0x^i \times n0y^i$

div_expt: LEMMA
 $(n0x/n0y)^i = n0x^i/n0y^i$

inv_expt: LEMMA
 $(1/n0x)^i = 1/n0x^i$

abs_expt: LEMMA $|n0x|^i = |n0x^i|$

abs_hat_nat: LEMMA $|x|^n = |x^n|$

expt_minus1_abs: LEMMA $|(-1)^i| = 1$

even_m1_pow: LEMMA $\text{even?}(i) \text{ IMPLIES } (-1)^i = 1$

not_even_m1_pow: LEMMA
NOT $\text{even?}(i) \text{ IMPLIES } (-1)^i = -1$

expt_lt1_bound1: LEMMA $\text{lt1}x^n \leq 1$

expt_lt1_bound2: LEMMA $\text{lt1}x^{pn} < 1$

expt_gt1_bound1: LEMMA $1 \leq \text{gt1}x^n$

expt_gt1_bound2: LEMMA $1 < \text{gt1}x^{pn}$

large_expt: LEMMA $1 < px \text{ IMPLIES } (\forall py: \exists n: py < px^n)$

small_expt: LEMMA $px < 1 \text{ IMPLIES } (\forall py: \exists n: px^n < py)$

exponent_adjust: LEMMA
$$\frac{b^i + b^{(i-pn)}}{b^{(i+1)}} <$$

exp_of_exists: LEMMA
 $\exists i: b^i \leq py \ \& \ py < b^{(i+1)}$

END exponentiation

euclidean_division: THEORY

BEGIN

a, i : VAR \mathbb{N}

b : VAR $\mathbb{N}_{>0}$

n : VAR \mathbb{Z}

$\text{mod}(b)$: TYPE+ = $\{i \mid i < b\}$

euclid_nat: LEMMA

$\exists (q: \mathbb{N}), (r: \text{mod}(b)): a = b \times q + r$

euclid_int: PROPOSITION

$\exists (q: \mathbb{Z}), (r: \text{mod}(b)): n = b \times q + r$

unique_quotient: PROPOSITION

$\forall (q_1, q_2: \mathbb{Z}), (r_1, r_2: \text{mod}(b)):$
 $b \times q_1 + r_1 = b \times q_2 + r_2$ IMPLIES $q_1 = q_2$

unique_remainder: COROLLARY

$\forall (q_1, q_2: \mathbb{Z}), (r_1, r_2: \text{mod}(b)):$
 $b \times q_1 + r_1 = b \times q_2 + r_2$ IMPLIES $r_1 = r_2$

unique_division: COROLLARY

$\forall (q_1, q_2: \mathbb{Z}), (r_1, r_2: \text{mod}(b)):$
 $b \times q_1 + r_1 = b \times q_2 + r_2$ IMPLIES
 $q_1 = q_2$ AND $r_1 = r_2$

END euclidean_division

```

divides: THEORY
BEGIN

  n, m, l, x, y: VAR ℤ

  p, q: VAR ℕ

  nz: VAR ℤ≠0

  divides(n, m): bool = ∃ x: m = n × x

  divides(n)(m): bool = divides(n, m)

  mult_divides1: JUDGEMENT ×(n, m) HAS_TYPE (divides(n))

  mult_divides2: JUDGEMENT ×(n, m) HAS_TYPE (divides(m))

  divides_sum: LEMMA
    divides(x, n) AND divides(x, m) IMPLIES divides(x, n + m)

  divides_diff: LEMMA
    divides(x, n) AND divides(x, m) IMPLIES divides(x, n - m)

  divides_opposite: LEMMA divides(x, -n) IFF divides(x, n)

  opposite_divides: LEMMA divides(-x, n) IFF divides(x, n)

  divides_prod1: LEMMA divides(x, n) IMPLIES divides(x, n × m)

  divides_prod2: LEMMA divides(x, n) IMPLIES divides(x, m × n)

  divides_prod_elim1: LEMMA
    divides(nz × n, nz × m) IFF divides(n, m)

  divides_prod_elim2: LEMMA
    divides(n × nz, m × nz) IFF divides(n, m)

  divides_reflexive: LEMMA divides(n, n)

  divides_transitive: LEMMA
    divides(n, m) AND divides(m, l) IMPLIES divides(n, l)

```

product_one: LEMMA

$x \times y = 1$ IFF

$(x = 1 \text{ AND } y = 1) \text{ OR } (x = -1 \text{ AND } y = -1)$

mutual_divisors: LEMMA

$\text{divides}(n, m) \text{ AND } \text{divides}(m, n) \text{ IMPLIES } n = m \text{ OR } n = -m$

mutual_divisors_nat: LEMMA

$\text{divides}(p, q) \text{ AND } \text{divides}(q, p) \text{ IMPLIES } p = q$

one_divides: LEMMA $\text{divides}(1, n)$

divides_zero: LEMMA $\text{divides}(n, 0)$

zero_div_zero: LEMMA $\text{divides}(0, n) \text{ IFF } n = 0$

divisors_of_one: LEMMA $\text{divides}(n, 1) \text{ IFF } n = 1 \text{ OR } n = -1$

one_div_one: LEMMA $\text{divides}(p, 1) \text{ IFF } p = 1$

divisor_smaller: LEMMA $\text{divides}(p, q) \text{ IMPLIES } q = 0 \text{ OR } p \leq q$

divides_next: LEMMA

$\text{divides}(n, n + 1) \text{ IFF } n = 1 \text{ OR } n = -1$

p_1 : VAR above(1)

divides_plus_1: LEMMA

$\text{divides}(p_1, nz) \Rightarrow \text{NOT } \text{divides}(p_1, nz + 1)$

END divides

modulo_arithmetic: THEORY

BEGIN

x, y, z, t, q, i : VAR \mathbb{Z}

n0x: VAR $\mathbb{Z}_{\neq 0}$

nx: VAR \mathbb{N}

px, py, b: VAR $\mathbb{N}_{>0}$

n: VAR \mathbb{N}

nrem(x, b): { r : mod(b) | $\exists q$: $x = b \times q + r$ }

rem(b)(x): { r : mod(b) | $\exists q$: $x = b \times q + r$ }

rem_def: LEMMA

$\forall (r$: mod(b)):
rem(b)(x) = r IFF $\exists q$: $x = b \times q + r$

rem_def2: LEMMA

$\forall (r$: mod(b)):
rem(b)(x) = r IFF divides($b, x - r$)

rem_def3: LEMMA

$\forall (r$: mod(b)):
rem(b)(x) = r IFF divides($b, r - x$)

rem_nrem0: LEMMA rem = λb : λx : nrem(x, b)

rem_nrem: LEMMA rem(b)(x) = nrem(x, b)

rem_mod: LEMMA $\forall (r$: mod(b)): rem(b)(r) = r

rem_mod2: LEMMA $0 \leq x$ AND $x < b$ IMPLIES rem(b)(x) = x

rem_zero: LEMMA rem(b)(0) = 0

rem_self: LEMMA rem(b)(b) = 0

rem_multiple1: LEMMA rem(b)($b \times x$) = 0

rem_multiple2: LEMMA $\text{rem}(b)(x \times b) = 0$

rem_one: LEMMA $b \neq 1$ IMPLIES $\text{rem}(b)(1) = 1$

rem_minus_one: LEMMA $\text{rem}(b)(-1) = b - 1$

same_remainder: LEMMA

$\text{rem}(b)(x) = \text{rem}(b)(y)$ IFF $\text{divides}(b, x - y)$

rem_rem: LEMMA $\text{rem}(b)(\text{rem}(b)(x)) = \text{rem}(b)(x)$

rem_sum: LEMMA

$\text{rem}(b)(x) = \text{rem}(b)(y)$ AND $\text{rem}(b)(z) = \text{rem}(b)(t)$ IMPLIES
 $\text{rem}(b)(x + z) = \text{rem}(b)(y + t)$

rem_sum1: LEMMA

$\text{rem}(b)(\text{rem}(b)(x) + y) = \text{rem}(b)(x + y)$

rem_sum2: LEMMA

$\text{rem}(b)(x + \text{rem}(b)(y)) = \text{rem}(b)(x + y)$

rem_diff: LEMMA

$\text{rem}(b)(x) = \text{rem}(b)(y)$ AND $\text{rem}(b)(z) = \text{rem}(b)(t)$ IMPLIES
 $\text{rem}(b)(x - z) = \text{rem}(b)(y - t)$

rem_diff1: LEMMA

$\text{rem}(b)(\text{rem}(b)(x) - y) = \text{rem}(b)(x - y)$

rem_diff2: LEMMA

$\text{rem}(b)(x - \text{rem}(b)(y)) = \text{rem}(b)(x - y)$

rem_prod1: LEMMA

$\text{rem}(b)(\text{rem}(b)(x) \times y) = \text{rem}(b)(x \times y)$

rem_prod2: LEMMA

$\text{rem}(b)(x \times \text{rem}(b)(y)) = \text{rem}(b)(x \times y)$

rem_prod: LEMMA

$\text{rem}(b)(x) = \text{rem}(b)(y)$ AND $\text{rem}(b)(z) = \text{rem}(b)(t)$ IMPLIES
 $\text{rem}(b)(x \times z) = \text{rem}(b)(y \times t)$

rem_expt: LEMMA

$$\begin{aligned} \text{rem}(b)(x) = \text{rem}(b)(y) & \text{ IMPLIES} \\ \text{rem}(b)(x^n) & = \text{rem}(b)(y^n) \end{aligned}$$

rem_expt1: LEMMA

$$\text{rem}(b)(\text{rem}(b)(x^n)) = \text{rem}(b)(x^n)$$

rem_sum_elim1: LEMMA

$$\begin{aligned} \text{rem}(b)(x + y) = \text{rem}(b)(x + z) & \text{ IFF} \\ \text{rem}(b)(y) & = \text{rem}(b)(z) \end{aligned}$$

rem_sum_elim2: LEMMA

$$\begin{aligned} \text{rem}(b)(y + x) = \text{rem}(b)(z + x) & \text{ IFF} \\ \text{rem}(b)(y) & = \text{rem}(b)(z) \end{aligned}$$

rem_diff_elim1: LEMMA

$$\begin{aligned} \text{rem}(b)(x - y) = \text{rem}(b)(x - z) & \text{ IFF} \\ \text{rem}(b)(y) & = \text{rem}(b)(z) \end{aligned}$$

rem_diff_elim2: LEMMA

$$\begin{aligned} \text{rem}(b)(y - x) = \text{rem}(b)(z - x) & \text{ IFF} \\ \text{rem}(b)(y) & = \text{rem}(b)(z) \end{aligned}$$

rem_opposite_elim: LEMMA

$$\begin{aligned} \text{rem}(b)(-x) = \text{rem}(b)(-y) & \text{ IFF} \\ \text{rem}(b)(x) & = \text{rem}(b)(y) \end{aligned}$$

$$\text{ndiv}(x, b): \{q: \mathbb{Z} \mid x = b \times q + \text{rem}(b)(x)\}$$

$$\text{ndiv_lt: LEMMA } \text{ndiv}(x, b) \leq x/b$$

JUDGEMENT $\text{ndiv}(n, b)$ HAS_TYPE upto(n)

rem_floor: LEMMA

$$\forall b, x: x = \text{rem}(b)(x) + b \times \lfloor x/b \rfloor$$

rem_base: LEMMA

$$\begin{aligned} \forall b, x, i, n: \\ \text{rem}(b)(x) = \text{rem}(b+n)(x+i) & \text{ IFF} \\ \text{divides}(b+n, i - n \times \lfloor x/b \rfloor) & \end{aligned}$$

rem_sum_floor: LEMMA

$$\forall b, x, i:$$

$$\text{rem}(b)(x + i) =$$

$$\text{rem}(b)(x) + i - b \times \lfloor (\text{rem}(b)(x) + i) / b \rfloor$$

rem_sum_assoc: COROLLARY

$$\forall b, x, n:$$

$$\text{rem}(b)(x + n) = \text{rem}(b)(x) + n \text{ IFF}$$

$$\text{rem}(b)(x) < b - n$$

rem_add_one: LEMMA

$$\forall b, x:$$

$$\text{rem}(b)(x + 1) = \text{rem}(b)(x) + 1 \text{ OR}$$

$$(\text{rem}(b)(x) = b - 1 \text{ AND}$$

$$\text{rem}(b)(x + 1) = 0)$$

rem_wrap: LEMMA

$$\forall b, x, (n: \text{below}(b)):$$

$$\text{rem}(b)(x) < \text{rem}(b)(x + n) \text{ IFF}$$

$$\text{rem}(b)(x) < b - n \text{ AND } n > 0$$

rem_wrap_eq: COROLLARY

$$\forall b, x, (n: \text{below}(b)):$$

$$\text{rem}(b)(x) \leq \text{rem}(b)(x + n) \text{ IFF}$$

$$\text{rem}(b)(x) < b - n \text{ OR } \text{divides}(b, n)$$

END modulo_arithmetic


```

subrange_inductions[ $i: \mathbb{Z}, j: \text{upfrom}(i)$ ]: THEORY
BEGIN

   $k, m$ : VAR subrange( $i, j$ )

   $p$ : VAR pred[subrange( $i, j$ )]

  subrange_induction: LEMMA
    ( $p(i)$  AND ( $\forall k: k < j$  AND  $p(k)$  IMPLIES  $p(k + 1)$ )) IMPLIES
      ( $\forall k: p(k)$ )

  SUBRANGE_induction: LEMMA
    ( $\forall k: (\forall m: m < k$  IMPLIES  $p(m))$  IMPLIES  $p(k)$ ) IMPLIES
      ( $\forall k: p(k)$ )

END subrange_inductions

```

```

bounded_int_inductions[m: ℤ]: THEORY
BEGIN

  pf: VAR pred[upfrom(m)]

  jf, kf: VAR upfrom(m)

  upfrom_induction: LEMMA
    (pf(m) AND (∀ jf: pf(jf) IMPLIES pf(jf + 1))) IMPLIES
    (∀ jf: pf(jf))

  UPFROM_induction: LEMMA
    (∀ jf: (∀ kf: kf < jf IMPLIES pf(kf)) IMPLIES pf(jf)) IMPLIES
    (∀ jf: pf(jf))

  pa: VAR pred[above(m)]

  ja, ka: VAR above(m)

  above_induction: LEMMA
    (pa(m + 1) AND (∀ ja: pa(ja) IMPLIES pa(ja + 1))) IMPLIES
    (∀ ja: pa(ja))

  ABOVE_induction: LEMMA
    (∀ ja: (∀ ka: ka < ja IMPLIES pa(ka)) IMPLIES pa(ja)) IMPLIES
    (∀ ja: pa(ja))

END bounded_int_inductions

```

```

bounded_nat_inductions[m: ℕ]: THEORY
BEGIN

  pt: VAR pred[upto(m)]

  jt, kt: VAR upto(m)

  upto_induction: LEMMA
    (pt(0) AND (∀ jt: jt < m AND pt(jt) IMPLIES pt(jt + 1))) IMPLIES
    (∀ jt: pt(jt))

  UPTO_induction: LEMMA
    (∀ jt: (∀ kt: kt < jt IMPLIES pt(kt)) IMPLIES pt(jt)) IMPLIES
    (∀ jt: pt(jt))

  pb: VAR pred[below(m)]

  jb, kb: VAR below(m)

  below_induction: LEMMA
    ((m > 0 IMPLIES pb(0)) AND
     (∀ jb: jb < m - 1 AND pb(jb) IMPLIES pb(jb + 1)))
    IMPLIES (∀ jb: pb(jb))

  BELOW_induction: LEMMA
    (∀ jb: (∀ kb: kb < jb IMPLIES pb(kb)) IMPLIES pb(jb)) IMPLIES
    (∀ jb: pb(jb))

END bounded_nat_inductions

```

```
subrange_type[m, n: ℤ]: THEORY
  BEGIN

    subrange: TYPE = subrange(m, n)

  END subrange_type
```

```
int_types[m:  $\mathbb{Z}$ ]: THEORY
  BEGIN

    upfrom: TYPE+ = upfrom(m)

    above: TYPE+ = above(m)

  END int_types
```

```
nat_types [m: ℕ]: THEORY
BEGIN

  upto: TYPE+ = upto(m)

  below: TYPE = below(m)

END nat_types
```

nat_fun_props: THEORY

BEGIN

n, m : VAR \mathbb{N}

injection_n_to_m: THEOREM

$(\exists (f: [\text{below}(n) \rightarrow \text{below}(m)]): \text{injective?}(f))$ IMPLIES
 $n \leq m$

injection_n_to_m_var: THEOREM

$(\exists (f: [\text{below}(n) \rightarrow \text{below}(m)]): \text{injective?}(f))$ IFF
 $n \leq m$

surjection_n_to_m: THEOREM

$(\exists (f: [\text{below}(n) \rightarrow \text{below}(m)]): \text{surjective?}(f))$ IMPLIES
 $m \leq n$

surjection_n_to_m_var: THEOREM

$(\exists (f: [\text{below}(n) \rightarrow \text{below}(m)]): \text{surjective?}(f))$ IFF
 $(m > 0 \text{ AND } m \leq n) \text{ OR } (m = 0 \text{ AND } n = 0)$

bijection_n_to_m: THEOREM

$(\exists (f: [\text{below}(n) \rightarrow \text{below}(m)]): \text{bijective?}(f))$ IFF
 $n = m$

injection_n_to_m2: THEOREM

$(\exists (f: [\text{upto}(n) \rightarrow \text{upto}(m)]): \text{injective?}(f))$ IFF
 $n \leq m$

surjection_n_to_m2: THEOREM

$(\exists (f: [\text{upto}(n) \rightarrow \text{upto}(m)]): \text{surjective?}(f))$ IFF
 $m \leq n$

bijection_n_to_m2: THEOREM

$(\exists (f: [\text{upto}(n) \rightarrow \text{upto}(m)]): \text{bijective?}(f))$ IFF
 $n = m$

surj_equiv_inj: THEOREM

$\forall (f: [\text{below}(n) \rightarrow \text{below}(n)]):$
 $\text{surjective?}(f)$ IFF $\text{injective?}(f)$

inj_equiv_bij: THEOREM

$\forall (f: [\text{below}(n) \rightarrow \text{below}(n)]):$
bijjective?(f) IFF injective?(f)

surj_equiv_bij: THEOREM

$\forall (f: [\text{below}(n) \rightarrow \text{below}(n)]):$
bijjective?(f) IFF surjective?(f)

surj_equiv_inj2: THEOREM

$\forall (f: [\text{upto}(n) \rightarrow \text{upto}(n)]):$
surjective?(f) IFF injective?(f)

inj_equiv_bij2: THEOREM

$\forall (f: [\text{upto}(n) \rightarrow \text{upto}(n)]):$
bijjective?(f) IFF injective?(f)

surj_equiv_bij2: THEOREM

$\forall (f: [\text{upto}(n) \rightarrow \text{upto}(n)]):$
bijjective?(f) IFF surjective?(f)

END nat_fun_props


```

finite_sets[T : TYPE] : THEORY
BEGIN

  x, y, z : VAR T

  s : VAR set[T]

  N : VAR ℕ

  is_finite(s) : bool =
    (∃ N, (f : [(s) → below[N]]): injective?(f))

  finite_set : TYPE = (is_finite) CONTAINING ∅[T]

  non_empty_finite_set : TYPE = {s : finite_set | NOT empty?(s)}

  is_finite_surj : LEMMA
    (∃ (N : ℕ), (f : [below[N] → (s)]): surjective?(f)) IFF
    is_finite(s)

  A, B : VAR finite_set

  NA, NB : VAR non_empty_finite_set

  finite_subset : LEMMA (s ⊆ A) IMPLIES is_finite(s)

  finite_intersection : LEMMA is_finite((A ∩ B))

  finite_add : LEMMA is_finite((A ∪ {x}))

  nonempty_finite_is_nonempty : JUDGEMENT non_empty_finite_set SUBTYPE_OF
    (nonempty?[T])

  finite_singleton : JUDGEMENT singleton(x) HAS_TYPE finite_set

  finite_union : JUDGEMENT union(A, B) HAS_TYPE finite_set

  finite_intersection1 : JUDGEMENT intersection(s, A) HAS_TYPE
    finite_set

  finite_intersection2 : JUDGEMENT intersection(A, s) HAS_TYPE
    finite_set

```

finite_difference: JUDGEMENT $\text{difference}(A, s)$ HAS_TYPE finite_set

nonempty_finite_union1: JUDGEMENT $\text{union}(NA, B)$ HAS_TYPE
non_empty_finite_set

nonempty_finite_union2: JUDGEMENT $\text{union}(A, NB)$ HAS_TYPE
non_empty_finite_set

nonempty_add_finite: JUDGEMENT $\text{add}(x, A)$ HAS_TYPE
non_empty_finite_set

finite_remove: JUDGEMENT $\text{remove}(x, A)$ HAS_TYPE finite_set

finite_rest: JUDGEMENT $\text{rest}(A)$ HAS_TYPE finite_set

finite_emptyset: JUDGEMENT \emptyset HAS_TYPE finite_set

nonempty_singleton_finite: JUDGEMENT $\text{singleton}(x)$ HAS_TYPE
non_empty_finite_set

is_finite_type: bool =
 $(\exists N, (g: [T \rightarrow \text{below}[N]]): \text{injective?}(g))$

finite_full: LEMMA is_finite_type IFF is_finite(fullset[T])

finite_type_set: LEMMA is_finite_type IMPLIES is_finite(s)

finite_complement: LEMMA is_finite_type IMPLIES is_finite(\bar{s})

S, S_2 : VAR finite_set

n, m : VAR \mathbb{N}

p : VAR $\mathbb{N}_{>0}$

inj_set(S): (nonempty?[\mathbb{N}]) =
 $\{n \mid \exists (f: [S \rightarrow \text{below}[n]]): \text{injective?}(f)\}$

Card(S): $\mathbb{N} = \min(\text{inj_set}(S))$

inj_Card: LEMMA

$\text{Card}(S) = n$ IMPLIES
 $(\exists (f: [(S) \rightarrow \text{below}[n]]): \text{injective?}(f))$

reduce_inj: LEMMA
 $(\forall (f: [(S) \rightarrow \text{below}[p]]):$
 $\text{injective?}(f)$ AND NOT $\text{surjective?}(f)$ IMPLIES
 $(\exists (g: [(S) \rightarrow \text{below}[p-1]]):$
 $\text{injective?}(g))$

Card_injection: LEMMA
 $(\exists (f: [(S) \rightarrow \text{below}[n]]): \text{injective?}(f))$ IMPLIES
 $\text{Card}(S) \leq n$

Card_surjection: LEMMA
 $(\exists (f: [(S) \rightarrow \text{below}[n]]): \text{surjective?}(f))$ IMPLIES
 $n \leq \text{Card}(S)$

Card_bijection: THEOREM
 $\text{Card}(S) = n$ IFF
 $(\exists (f: [(S) \rightarrow \text{below}[n]]): \text{bijective?}(f))$

Card_disj_union: THEOREM
 $\text{disjoint?}(S, S_2)$ IMPLIES
 $\text{Card}((S \cup S_2)) = \text{Card}(S) + \text{Card}(S_2)$

$\text{card}(S): \{n: \mathbb{N} \mid n = \text{Card}(S)\}$

card_def: THEOREM $\text{card}(S) = \text{Card}(S)$

card_emptyset: THEOREM $\text{card}(\emptyset[T]) = 0$

empty_card: THEOREM $\text{empty?}(S)$ IFF $\text{card}(S) = 0$

card_empty?: THEOREM $(\text{card}(S) = 0) = \text{empty?}(S)$

card_is_0: THEOREM $(\text{card}(S) = 0) = (S = \emptyset)$

nonempty_card: THEOREM $\text{nonempty?}(S)$ IFF $\text{card}(S) > 0$

card_singleton: THEOREM $\text{card}(\text{singleton}(x)) = 1$

card_one: THEOREM $\text{card}(S) = 1$ IFF $(\exists x: S = \text{singleton}(x))$

card_disj_union: THEOREM
 $\text{disjoint?}(A, B) \text{ IMPLIES}$
 $\text{card}((A \cup B)) = \text{card}(A) + \text{card}(B)$

card_diff_subset: THEOREM
 $(A \subseteq B) \text{ IMPLIES}$
 $\text{card}((B \setminus A)) = \text{card}(B) - \text{card}(A)$

card_subset: THEOREM $(A \subseteq B) \text{ IMPLIES } \text{card}(A) \leq \text{card}(B)$

card_plus: THEOREM
 $\text{card}(A) + \text{card}(B) =$
 $\text{card}((A \cup B)) + \text{card}((A \cap B))$

card_union: THEOREM
 $\text{card}((A \cup B)) =$
 $\text{card}(A) + \text{card}(B) - \text{card}((A \cap B))$

card_add: THEOREM
 $\text{card}((S \cup \{x\})) =$
 $\text{card}(S) + \text{IF } S(x) \text{ THEN } 0 \text{ ELSE } 1 \text{ ENDIF}$

card_add_gt0: THEOREM $\text{card}((S \cup \{x\})) > 0$

card_remove: THEOREM
 $\text{card}((S \setminus \{x\})) =$
 $\text{card}(S) - \text{IF } S(x) \text{ THEN } 1 \text{ ELSE } 0 \text{ ENDIF}$

card_rest: THEOREM
 $\text{NOT empty?}(S) \text{ IMPLIES } \text{card}(\text{rest}(S)) = \text{card}(S) - 1$

same_card_subset: THEOREM
 $(A \subseteq B) \text{ AND } \text{card}(A) = \text{card}(B) \text{ IMPLIES } A = B$

smaller_card_subset: THEOREM
 $(A \subseteq B) \text{ AND } \text{card}(A) < \text{card}(B) \text{ IMPLIES}$
 $(\exists x: (x \in B) \text{ AND NOT } (x \in A))$

card_strict_subset: THEOREM $(A \subset B) \text{ IMPLIES } \text{card}(A) < \text{card}(B)$

card_1_has_1: THEOREM $\text{card}(S) \geq 1 \text{ IMPLIES } (\exists (x: T): S(x))$

card_2_has_2: THEOREM

card(S) \geq 2 IMPLIES

$(\exists (x, y: T): x \neq y \text{ AND } S(x) \text{ AND } S(y))$

card_intersection_le: THEOREM

card($(A \cap B)$) \leq card(A) AND

card($(A \cap B)$) \leq card(B)

card_bij: THEOREM

card(S) = N IFF

$(\exists (f: [(S) \rightarrow \text{below}[N]]): \text{bijective?}(f))$

card_bij_inv: THEOREM

card(S) = N IFF

$(\exists (f: [\text{below}[N] \rightarrow (S)]): \text{bijective?}(f))$

bij_exists: COROLLARY

$(\exists (f: [(S) \rightarrow \text{below}(\text{card}(S))]): \text{bijective?}(f))$

bij(S : finite_set):

$\{f: [(S) \rightarrow \text{below}(\text{card}(S))] \mid \text{bijective?}(f)\}$

ibij(S : non_empty_finite_set): $\{f: [\text{below}(\text{card}(S)) \rightarrow (S)] \mid \text{bijective?}(f)\} =$
inverse(bij(S))

bij_ibij: LEMMA

$\forall (S: \text{non_empty_finite_set}, ii: \text{below}(\text{card}(S))):$

bij(S)(ibij(S)(ii)) = ii

ibij_bij: LEMMA

$\forall (S: \text{non_empty_finite_set}, x: T):$

$S(x)$ IMPLIES ibij(S)(bij(S)(x)) = x

is_finite_exists_N: LEMMA

$\forall (g: [\text{below}[N] \rightarrow T]):$

is_finite($\{r: T \mid \exists (n: \text{below}[N]): r = g(n)\}$)

$P, P_1, P_2: \text{VAR pred}[T]$

finite_pred: LEMMA

is_finite(fullset $[T]$) IMPLIES

is_finite[T]({x: T | P(x)})

finite_pred2: LEMMA

is_finite(P) IMPLIES is_finite[T]({x: T | P(x)})

card_implies: LEMMA

is_finite(fullset[T]) AND ($\forall (x: T): P_1(x) \text{ IMPLIES } P_2(x)$) IMPLIES
card({x: T | P₁(x)}) ≤ card({x: T | P₂(x)})

finite_induction: THEOREM

$\forall (p: \text{pred}[\text{set}[T]]):$
 $(\forall (n: \mathbb{N}), (S: \text{set}[T])):$
 $(\exists (f: [(S) \rightarrow \text{below}[n]]): \text{injective?}(f)) \Rightarrow p(S)$
 $\Rightarrow (\forall (\text{FS}: \text{finite_set}): p(\text{FS}))$

END finite_sets

```

restrict_set_props [T : TYPE, S : TYPE FROM T] : THEORY
BEGIN

  restrict_finite : LEMMA
     $\forall (a : \text{set}[T]) :$ 
       $\text{is\_finite}(a) \Rightarrow \text{is\_finite}(\text{restrict}[T, S, \text{bool}](a))$ 

  finite_restrict : JUDGEMENT  $\text{restrict}[T, S, \text{bool}](a : \text{finite\_set}[T]) \text{ HAS\_TYPE}$ 
     $\text{finite\_set}[S]$ 

  empty_restrict : JUDGEMENT  $\text{restrict}[T, S, \text{bool}](a : (\text{empty?}[T])) \text{ HAS\_TYPE}$ 
     $(\text{empty?}[S])$ 

  card_restrict : LEMMA
     $\forall (a : \text{finite\_set}[T]) :$ 
       $\text{card}(\text{restrict}[T, S, \text{bool}](a)) \leq \text{card}(a)$ 

END restrict_set_props

```

```

extend_set_props [T : TYPE, S : TYPE FROM T] : THEORY
BEGIN

finite_extension : LEMMA
  ∀ (a : set [S]) :
    is_finite(extend [T, S, bool, FALSE] (a)) IFF is_finite(a)

finite_extend : JUDGEMENT extend [T, S, bool, FALSE] (a : finite_set [S]) HAS_TYPE
  finite_set [T]

empty_extend : JUDGEMENT extend [T, S, bool, FALSE] (a : (empty? [S])) HAS_TYPE
  (empty? [T])

nonempty_extend : JUDGEMENT extend [T, S, bool, FALSE] (a : (nonempty? [S])) HAS_TYPE
  (nonempty? [T])

singleton_extend : JUDGEMENT extend [T, S, bool, FALSE] (a : (singleton? [S])) HAS_TYPE
  (singleton? [T])

card_extend : LEMMA
  ∀ (a : finite_set [S]) :
    card(extend [T, S, bool, FALSE] (a)) = card(a)

empty?_extend : LEMMA
  ∀ (a : set [S]) :
    empty?(extend [T, S, bool, FALSE] (a)) IFF empty?(a)

nonempty?_extend : LEMMA
  ∀ (a : set [S]) :
    nonempty?(extend [T, S, bool, FALSE] (a)) IFF nonempty?(a)

singleton?_extend : LEMMA
  ∀ (a : set [S]) :
    singleton?(extend [T, S, bool, FALSE] (a)) IFF singleton?(a)

subset_extend : LEMMA
  ∀ (a, b : set [S]) :
    (extend [T, S, bool, FALSE] (a) ⊆ extend [T, S, bool, FALSE] (b)) IFF
    (a ⊆ b)

union_extend : LEMMA
  ∀ (a, b : set [S]) :

```


$$(\text{extend}[T, S, \text{bool}, \text{FALSE}](a) \cup \text{extend}[T, S, \text{bool}, \text{FALSE}](b)) = \text{extend}[T, S, \text{bool}, \text{FALSE}]((a \cup b))$$

intersection_extend: LEMMA

$$\begin{aligned} &\forall (a, b: \text{set}[S]): \\ &(\text{extend}[T, S, \text{bool}, \text{FALSE}](a) \cap \text{extend}[T, S, \text{bool}, \text{FALSE}](b)) = \\ &\text{extend}[T, S, \text{bool}, \text{FALSE}]((a \cap b)) \end{aligned}$$

difference_extend: LEMMA

$$\begin{aligned} &\forall (a, b: \text{set}[S]): \\ &(\text{extend}[T, S, \text{bool}, \text{FALSE}](a) \setminus \text{extend}[T, S, \text{bool}, \text{FALSE}](b)) = \\ &\text{extend}[T, S, \text{bool}, \text{FALSE}]((a \setminus b)) \end{aligned}$$

add_extend: LEMMA

$$\begin{aligned} &\forall (x: S, a: \text{set}[S]): \\ &(\text{extend}[T, S, \text{bool}, \text{FALSE}](a) \cup \{x\}) = \\ &\text{extend}[T, S, \text{bool}, \text{FALSE}]((a \cup \{x\})) \end{aligned}$$

remove_extend: LEMMA

$$\begin{aligned} &\forall (x: S, a: \text{set}[S]): \\ &(\text{extend}[T, S, \text{bool}, \text{FALSE}](a) \setminus \{x\}) = \\ &\text{extend}[T, S, \text{bool}, \text{FALSE}]((a \setminus \{x\})) \end{aligned}$$

END extend_set_props

```

function_image_aux[D: TYPE, R: TYPE]: THEORY
BEGIN

  S: VAR finite_set[D]

  f: VAR [D → R]

  inj: VAR (injective?[D, R])

  finite_image: JUDGEMENT image(f, S) HAS_TYPE finite_set[R]

  card_image: LEMMA ∀ f, S: card(image(f, S)) ≤ card(S)

  card_injective_image: LEMMA
    ∀ inj, S: card(image(inj, S)) = card(S)

  bijective_image: LEMMA
    ∀ inj: bijective?[D, (image(inj, fullset[D]))](inj)

END function_image_aux

```

```

function_iterate [T: TYPE]: THEORY
BEGIN

  f: VAR [T → T]

  m, n: VAR ℕ

  x: VAR T

  iterate(f, n)(x): RECURSIVE T =
    IF n = 0 THEN x ELSE f(iterate(f, n - 1)(x)) ENDIF
    MEASURE n

  iterate_add: LEMMA
    iterate(f, m) ∘ iterate(f, n) =
      iterate(f, m + n)

  iterate_add_applied: LEMMA
    iterate(f, m)(iterate(f, n)(x)) =
      iterate(f, m + n)(x)

  iterate_add_one: LEMMA
    iterate(f, n)(f(x)) = iterate(f, n + 1)(x)

  iterate_mult: LEMMA
    iterate(iterate(f, m), n) = iterate(f, m × n)

  iterate_invariant: LEMMA
    f(iterate(f, n)(x)) = iterate(f, n)(f(x))

END function_iterate

```

```

sequences[T: TYPE]: THEORY
BEGIN

  sequence: TYPE = [N → T]

  i, n: VAR N

  x: VAR T

  p: VAR pred[T]

  seq: VAR sequence

  Trel: VAR PRED[[T]]

  nth(seq, n): T = seq(n)

  suffix(seq, n): sequence = (λ i: seq(i + n))

  first(seq): T = nth(seq, 0)

  rest(seq): sequence = suffix(seq, 1)

  delete(n, seq): sequence =
    (λ i:
      (IF i < n THEN seq(i) ELSE seq(i + 1) ENDIF))

  insert(x, n, seq): sequence =
    (λ i:
      (IF i < n
        THEN seq(i)
        ELSIF i = n THEN x
        ELSE seq(i - 1)
        ENDIF))

  add(x, seq): sequence = insert(x, 0, seq)

  insert_delete: LEMMA
    insert(nth(seq, n), n, delete(n, seq)) = seq

  add_first_rest: LEMMA add(first(seq), rest(seq)) = seq

```

every(p)(seq): bool = ($\forall n: p(\text{nth}(\text{seq}, n))$)

every(p , seq): bool = ($\forall n: p(\text{nth}(\text{seq}, n))$)

some(p)(seq): bool = ($\exists n: p(\text{nth}(\text{seq}, n))$)

some(p , seq): bool = ($\exists n: p(\text{nth}(\text{seq}, n))$)

sequence_induction: LEMMA

$p(\text{nth}(\text{seq}, 0))$ AND ($\forall n: p(\text{nth}(\text{seq}, n))$ IMPLIES $p(\text{nth}(\text{seq}, n + 1))$) IMPLIES
every(p)(seq)

ascends?(seq, Trel): bool =
preserves(seq, ($\lambda i, n: i \leq n$), Trel)

descends?(seq, Trel): bool =
inverts(seq, ($\lambda i, n: i \leq n$), Trel)

END sequences

```

seq_functions[D, R: TYPE]: THEORY
BEGIN

  f: VAR [D → R]

  s: VAR sequence[D]

  n: VAR ℕ

  map(f)(s): sequence[R] = (λ n: f(nth(s, n)))

  map(f, s): sequence[R] = (λ n: f(nth(s, n)))

END seq_functions

```

```

finite_sequences[T: TYPE]: THEORY
BEGIN

  finite_sequence: TYPE = [#length: ℕ, seq: [below[length] → T]#]

  finseq: TYPE = finite_sequence

  fs, fs1, fs2, fs3: VAR finseq

  m, n: VAR ℕ

  empty_seq: finseq =
    (#length := 0,
     seq := (λ (x: below[0]): ε! (t: T): TRUE)#)

  finseq_appl(fs): [below[length(fs)] → T] = fs`seq;

  CONVERSION finseq_appl

  fs1 ∘ fs2: finseq =
    LET l1 = fs1`length, lsum = l1 + fs2`length IN
    (#length := lsum,
     seq
     := (λ (n: below[lsum]):
         IF n < l1
         THEN fs1`seq(n)
         ELSE fs2`seq(n - l1)
         ENDIF)#);

  p: VAR [ℕ]

  fsp: finseq =
    LET (m, n) = p IN
    IF m > n OR m ≥ fs`length
    THEN empty_seq
    ELSE LET len = min(n - m + 1, fs`length - m) IN
         (#length := len, seq := (λ (x: below[len]): fs`seq(x + m))#)
    ENDIF;

  ^^ (fs, p): finseq =
    LET (m, n) = p IN

```

```

IF  $m \geq n$  OR  $m \geq \text{fs`length}$ 
  THEN empty_seq
ELSE LET len =  $\min(n - m, \text{fs`length} - m)$  IN
  (#length := len, seq := ( $\lambda (x: \text{below}[\text{len}]$ ): fs`seq( $x + m$ ))#)
ENDIF

```

```
extract1(fs: {fs | fs`length = 1}): T = fs`seq(0)
```

```
CONVERSION extract1
```

```
O_ASSOC: LEMMA
```

```
fs1  $\circ$  (fs2  $\circ$  fs3) = (fs1  $\circ$  fs2)  $\circ$  fs3
```

```
END finite_sequences
```



```

more_finseq[T: TYPE]: THEORY
BEGIN

  seq: TYPE = finseq[T]

  rr, ss, tt: VAR seq

  x, y, z: VAR T

  prefix?(rr, ss): bool =
    rr`length ≤ ss`length AND
    (∀ (i: below(rr`length)): rr`seq(i) = ss`seq(i))

  prefix_closed?(X: set[seq]): bool =
    ∀ ss: X(ss) IMPLIES (∀ (rr: seq | prefix?(rr, ss)): X(rr))

  add(x, rr): finseq[T] =
    rr WITH [ `length := rr`length + 1, `seq(rr`length) := x ]

END more_finseq

```

```
ordstruct: DATATYPE
BEGIN
  zero: zero?
  add(coef:  $\mathbb{N}_{>0}$ , exp: ordstruct,
      rest: ordstruct)
      : nonzero?
END ordstruct
```

ordinals: THEORY

BEGIN

i, j, k : VAR $\mathbb{N}_{>0}$

m, n, O : VAR \mathbb{N}

u, v, w, x, y, z : VAR ordstruct

size: [ordstruct \rightarrow \mathbb{N}] =
 reduce_nat(0, ($\lambda i, m, n$: 1 + $m + n$));

$<(x, y)$: RECURSIVE bool =

 CASES x OF

 zero: NOT zero?(y),

 add(i, u, v):

 CASES y OF

 zero: FALSE,

 add(j, z, w):

 ($u < z$) OR

 ($u = z$) AND ($i < j$) OR ($u = z$) AND ($i = j$) AND ($v < w$)

 ENDCASES

 ENDCASES

 MEASURE size(x);

$>(x, y)$: bool = $y < x$;

$\leq(x, y)$: bool = $x < y$ OR $x = y$;

$\geq(x, y)$: bool = $y < x$ OR $y = x$

ordinal?(x): RECURSIVE bool =

 CASES x OF

 zero: TRUE,

 add(i, u, v):

 (ordinal?(u) AND

 ordinal?(v) AND CASES v OF zero: TRUE, add(k, r, s): $r < u$ ENDCASES)

 ENDCASES

 MEASURE size

ordinal: TYPE+ = (ordinal?)

r, s, t : VAR ordinal

ordinal_irreflexive: LEMMA NOT $r < r$

ordinal_antisym: LEMMA $r < s$ IMPLIES NOT $s < r$

ordinal_antisymmetric: LEMMA $r \leq s$ AND $s \leq r$ IMPLIES $r = s$

ordinal_transitive: LEMMA $r < s$ AND $s < t$ IMPLIES $r < t$

ordinal_trichotomy: LEMMA $r < s$ OR $r = s$ OR $s < r$

p : VAR pred[ordinal]

ordinal_induction: AXIOM

$(\forall r: (\forall s: s < r \text{ IMPLIES } p(s)) \text{ IMPLIES } p(r)) \text{ IMPLIES } (\forall r: p(r))$

well_founded_le: LEMMA

$\text{well_founded?}(\lambda (r, s: (\text{ordinal?})): r < s)$

END ordinals

lex2: THEORY

BEGIN

i, j, m, n : VAR \mathbb{N}

lex2(m, n): ordinal =

(IF $m = 0$

THEN IF $n = 0$ THEN zero ELSE add($n, zero, zero$) ENDIF

ELSIF $n = 0$ THEN add($m, add(1, zero, zero), zero$)

ELSE add($m, add(1, zero, zero), add(n, zero, zero)$)

ENDIF)

lex2_lt: LEMMA

(lex2(i, j) < lex2(m, n)) =

($i < m$ OR ($i = m$ AND $j < n$))

END lex2

```
list[T: TYPE]: DATATYPE
  BEGIN
    null: null?
    cons(car: T, cdr: list): cons?
  END list
```

```

list_props[T: TYPE]: THEORY
BEGIN

  l, l1, l2, l3: VAR list[T]

  x: VAR T

  P, Q: VAR PRED[T]

  length(l): RECURSIVE ℕ =
    CASES l OF null: 0, cons(x, y): length(y) + 1 ENDCASES
    MEASURE reduce_nat(0, (λ (x: T), (n: ℕ): n + 1))

  member(x, l): RECURSIVE bool =
    CASES l OF null: FALSE, cons(hd, tl): x = hd OR member(x, tl) ENDCASES
    MEASURE length(l)

  member_null: LEMMA member(x, l) IMPLIES NOT null?(l)

  nth(l, (n: below[length(l)])): RECURSIVE T =
    IF n = 0 THEN car(l) ELSE nth(cdr(l), n - 1) ENDIF
    MEASURE length(l)

  append(l1, l2): RECURSIVE list[T] =
    CASES l1 OF null: l2, cons(x, y): cons(x, append(y, l2)) ENDCASES
    MEASURE length(l1)

  reverse(l): RECURSIVE list[T] =
    CASES l OF null: l, cons(x, y): append(reverse(y), cons(x, null)) ENDCASES
    MEASURE length

  append_null: LEMMA append(l, null) = l

  append_assoc: LEMMA
    append(append(l1, l2), l3) = append(l1, append(l2, l3))

  reverse_append: LEMMA
    reverse(append(l1, l2)) = append(reverse(l2), reverse(l1))

  reverse_reverse: LEMMA reverse(reverse(l)) = l

  length_append: LEMMA

```

$\text{length}(\text{append}(l_1, l_2)) = \text{length}(l_1) + \text{length}(l_2)$

`length_reverse: LEMMA` $\text{length}(\text{reverse}(l)) = \text{length}(l)$

$a, b, c: \text{VAR } T$

`list_rep: LEMMA`
 $(:a, b, c:) = \text{cons}(a, \text{cons}(b, \text{cons}(c, \text{null})))$

`every_append: LEMMA`
 $\text{every}(P)(\text{append}(l_1, l_2)) \text{ IFF}$
 $(\text{every}(P)(l_1) \text{ AND } \text{every}(P)(l_2))$

`every_disjunct1: LEMMA`
 $\text{every}(P)(l) \text{ IMPLIES}$
 $\text{every}(\lambda (x: T): \mathbb{P}(x) \text{ OR } Q(x))(l)$

`every_disjunct2: LEMMA`
 $\text{every}(Q)(l) \text{ IMPLIES}$
 $\text{every}(\lambda (x: T): \mathbb{P}(x) \text{ OR } Q(x))(l)$

`every_conjunct: LEMMA`
 $\text{every}(\lambda (x: T): \mathbb{P}(x) \text{ AND } Q(x))(l) \Rightarrow$
 $(\text{every}(P)(l) \text{ AND } \text{every}(Q)(l))$

`every_conjunct2: LEMMA`
 $(\text{every}(P)(l) \text{ AND } \text{every}(Q)(l)) \Rightarrow$
 $\text{every}(\lambda (x: T): \mathbb{P}(x) \text{ AND } Q(x))(l)$

`every_member: LEMMA` $\text{every}(\{c: T \mid \text{member}(c, l)\})(l)$

`every_nth: LEMMA`
 $\text{every}(P)(l) \text{ IFF}$
 $\forall (i: \text{below}(\text{length}(l))): \mathbb{P}(\text{nth}(l, i))$

`END list_props`


```

map_props [T1, T2, T3: TYPE]: THEORY
BEGIN

  f1: VAR [T1 → T2]

  f2: VAR [T2 → T3]

  s: VAR sequence [T1]

  l: VAR list [T1]

  map_list_composition: LEMMA
    map(f2)(map(f1)(l)) = map(f2 ∘ f1)(l)

  map_seq_composition: LEMMA
    map(f2)(map(f1)(s)) = map(f2 ∘ f1)(s)

END map_props

```

```

more_map_props [T1, T2: TYPE]: THEORY
BEGIN

  f: VAR [T1 → T2]

  l: VAR list[T1]

  map_length: LEMMA length(map(f)(l)) = length(l)

  map_nth_rw: LEMMA
    ∀ (i: below(length(l))):
      nth(map(f)(l), i) = f(nth(l, i))

END more_map_props

```

```

filters[T: TYPE]: THEORY
BEGIN

  s: VAR set[T]

  l: VAR list[T]

  p: VAR pred[T]

  filter(s, p): set[T] = {x: T | s(x) & p(x)}

  filter(p)(s): set[T] = {x: T | s(x) & p(x)}

  filter(l, p): RECURSIVE list[T] =
    CASES l OF
      null: null,
      cons(x, y): IF p(x) THEN cons(x, filter(y, p)) ELSE filter(y, p) ENDIF
    ENDCASES
    MEASURE length(l)

  filter(p)(l): RECURSIVE list[T] =
    CASES l OF
      null: null,
      cons(x, y): IF p(x) THEN cons(x, filter(p)(y)) ELSE filter(p)(y) ENDIF
    ENDCASES
    MEASURE length(l)

END filters

```

```

list2finseq[T : TYPE] : THEORY
BEGIN

  l : VAR list[T]

  fs : VAR finseq[T]

  n : VAR ℕ

  list2finseq(l) : finseq[T] =
    (#length := length(l),
     seq := (λ (x : below[length(l)]): nth(l, x))#)

  finseq2list_rec(fs, (n : ℕ | n ≤ length(fs))) : RECURSIVE list[T] =
    IF n = 0
      THEN null
    ELSE cons(fs`seq(length(fs) - n), finseq2list_rec(fs, n - 1))
    ENDIF
    MEASURE n

  finseq2list(fs) : list[T] = finseq2list_rec(fs, length(fs))

  CONVERSION list2finseq

  CONVERSION finseq2list

END list2finseq

```

```

list2set[T: TYPE]: THEORY
BEGIN

  l: VAR list[T]

  x: VAR T

  list2set(l): RECURSIVE set[T] =
    CASES l OF null:  $\emptyset[T]$ , cons(x, y): (list2set(y)  $\cup$  {x}) ENDCASES
    MEASURE length

  CONVERSION list2set

END list2set

```

```
disjointness: THEORY
BEGIN

  l: VAR list[bool]

  pairwise_disjoint?(l): RECURSIVE boolean =
    CASES l OF
      null: TRUE,
      cons(x, y): every( $\lambda (z: \text{bool}): \text{NOT } (x \text{ AND } z)$ )(y) AND pairwise_disjoint?(y)
    ENDCASES
    MEASURE length(l)

END disjointness
```

```
character : DATATYPE
BEGIN
  char(code : below[256]): char?
END character
```

```

strings: THEORY
BEGIN

char: TYPE = (char?)

string: TYPE = finite_sequence[char]

l1, l2: VAR list[char]

c1, c2: VAR char

fseq_lem: LEMMA
(list2finseq(l1) = list2finseq(l2)) = (l1 = l2)

cons_lem: LEMMA
(cons(c1, l1) = cons(c2, l2)) = (c1 = c2 & l1 = l2)

char_lem: LEMMA (c1 = c2) = (code(c1) = code(c2))

END strings

```



```
lift[T: TYPE]: DATATYPE
BEGIN
  bottom: bottom?
  up(down: T): up?
END lift
```

```
union[T1, T2: TYPE]: DATATYPE
  BEGIN
    inl(left: T1): inl?
    inr(right: T2): inr?
  END union
```

```

mucalculus[T: TYPE]: THEORY
BEGIN

  s: VAR T

  p, p1, p2: VAR pred[T]

  predicate_transformer: TYPE = [pred[T] → pred[T]]

  pt: VAR predicate_transformer

  setofpred: VAR pred[pred[T]]

  ≤(p1, p2): bool = ∀ s: p1(s) IMPLIES p2(s)

  monotonic?(pt): bool = ∀ p1, p2: p1 ≤ p2 IMPLIES pt(p1) ≤ pt(p2)

  pp: VAR (monotonic?)

  fixpoint?(pp, p): bool = (pp(p) = p)

  fixpoint?(pp)(p): bool = fixpoint?(pp, p)

  glb(setofpred): pred[T] =
    λ s: (∀ p: (p ∈ setofpred) IMPLIES p(s))

  lfp(pp): pred[T] = glb({p | pp(p) ≤ p})

  lfp_induction: FORMULA pp(p) ≤ p IMPLIES lfp(pp) ≤ p

  μ(pp): pred[T] = lfp(pp)

  lfp?(pp, p1): bool =
    fixpoint?(pp, p1) AND ∀ p2: fixpoint?(pp, p2) IMPLIES p1 ≤ p2

  lfp?(pp)(p1): bool = lfp?(pp, p1)

  lub(setofpred): pred[T] =
    λ s: ∃ p: (p ∈ setofpred) AND p(s)

  gfp(pp): pred[T] = lub({p | p ≤ (pp(p))})

```

gfp_induction: FORMULA $p \leq pp(p)$ IMPLIES $p \leq \text{gfp}(pp)$

$\nu(pp)$: $\text{pred}[T] = \text{gfp}(pp)$

$\text{gfp?}(pp, p_1)$: bool =
fixpoint?(pp, p_1) AND $\forall p_2$: fixpoint?(pp, p_2) IMPLIES $p_2 \leq p_1$

$\text{gfp?}(pp)(p_1)$: bool = $\text{gfp?}(pp, p_1)$

END mu calculus

```

ctlops[state: TYPE]: THEORY
BEGIN

```

```

  u, v, w: VAR state

```

```

  f, g, Q, P, p1, p2: VAR pred[state]

```

```

  Z: VAR pred[[state]]

```

```

  N: VAR [state, state → bool]

```

```

  CONVERSION+ K_conversion

```

```

  EX(N, f)(u): bool = (∃ v: (f(v) AND N(u, v)))

```

```

  EG(N, f): pred[state] =
    ν(λ Q: (λ (s: state): f(s) AND EX(N, Q)(s)))

```

```

  EU(N, f, g): pred[state] =
    μ(λ Q:
      (λ (s1: state):
        g(s1) OR
        (λ (s: state): f(s) AND EX(N, Q)(s)
          (s1))))

```

```

  EF(N, f): pred[state] =
    EU(N, K_conversion[boolean, state](TRUE), f)

```

```

  AX(N, f): pred[state] =
    λ (s1: state): NOT EX(N, λ (s: state): NOT f(s))(s1)

```

```

  AF(N, f): pred[state] =
    λ (s1: state): NOT EG(N, λ (s: state): NOT f(s))(s1)

```

```

  AG(N, f): pred[state] =
    λ (s1: state): NOT EF(N, λ (s: state): NOT f(s))(s1)

```

```

  AU(N, f, g): pred[state] =
    λ (s3: state):
      (λ (s2: state):
        NOT EU(N, λ (s: state): NOT g(s),

```

```
(λ (s1: state):  
  (λ (s: state): NOT f(s))(s1) AND  
  (λ (s: state): NOT g(s))(s1)))  
  (s2)  
  (s3)  
  AND AF(N, g)(s3)
```

CONVERSION- K_conversion

END ctlops

```

fairctlops[state: TYPE]: THEORY
BEGIN

  u, v, w: VAR state

  f, g, Q, P, p1, p2: VAR pred[state]

  N: VAR [state, state → bool]

  Ff: VAR pred[state]

  CONVERSION+ K_conversion

  fairEG(N, f)(Ff): pred[state] =
    ν(λ P:
      EU(N, f,
        λ (s1: state):
          f(s1) AND
          (λ (s: state): Ff(s) AND EX(N, P)(s)
            (s1)))

  fairAF(N, f)(Ff): pred[state] =
    λ (s1: state):
      NOT fairEG(N, λ (s: state): NOT f(s))(Ff)(s1)

  fair?(N, Ff): pred[state] = fairEG(N, λ u: TRUE)(Ff)

  fairEX(N, f)(Ff): pred[state] =
    EX(N, λ (s: state): f(s) AND fair?(N, Ff)(s))

  fairEU(N, f, g)(Ff): pred[state] =
    EU(N, f, λ (s: state): g(s) AND fair?(N, Ff)(s))

  fairEF(N, f)(Ff): pred[state] =
    EF(N, λ (s: state): f(s) AND fair?(N, Ff)(s))

  fairAX(N, f)(Ff): pred[state] =
    λ (s1: state):
      NOT fairEX(N, λ (s: state): NOT f(s))(Ff)(s1)

  fairAG(N, f)(Ff): pred[state] =

```

```

λ (s1: state):
  NOT fairEF(N, λ (s: state): NOT f(s))(Ff)(s1)

fairAU(N, f, g)(Ff): pred[state] =
  λ (s3: state):
    (λ (s2: state):
      NOT fairEU(N, λ (s: state): NOT g(s),
        λ (s1: state):
          (λ (s: state): NOT f(s))(s1) AND
            (λ (s: state): NOT g(s))(s1))
        (Ff)(s2))
      (s3)
    AND fairAF(N, g)(Ff)(s3)

CONVERSION- K_conversion

END fairctlops

```


Fairctlops[state : TYPE] : THEORY

BEGIN

u, v, w : VAR state

f, g, Q, P, p_1, p_2 : VAR pred[state]

N : VAR [state, state \rightarrow bool]

Fflist, Gflist : VAR list[pred[state]]

CONVERSION+ K_conversion

CheckFair(Q, N, f, Fflist) : RECURSIVE pred[state] =

(CASES Fflist OF

cons(Ff, Gflist) :

EU($N, f,$

$\lambda (s_1 : \text{state}) :$

$f(s_1)$ AND

$(\lambda (s : \text{state}) : \text{Ff}(s)$ AND $\text{EX}(N, \text{CheckFair}(Q, N, f, \text{Gflist}))(s)$
 $(s_1)),$

null : Q

ENDCASES)

MEASURE length(Fflist)

FairEG(N, f)(Fflist) : pred[state] =

$\nu(\lambda P : \text{CheckFair}(P, N, f, \text{Fflist}))$

FairAF(N, f)(Fflist) : pred[state] =

$\lambda (s_1 : \text{state}) :$

NOT FairEG($N, \lambda (s : \text{state}) : \text{NOT } f(s)$)(Fflist)(s_1)

Fair?(N, Fflist) : pred[state] = FairEG($N, \lambda u : \text{TRUE}$)(Fflist)

FairEX(N, f)(Fflist) : pred[state] =

$\text{EX}(N, \lambda (s : \text{state}) : f(s)$ AND $\text{Fair?}(N, \text{Fflist})(s))$

FairEU(N, f, g)(Fflist) : pred[state] =

$\text{EU}(N, f, \lambda (s : \text{state}) : g(s)$ AND $\text{Fair?}(N, \text{Fflist})(s))$

FairEF(N, f)(Fflist) : pred[state] =

EF(N , $\lambda (s : \text{state}) : f(s)$) AND Fair?(N , Fflist)(s)

FairAX(N , f)(Fflist): pred[state] =
 $\lambda (s_1 : \text{state}) :$
NOT FairEX(N , $\lambda (s : \text{state}) : \text{NOT } f(s)$)(Fflist)(s_1)

FairAG(N , f)(Fflist): pred[state] =
 $\lambda (s_1 : \text{state}) :$
NOT FairEF(N , $\lambda (s : \text{state}) : \text{NOT } f(s)$)(Fflist)(s_1)

FairAU(N , f , g)(Fflist): pred[state] =
 $\lambda (s_3 : \text{state}) :$
($\lambda (s_2 : \text{state}) :$
NOT FairEU(N , $\lambda (s : \text{state}) : \text{NOT } g(s)$,
 $\lambda (s_1 : \text{state}) :$
($\lambda (s : \text{state}) : \text{NOT } f(s)$)(s_1) AND
($\lambda (s : \text{state}) : \text{NOT } g(s)$)(s_1))
(Fflist)(s_2))
(s_3)
AND FairAF(N , g)(Fflist)(s_3))

CONVERSION- K_conversion

END Fairctlops

```

bit: THEORY
BEGIN

  bit: TYPE = bool

  nbit: TYPE = below(2)

  b: VAR bit

  bit_cases: LEMMA  $b = \text{FALSE}$  OR  $b = \text{TRUE}$ 

  b0: [below(1) → bit] = (λ (i: below(1)): FALSE)

  b1: [below(1) → bit] = (λ (i: below(1)): TRUE)

  b2n(b: bool): nbit = IF  $b$  THEN 1 ELSE 0 ENDIF

  n2b(nb: nbit): bool = (nb = 1)

END bit

```

```

bv[N : ℕ] : THEORY
BEGIN

  CONVERSION+ b2n

  bvec : TYPE = [below(N) → bit]

  b : VAR bit

  bv : VAR bvec

  i : VAR below[N]

  bvec0(i) : bit = FALSE

  bvec1(i) : bit = TRUE

  fill(b)(i) : bit = b;

  bvi : bit = bv(i)

  CONVERSION- b2n

END bv

```

```

exp2: THEORY
BEGIN

  n, m, x1, x2: VAR ℕ

  exp2(n: ℕ): RECURSIVE ℕ>0 = IF n = 0 THEN 1 ELSE 2 × exp2(n - 1) ENDIF
    MEASURE n

  JUDGEMENT exp2(n: ℕ) HAS_TYPE above(n)

  exp2_def: LEMMA exp2(n) = 2n

  exp2_pos: LEMMA exp2(n) > 0

  exp2_n: LEMMA exp2(n + 1) = 2 × exp2(n)

  exp2_sum: LEMMA exp2(n + m) = exp2(n) × exp2(m)

  exp2_minus: LEMMA
    (∀ n, (k: upto(n))):
      exp2(n - k) = exp2(n)/exp2(k)

  exp2_strictpos: LEMMA n > 0 IMPLIES exp2(n) > 1

  exp2_lt: LEMMA n < m IMPLIES exp2(n) < exp2(m)

  exp_prop: LEMMA
    x1 < exp2(n) AND x2 < exp2(m) IMPLIES
      x1 × exp2(m) + x2 < exp2(n + m)

END exp2

```

```

bv_concat_def [n : ℕ, m : ℕ] : THEORY
BEGIN

  bvn : bvec [n] ∘ bvm : bvec [m] : bvec [n + m] =
    (λ (nm : below (n + m)) :
      IF nm < m THEN bvm (nm) ELSE bvn (nm - m) ENDIF)

END bv_concat_def

```

```

bv_bitwise[N: ℕ]: THEORY
BEGIN

  i: VAR below(N)

  OR(bv1, bv2: bvec[N]): bvec[N] =
    (λ i: bv1(i) OR bv2(i));

  AND(bv1, bv2: bvec[N]): bvec[N] =
    (λ i: bv1(i) AND bv2(i));

  IFF(bv1, bv2: bvec[N]): bvec[N] =
    (λ i: bv1(i) IFF bv2(i));

  NOT(bv: bvec[N]): bvec[N] = (λ i: NOT bv(i));

  XOR(bv1, bv2: bvec[N]): bvec[N] =
    (λ i: XOR(bv1(i), bv2(i)));

  bv, bv1, bv2: VAR bvec[N]

  bv_OR: LEMMA
    (bv1 OR bv2)i = (bv1i OR bv2i)

  bv_AND: LEMMA
    (bv1 AND bv2)i = (bv1i AND bv2i)

  bv_IFF: LEMMA
    (bv1 IFF bv2)i = (bv1i IFF bv2i)

  bv_XOR: LEMMA
    XOR(bv1, bv2)i = XOR(bv1i, bv2i)

  bv_NOT: LEMMA (NOT bv)i = NOT (bvi)

END bv_bitwise

```

bv_nat[$N : \mathbb{N}$]: THEORY

BEGIN

CONVERSION+ b2n

bv2nat_rec($n : \text{upto}(N)$, $\text{bv} : \text{bvec}[N]$): RECURSIVE $\mathbb{N} =$
IF $n = 0$
THEN 0
ELSE $\text{exp2}(n - 1) \times \text{b2n}(\text{bv}^{(n-1)}) + \text{bv2nat_rec}(n - 1, \text{bv})$
ENDIF
MEASURE n

bv_lem: LEMMA

$\forall (n : \text{below}(N), \text{bv} : \text{bvec}[N]) :$
 $\text{bv}(n) = \text{FALSE} \text{ OR } \text{bv}(n) = \text{TRUE}$

bv2nat_rec_bound: LEMMA

$\forall (n : \text{upto}(N), \text{bv} : \text{bvec}[N]) :$
 $\text{bv2nat_rec}(n, \text{bv}) < \text{exp2}(n)$

bv2nat($\text{bv} : \text{bvec}[N]$): $\text{below}(\text{exp2}(N)) = \text{bv2nat_rec}(N, \text{bv})$

n : VAR $\text{upto}(N)$

val: VAR $\text{below}(\text{exp2}(N))$

$\text{bv}, \text{bv1}, \text{bv2}$: VAR $\text{bvec}[N]$

bv2nat_inj_rec: LEMMA

$\text{bv2nat_rec}(n, \text{bv1}) = \text{bv2nat_rec}(n, \text{bv2}) \Leftrightarrow$
 $(\forall (m : \text{below}(N)) : (m < n) \text{ IMPLIES } \text{bv1}(m) = \text{bv2}(m))$

bv2nat_surj_rec: LEMMA

$\forall (y : \text{below}(\text{exp2}(n))) : \exists \text{bv} : \text{bv2nat_rec}(n, \text{bv}) = y$

bv2nat_inj: THEOREM

$(\forall (x, y : \text{bvec}[N]) :$
 $(\text{bv2nat}(x) = \text{bv2nat}(y) \text{ IMPLIES } (x = y)))$

bv2nat_surj: THEOREM

$(\forall (y : \text{below}(\text{exp2}(N))) :$

$(\exists (x : \text{bvec}[N]): \text{bv2nat}(x) = y)$

`bv2nat_bij`: THEOREM `bijjective?(bv2nat)`

`bv2nat_rec_fill_F`: LEMMA `bv2nat_rec(n, fill[N](FALSE)) = 0`

`bv2nat_rec_fill_T`: LEMMA
`bv2nat_rec(n, fill[N](TRUE)) = exp2(n) - 1`

`bv2nat_fill_F`: LEMMA `bv2nat(fill[N](FALSE)) = 0`

`bv2nat_fill_T`: LEMMA
`bv2nat(fill[N](TRUE)) = exp2(N) - 1`

`bv2nat_eq0`: LEMMA `bv2nat(bv) = 0 IMPLIES (bv = fill[N](FALSE))`

`bv2nat_eq_max`: LEMMA
`bv2nat(bv) = exp2(N) - 1 IMPLIES bv = (fill[N](TRUE))`

`bv2nat_top_bit`: THEOREM
`N > 0 IMPLIES`
`IF bv2nat(bv) < exp2(N - 1)`
`THEN bv(N-1) = FALSE`
`ELSE bv(N-1) = TRUE`
`ENDIF`

`bv2nat_topbit`: THEOREM
`N > 0 IMPLIES`
`bv(N-1) = (bv2nat(bv) ≥ exp2(N - 1))`

`nat2bv(val : below(exp2(N)))`:
`{bv : bvec[N] | bv2nat(bv) = val}`

`nat2bv_def`: LEMMA `nat2bv = inverse(bv2nat)`

`nat2bv_bij`: THEOREM
`bijjective?[below(exp2(N)), bvec[N]](nat2bv)`

`nat2bv_inv`: THEOREM `nat2bv(bv2nat(bv)) = bv`

`nat2bv_rew`: LEMMA `nat2bv(val) = bv IFF bv2nat(bv) = val`

bv2nat_inv: THEOREM $\text{bv2nat}(\text{nat2bv}(\text{val})) = \text{val}$

CONVERSION- b2n

END bv_nat

```
empty_bv: THEORY
BEGIN

  empty_bv: [below[0] → bool] = (λ (x: below[0]): TRUE);

END empty_bv
```

```

bv_caret[N : ℕ] : THEORY
BEGIN

  bv : bvec[N]sp: [upto(i1)] : bvec[PROJ_1(sp) - PROJ_2(sp) + 1] =
    (λ (ii : below(PROJ_1(sp) - PROJ_2(sp) + 1)) :
      bv(ii + PROJ_2(sp)));

  bv : VAR bvec[N]

  bv_caret_all : LEMMA N > 0 IMPLIES bv(N-1, 0) = bv

  bv_caret_ii_0 : LEMMA
    (∀ (i : below(N)) :
      bv(i, i)0 = bvi)

  bv_caret_elim : LEMMA
    (∀ (i : below(N), j : upto(i), k : below(i - j + 1)) :
      bv(i, j)k = bv(j+k))

END bv_caret

```

mod: THEORY

BEGIN

i, k, cc : VAR \mathbb{Z}

m : VAR $\mathbb{N}_{>0}$

n, a, b, c, x : VAR \mathbb{N}

j : VAR nonzero_integer

m13: LEMMA $|i - m \times \lfloor i/m \rfloor| < m$

m14: LEMMA $|i + m \times \lfloor -i/m \rfloor| < m$

mod(i, j): $\{k \mid |k| < |j|\}$ =
 $i - j \times \lfloor i/j \rfloor$

mod_pos: LEMMA mod(i, m) ≥ 0 AND mod(i, m) $< m$

JUDGEMENT mod(i : \mathbb{Z}, m : $\mathbb{N}_{>0}$) HAS_TYPE below(m)

mod_even: LEMMA integer_pred(i/j) IMPLIES mod(i, j) = 0

mod_neg: LEMMA

mod($-i, j$) =
IF integer_pred(i/j)
THEN 0
ELSE $j - \text{mod}(i, j)$
ENDIF

mod_neg_d: LEMMA

mod($i, -j$) =
IF integer_pred(i/j)
THEN 0
ELSE mod(i, j) - j
ENDIF

mod_eq_arg: LEMMA mod(j, j) = 0

mod_lt: LEMMA

$|i| < |j|$ IMPLIES

```

mod(i, j) =
  IF sgn(i) = sgn(j) OR i = 0
    THEN i
  ELSE i + j
  ENDIF

```

mod_lt_nat: LEMMA $n < m$ IMPLIES $\text{mod}(n, m) = n$

```

mod_lt_int: LEMMA
  -m < i AND i < m IMPLIES
  mod(i, m) = IF i ≥ 0 THEN i ELSE i + m ENDIF

```

mod_sum_pos: LEMMA $\text{mod}(i + k \times m, m) = \text{mod}(i, m)$

```

mod_gt: LEMMA
  m ≤ i AND i < 2 × m IMPLIES mod(i, m) = i - m

```

mod_sum: LEMMA $\text{mod}(i + k \times j, j) = \text{mod}(i, j)$

```

mod_sum_nat: LEMMA
  (∀ (n1, n2: below(m)):
    mod(n1 + n2, m) =
      IF n1 + n2 < m
        THEN n1 + n2
      ELSE n1 + n2 - m
    ENDIF)

```

```

mod_it_is: LEMMA
  a = b + m × c AND b < m IMPLIES b = mod(a, m)

```

mod_zero: LEMMA $\text{mod}(0, j) = 0$

```

mod_one: LEMMA
  mod(1, j) =
    IF |j| = 1
      THEN 0
    ELSIF j > 0 THEN 1
    ELSE j + 1
  ENDIF

```

```

mod_of_mod: LEMMA
  mod(i + mod(k, m), m) = mod(i + k, m)

```

mod_of_mod_neg: LEMMA

$$\text{mod}(i - \text{mod}(k, m), m) = \text{mod}(i - k, m)$$

mod_inj_plus: LEMMA

$$a < m \text{ AND } n < m \text{ AND } c < m \text{ AND } \text{mod}(a + n, m) = \text{mod}(a + c, m) \text{ IMPLIES } n = c$$

mod_inj_minus: LEMMA

$$a < m \text{ AND } n < m \text{ AND } c < m \text{ AND } \text{mod}(a - n, m) = \text{mod}(a - c, m) \text{ IMPLIES } n = c$$

mod_wrap_around: LEMMA

$$n < m \text{ AND } (c \leq m \text{ AND } c \geq m - n) \text{ IMPLIES } \text{mod}(n + c, m) = n - (m - c)$$

mod_wrap2: LEMMA $c < m$ IMPLIES $\text{mod}(m + c, m) = c$

mod_inj1: LEMMA

$$x < m \text{ AND } n < m \text{ AND } c < m \text{ AND } \text{mod}(x + n, m) = \text{mod}(x + c, m) \text{ IMPLIES } n = c$$

mod_inj2: LEMMA

$$x < m \text{ AND } n < m \text{ AND } c < m \text{ AND } \text{mod}(x - n, m) = \text{mod}(x - c, m) \text{ IMPLIES } n = c$$

mod_wrap_inj: LEMMA

$$n < m \text{ AND } a < m \text{ AND } b < m \text{ AND } a > 0 \text{ AND } \text{mod}(n + a, m) = \text{mod}(n - b, m) \text{ IMPLIES } a + b = m$$

mod_wrap_inj_eq: LEMMA

$$x < m \text{ AND } a < m \text{ AND } b < m \text{ AND } a > 0 \text{ IMPLIES } (\text{mod}(x + a, m) = \text{mod}(x - b, m)) = (a + b = m)$$

kk, vv: VAR \mathbb{N}

mod_neg_limited: LEMMA

$$0 \leq \text{kk} \text{ AND } \text{kk} < m \text{ AND } \text{vv} < m \text{ AND } \text{vv} - \text{kk} < 0 \text{ IMPLIES } \text{mod}(\text{vv} - \text{kk}, m) = m + \text{vv} - \text{kk}$$

odd_mod: LEMMA $\text{even?}(m) \Rightarrow (\text{odd?}(\text{mod}(i, m)) \text{ IFF } \text{odd?}(i))$

even_mod: LEMMA
 $\text{even?}(m) \Rightarrow (\text{even?}(\text{mod}(i, m)) \text{ IFF } \text{even?}(i))$

mj: VAR $\mathbb{N}_{>0}$

mod_mult: LEMMA $\text{mod}(\text{mod}(i, \text{mj} \times m), m) = \text{mod}(i, m)$

END mod


```

bv_arith_nat_defs[N: ℕ]: THEORY
BEGIN

  <(bv1, bv2: bvec[N]): bool = bv2nat(bv1) < bv2nat(bv2);

  ≤(bv1, bv2: bvec[N]): bool = bv2nat(bv1) ≤ bv2nat(bv2);

  >(bv1, bv2: bvec[N]): bool = bv2nat(bv1) > bv2nat(bv2);

  ≥(bv1, bv2: bvec[N]): bool = bv2nat(bv1) ≥ bv2nat(bv2);

  bv : bvec[N] + i : ℤ:
    {bv: bvec[N] |
      bv2nat(bv) =
        mod(bv2nat(bv) + i, exp2(N))}

  bv_plus: LEMMA
    (∀ (bv: bvec[N], i: ℤ):
      bv2nat(bv + i) =
        mod(bv2nat(bv) + i, exp2(N)));

  bv : bvec[N] - i : ℤ: bvec[N] = bv + (-i);

  bv_minus: LEMMA
    (∀ (bv: bvec[N], i: ℤ):
      bv2nat(bv - i) =
        mod(bv2nat(bv) - i, exp2(N)));

  bv1 : bvec[N] + bv2 : bvec[N]:
    {bv: bvec[N] |
      bv2nat(bv) =
        IF bv2nat(bv1) + bv2nat(bv2) < exp2(N)
          THEN bv2nat(bv1) + bv2nat(bv2)
          ELSE bv2nat(bv1) + bv2nat(bv2) - exp2(N)
          ENDIF}

  bv1 : bvec[N] × bv2 : bvec[N]:
    {bv: bvec[2 × N] |
      bv2nat(bv) = bv2nat(bv1) × bv2nat(bv2)};

END bv_arith_nat_defs

```

```

bv_int_defs[N : ℕ>0]: THEORY
BEGIN

  minint: ℤ = -exp2(N - 1)

  maxint: ℤ = exp2(N - 1) - 1

  bv_maxint_to_minint: LEMMA maxint = -minint - 1

  bv_minint_to_maxint: LEMMA minint = -maxint - 1

  in_rng_2s_comp(i: ℤ): bool = (minint ≤ i AND i ≤ maxint)

  rng_2s_comp: TYPE = {i: ℤ | minint ≤ i AND i ≤ maxint}

  bv2int(bv: bvec[N]): rng_2s_comp =
    IF bv2nat(bv) < exp2(N - 1)
      THEN bv2nat(bv)
    ELSE bv2nat(bv) - exp2(N)
    ENDIF

  int2bv(iv: rng_2s_comp): {bv: bvec[N] | bv2int(bv) = iv}

END bv_int_defs

```

```

bv_arithmetic_defs[N: ℕ>0]: THEORY
BEGIN

bv, bv1, bv2: VAR bvec[N]

-bv : bvec[N] :
    {bvn: bvec[N] |
      bv2int(bvn) =
        IF bv2int(bv) = minint[N]
          THEN bv2int(bv)
          ELSE -(bv2int(bv))
        ENDIF}

bv1 - bv2: bvec[N] = (bv1 + (-bv2))

overflow(bv1, bv2): bool =
  (bv2int(bv1) + bv2int(bv2)) > maxint[N] OR
  (bv2int(bv1) + bv2int(bv2)) < minint[N]

bv_slt(bv1, bv2): bool = bv2int(bv1) < bv2int(bv2)

bv_sle(bv1, bv2): bool = bv2int(bv1) ≤ bv2int(bv2)

bv_sgt(bv1, bv2): bool = bv2int(bv1) > bv2int(bv2)

bv_sge(bv1, bv2): bool = bv2int(bv1) ≥ bv2int(bv2)

bv_splus(bv1, (bv2: bvec[N] | NOT overflow(bv1, bv2))): bvec[N] =
  int2bv(bv2int(bv1) + bv2int(bv2))

mult_overflow(bv1, bv2): bool =
  (bv2int(bv1) × bv2int(bv2)) > maxint[N] OR
  (bv2int(bv1) × bv2int(bv2)) < minint[N]

bv_stimes(bv1, (bv2: bvec[N] | NOT mult_overflow(bv1, bv2))): bvec[N] =
  int2bv(bv2int(bv1) × bv2int(bv2))

END bv_arithmetic_defs

```

```

bv_extend_defs [N : ℕ>0] : THEORY
BEGIN

  bv : VAR bvec [N]

  k : VAR above(N)

  zero_extend(k : above(N)) : [bvec [N] → bvec [k]] =
    (λ bv : fill [k - N] (FALSE) ∘ bv)

  sign_extend(k : above(N)) : [bvec [N] → bvec [k]] =
    (λ bv : fill [k - N] (bv(N-1)) ∘ bv)

  zero_extend_lsend(k : above(N)) : [bvec [N] → bvec [k]] =
    (λ bv : bv ∘ fill [k - N] (FALSE))

  lsb_extend(k : above(N)) : [bvec [N] → bvec [k]] =
    (λ bv : bv ∘ fill [k - N] (bv0))

  pad_left(k : above(N), b : bit)(bv) : bvec [k] =
    fill [k - N] (b) ∘ bv

  pad_right(k : above(N), b : bit)(bv) : bvec [k] =
    bv ∘ fill [k - N] (b)

END bv_extend_defs

```

```

infinite_sets_def[T : TYPE] : THEORY
BEGIN

  S, R : VAR set[T]

  is_infinite(S) : MACRO bool = NOT is_finite(S)

  infinite_set : TYPE = {S | NOT is_finite(S)}

  Inf : VAR infinite_set

  Fin : VAR finite_set[T]

  t : VAR T

  infinite_nonempty : JUDGEMENT infinite_set SUBTYPE_OF (nonempty?[T])

  infinite_add : JUDGEMENT add(t, Inf) HAS_TYPE infinite_set

  infinite_remove : JUDGEMENT remove(t, Inf) HAS_TYPE infinite_set

  infinite_superset : THEOREM
     $\forall \text{ Inf}, S : (\text{Inf} \subseteq S) \Rightarrow \text{NOT is\_finite}(S)$ 

  infinite_union_left : JUDGEMENT union(Inf, S) HAS_TYPE infinite_set

  infinite_union_right : JUDGEMENT union(S, Inf) HAS_TYPE infinite_set

  infinite_union : THEOREM
     $\forall S, R :$ 
     $\text{NOT is\_finite}((S \cup R)) \Rightarrow$ 
     $\text{NOT is\_finite}(S) \text{ OR NOT is\_finite}(R)$ 

  infinite_intersection : THEOREM
     $\forall S, R :$ 
     $\text{NOT is\_finite}((S \cap R)) \Rightarrow$ 
     $\text{NOT is\_finite}(S) \text{ AND NOT is\_finite}(R)$ 

  infinite_difference : JUDGEMENT difference(Inf, Fin) HAS_TYPE
    infinite_set

  infinite_rest : JUDGEMENT rest(Inf) HAS_TYPE infinite_set

```

infinite_fullset: THEOREM

$(\exists S: \text{NOT is_finite}(S)) \Rightarrow \text{NOT is_finite}(\text{fullset}[T])$

END infinite_sets_def

```

finite_sets_of_sets[T : TYPE] : THEORY
BEGIN

  a : VAR set[T]

  powerset_natfun_rec(A : finite_set[T], n : upto(card(A)),
                    f : (bijective?[A, below(card(A))]),
                    B : (powerset(A))) : RECURSIVE

     $\mathbb{N} =$ 
    IF n = 0
    THEN 0
    ELSE LET nval =
          exp2(n - 1)  $\times$  IF (inverse(f)(n - 1)  $\in$  B) THEN 1 ELSE 0 ENDIF
        IN nval + powerset_natfun_rec(A, n - 1, f, B)
    ENDIF

    MEASURE n

  powerset_natfun_rec_bound : LEMMA
     $\forall$  (A : finite_set[T], n : upto(card(A)),
        f : (bijective?[A, below(card(A))]), B : (powerset(A))) :
      powerset_natfun_rec(A, n, f, B) < exp2(n)

  powerset_natfun(A : finite_set[T])(B : (powerset(A))) : below(exp2(card(A))) =
    LET f = choose(bijective?[A, below(card(A))]) IN
      powerset_natfun_rec(A, card(A), f, B)

  powerset_natfun_inj_rec : LEMMA
     $\forall$  (A : finite_set[T], n : upto(card(A)),
        f : (bijective?[A, below(card(A))]), B1, B2 : (powerset(A))) :
      powerset_natfun_rec(A, n, f, B1) = powerset_natfun_rec(A, n, f, B2)  $\Leftrightarrow$ 
      ( $\forall$  (m : upto(card(A))) :
        (m < n) IMPLIES
          ((inverse(f)(m)  $\in$  B1) IFF
            (inverse(f)(m)  $\in$  B2)))

  powerset_natfun_inj : LEMMA
     $\forall$  (A : finite_set[T]) :
       $\forall$  (B1, B2 : (powerset(A))) :
        powerset_natfun(A)(B1) = powerset_natfun(A)(B2) IMPLIES
          B1 = B2

  powerset_finite : JUDGEMENT powerset(A : finite_set[T]) HAS_TYPE

```

finite_set[set[T]]

SS: VAR setofsets[T]

Union_finite: THEOREM

∀ SS:
is_finite(∪ SS) IFF
is_finite(SS) AND every(is_finite)(SS)

finite_Union_finite: LEMMA

is_finite(SS) AND every(is_finite[T])(SS) IMPLIES
is_finite(∪ SS)

Union_infinite: COROLLARY

∀ SS:
NOT is_finite(∪ SS) IFF
NOT is_finite(SS) OR
some(λ (S: set[T]): NOT is_finite(S))(SS)

Intersection_finite: THEOREM

∀ SS:
nonempty?(SS) AND every(is_finite)(SS) ⇒
is_finite(∩ SS)

Intersection_infinite: COROLLARY

∀ SS:
NOT is_finite(∩ SS) ⇒
every(λ (S: set[T]): NOT is_finite(S))(SS)

Complement_finite: THEOREM

∀ SS: is_finite(Complement(SS)) IFF is_finite(SS)

Complement_is_finite: JUDGEMENT Complement(SS: finite_set[set[T]]) HAS_TYPE
finite_set[set[T]]

Complement_infinite: COROLLARY

∀ SS: NOT is_finite(Complement(SS)) IFF NOT is_finite(SS)

Complement_is_infinite: JUDGEMENT Complement(SS: infinite_set[set[T]]) HAS_TYPE
infinite_set[set[T]]

END finite_sets_of_sets


```

EquivalenceClosure [T: TYPE]: THEORY
BEGIN

  R, S: VAR PRED [[T]]

  x, y: VAR T

  EquivClos(R): equivalence [T] =
    {x, y |
      ∀ (S: equivalence [T]): (R ⊆ S) IMPLIES S(x, y)}

  EquivClosSuperset: LEMMA (R ⊆ EquivClos(R))

  EquivClosMonotone: LEMMA
    (R ⊆ S) IMPLIES (EquivClos(R) ⊆ EquivClos(S))

  EquivClosLeast: LEMMA
    equivalence?(S) AND (R ⊆ S) IMPLIES (EquivClos(R) ⊆ S)

  EquivClosIdempotent: LEMMA
    EquivClos(EquivClos(R)) = EquivClos(R)

  EquivalenceCharacterization: LEMMA
    equivalence?(S) IFF (S = EquivClos(S))

END EquivalenceClosure

```

```

QuotientDefinition [T: TYPE]: THEORY
BEGIN

  R: VAR set[[T]]

  S: VAR equivalence[T]

  x, y, z: VAR T

  EquivClass(R)(x): set[T] = {z | R(x, z)}

  EquivClassNonEmpty: LEMMA nonempty?[T](EquivClass(S)(x))

  EquivClassEq: LEMMA
    EquivClass(S)(x) = EquivClass(S)(y) IFF S(x, y)

  repEC(S)(x): T = choose(EquivClass(S)(x))

  EquivClassChoose: LEMMA S(x, repEC(S)(x))

  ChooseEquivClassChoose: LEMMA
    EquivClass(S)(repEC(S)(x)) = EquivClass(S)(x)

  Quotient(S): TYPE =
    {P: set[T] | ∃ x: P = EquivClass(S)(x)}

  rep(S)(P: Quotient(S)): T = choose(P)

  rep_is_repEC: LEMMA
    rep(S)(EquivClass(S)(x)) = repEC(S)(x)

  rep_lemma: LEMMA
    EquivClass(S)(x)(rep(S)(EquivClass(S)(x)))

  quotient_map(S)(x): Quotient(S) = EquivClass(S)(x)

  quotient_map_surjective: LEMMA surjective?(quotient_map(S))

  ECQuotient(R): TYPE = Quotient(EquivClos(R))

  ECquotient_map(R)(x): ECQuotient(R) =
    quotient_map(EquivClos(R))(x)

```

END QuotientDefinition

```

KernelDefinition[X : TYPE, X1 : TYPE FROM X, Y : TYPE] : THEORY
BEGIN

  f : VAR [X1 → Y]

  R : VAR PRED[[X]]

  x1, x2 : VAR X1

  EquivalenceKernel(f) : equivalence[X1] =
    {x1, x2 | f(x1) = f(x2)}

  PreservesEq(R)(f) : bool =
    (restrict [[X], [X1], bool](R) ⊆ EquivalenceKernel(f))

  PreservesEqClosure : LEMMA
  PreservesEq(R) =
    PreservesEq(extend[[X], [X1], bool, FALSE]
      (EquivClos[X1]
        (restrict[[X], [X1], boolean]
          (R))))

  PreservesEq_is_preserving : LEMMA
  PreservesEq(R) =
    preserves(restrict[[X], [X1], bool](R), =[Y])

END KernelDefinition

```

```

QuotientKernelProperties[X : TYPE, X1 : TYPE FROM X] : THEORY
BEGIN

  S : VAR equivalence[X1]

  R : VAR PRED[[X]]

  Kernel_quotient_map : LEMMA
    EquivalenceKernel[X, X1, Quotient(S)](quotient_map(S)) = S

  PreservesEq_quotient_map : LEMMA
    PreservesEq[X, X1, Quotient(S)]
      (extend[[X], [X1], bool, FALSE](S))
      (quotient_map(S))

  quotient_map_is_Quotient_EquivalenceRespecting : JUDGEMENT quotient_map(S) HAS_TYPE
    (PreservesEq[X, X1, Quotient(S)]
      (extend[[X], [X1], bool, FALSE](S)))

  Kernel_ECquotient_map : LEMMA
    EquivalenceKernel[X, X1, ECQuotient(S)](quotient_map(S)) =
      S

  PreservesEq_ECquotient_map : LEMMA
    PreservesEq[X, X1, ECQuotient(S)]
      (extend[[X], [X1], bool, FALSE](S))
      (quotient_map(S))

  quotient_map_is_ECQuotient_EquivalenceRespecting : JUDGEMENT quotient_map(S) HAS_TYPE
    (PreservesEq[X, X1, ECQuotient(S)]
      (extend[[X], [X1], bool, FALSE](S)))

END QuotientKernelProperties

```

```

QuotientSubDefinition[X : TYPE, X1 : TYPE FROM X] : THEORY
BEGIN

  x : VAR X1

  S : VAR
    {R : equivalence[X] |
      PreservesEq[X, X, bool](R)(X1_pred)}

  QuotientSub(S) : TYPE =
    {P : set[X] | ∃ x : P = EquivClass(S)(x)}

  quotient_sub_map(S)(x) : QuotientSub(S) = EquivClass(S)(x)

END QuotientSubDefinition

```

```

QuotientExtensionProperties[X : TYPE, X1 : TYPE FROM X, Y : TYPE] : THEORY
BEGIN

  S : VAR
    {R : equivalence[X] |
      PreservesEq[X, X, bool](R)(X1_pred)}

  lift(S)(g : (PreservesEq[X, X1, Y](S)))(P : QuotientSub[X, X1](S)) : Y =
    g(rep(S)(P))

  lift_commutation : LEMMA
    ∀ S, (g : (PreservesEq[X, X1, Y](S))):
      lift(S)(g) ∘ quotient_sub_map[X, X1](S) = g

  lift_unicity : LEMMA
    ∀ ((S : {R : equivalence[X] | PreservesEq[X, X, bool](R)(X1_pred)}
      | PreservesEq[X, X, bool](S)(X1_pred)),
      (g : (PreservesEq[X, X1, Y](S))):
      ∀ (h : [QuotientSub[X, X1](S) → Y]):
        h ∘ quotient_sub_map[X, X1](S) = g IMPLIES
          h = lift(S)(g)

END QuotientExtensionProperties

```

```

QuotientDistributive[X, Y : TYPE] : THEORY
BEGIN

  S : VAR equivalence[X]

  z, w : VAR [Y]

  EqualityExtension(S) : set[[[Y]]] =
    {z, w | S(z`1, w`1) AND z`2 = w`2}

  EqualityExtension_is_equivalence : JUDGEMENT EqualityExtension(S) HAS_TYPE
    equivalence[[Y]]

  EqualityExtensionPreservesEq : LEMMA
    PreservesEq[[Y], [Y], [Y]]
      (EqualityExtension(S))
      (λ (x : X, y : Y) : (quotient_map(S)(x), y))

  QuotientDistributive : LEMMA
    bijective?[Quotient(EqualityExtension(S)), [Y]]
      (lift[[Y], [Y], [Y]]
        (EqualityExtension(S))
        (λ (x : X, y : Y) : (quotient_map(S)(x), y)))

  R : VAR equivalence[Y]

  RelExtension(S, R) : equivalence[[Y]] =
    {z, w | S(z`1, w`1) AND R(z`2, w`2)}

  RelExtensionPreservesEq : LEMMA
    PreservesEq[[Y], [Y], [Quotient(R)]]
      (RelExtension(S, R))
      (λ (x : X, y : Y) :
        (quotient_map(S)(x), quotient_map(R)(y)))

  RelQuotientDistributive : LEMMA
    bijective?[Quotient(RelExtension(S, R)), [Quotient(R)]]
      (lift[[Y], [Y], [Quotient(R)]]
        (RelExtension(S, R))
        (λ (x : X, y : Y) :
          (quotient_map(S)(x),
            quotient_map(R)(y))))

```


$F: \text{VAR } [X \rightarrow \text{equivalence}[Y]]$

$f, g: \text{VAR } [X \rightarrow Y]$

$\text{FunExtension}(F): \text{equivalence}[[X \rightarrow Y]] =$
 $\{f, g \mid \forall (x: X): F(x)(f(x), g(x))\}$

FunExtensionPreservesEq: LEMMA

$\text{PreservesEq}[[X \rightarrow Y], [X \rightarrow Y], [x: X \rightarrow \text{Quotient}(F(x))]]$
 $(\text{FunExtension}(F))$
 $(\lambda (f: [X \rightarrow Y]):$
 $\lambda (x: X): \text{quotient_map}(F(x))(f(x)))$

FunQuotientDistributive: LEMMA

$\text{bijective?}[\text{Quotient}(\text{FunExtension}(F)), [x: X \rightarrow \text{Quotient}(F(x))]]$
 $(\text{lift}[[X \rightarrow Y], [X \rightarrow Y], [x: X \rightarrow \text{Quotient}(F(x))]]$
 $(\text{FunExtension}(F))$
 $(\lambda (f: [X \rightarrow Y]):$
 $(\lambda (x: X):$
 $(\text{quotient_map}(F(x))(f(x))))))$

END QuotientDistributive

```

QuotientIteration [X : TYPE] : THEORY
BEGIN

  S : VAR equivalence [X]

  x, y : VAR X

  action(S)(R : equivalence [Quotient(S)])(x, y) : bool =
    R(EquivClass(S)(x), EquivClass(S)(y))

  action_equivalence_is_equivalence : JUDGEMENT action(S)(R : equivalence [Quotient(S)])
    HAS_TYPE equivalence [X]

  QuotientAction : LEMMA
    ∀ (R : equivalence [Quotient(S)]):
      bijective? [Quotient(R), Quotient(action(S)(R))]
        (lift [Quotient(S), Quotient(S), Quotient(action(S)(R))]
          (R)
          (lift [X, X, Quotient(action(S)(R))]
            (S)
            (quotient_map [X] (action(S)(R))))))

END QuotientIteration

```

```

PartialFunctionDefinitions[X, Y: TYPE]: THEORY
BEGIN

  SubsetPartialFunction: TYPE = [#dom: PRED[X], fun: [(dom) → Y]#]

  LiftPartialFunction: TYPE = [X → lift[Y]]

  f: VAR LiftPartialFunction

  g: VAR SubsetPartialFunction

  h: VAR [X → Y]

  SPartFun_appl(g): [(dom(g)) → Y] = g`fun

  SPartFun_to_LPartFun(g): LiftPartialFunction =
    λ (x: X):
      IF dom(g)(x) THEN up(fun(g)(x)) ELSE bottom ENDIF

  LPartFun_to_SPartFun(f): SubsetPartialFunction =
    (#dom := {x: X | up?(f(x))},
     fun := λ (y: {x: X | up?(f(x))}): down(f(y))#)

  TotalFun_to_SPartFun(h): SubsetPartialFunction =
    (#dom := {x: X | TRUE}, fun := h#)

  TotalFun_to_LPartFun(h): LiftPartialFunction =
    λ (x: X): up(h(x))

  CONVERSION SPartFun_appl

  CONVERSION SPartFun_to_LPartFun

  CONVERSION LPartFun_to_SPartFun

  CONVERSION TotalFun_to_SPartFun

  CONVERSION TotalFun_to_LPartFun

```

SPartFun_to_LPartFun_to_SPartFun: LEMMA

$$\text{LPartFun_to_SPartFun}(\text{SPartFun_to_LPartFun}(g)) = g$$

LPartFun_to_SPartFun_to_LPartFun: LEMMA

$$\text{SPartFun_to_LPartFun}(\text{LPartFun_to_SPartFun}(f)) = f$$

END PartialFunctionDefinitions

```

PartialFunctionComposition[X, Y, Z: TYPE]: THEORY
BEGIN

  f: VAR LiftPartialFunction[X, Y]

  g: VAR LiftPartialFunction[Y, Z]

  g ∘ f: LiftPartialFunction[X, Z] =
    λ (x: X):
      CASES f(x) OF bottom: bottom, up(y): g(y) ENDCASES

  h: VAR SubsetPartialFunction[X, Y]

  k: VAR SubsetPartialFunction[Y, Z]

  CompDom(k, h): PRED[X] =
    {x: X | dom(h)(x) AND dom(k)(fun(h)(x))};

  k ∘ h: SubsetPartialFunction[X, Z] =
    (#dom := CompDom(k, h),
     fun
      := λ (x: (CompDom(k, h))):
         fun(k)(fun(h)(x))#)

  SPartFun_to_LPartFun_CompositionPreservation: LEMMA
    SPartFun_to_LPartFun(k ∘ h) =
      SPartFun_to_LPartFun(k) ∘ SPartFun_to_LPartFun(h)

  LPartFun_to_SPartFun_CompositionPreservation: LEMMA
    LPartFun_to_SPartFun(g ∘ f) =
      LPartFun_to_SPartFun(g) ∘ LPartFun_to_SPartFun(f)

END PartialFunctionComposition

```

```
stdlang: THEORY
BEGIN

  void: TYPE = bool

  skip: void = TRUE

  fail: void = FALSE

  try( $s_1, s_2$ : void): MACRO void =  $s_1$  OR  $s_2$ 

  try( $s$ : void): MACRO void =  $s$  OR skip

  ifthen( $b$ : bool,  $s$ : void): MACRO void = IF  $b$  THEN  $s$  ELSE skip ENDIF

  ifelse( $b$ : bool,  $s$ : void): MACRO void = IF  $b$  THEN skip ELSE  $s$  ENDIF

  Dummy: TYPE = bool

  dummy: MACRO Dummy = FALSE

END stdlang
```

```
stdexc[T: TYPE+]: THEORY
BEGIN

  ExceptionTag: TYPE = string

  Exception: TYPE = [#tag: ExceptionTag, val: T#]

  make_exc(e: ExceptionTag, t: T): Exception =
    (#tag := e, val := t#)

END stdexc
```

```

stdcatch[T1, T2: TYPE+]: THEORY
BEGIN

  catch_lift(tag: ExceptionTag[T2], t1: [Dummy → T1], t2: [Exception[T2] → T1]):
    T1

  catch(tag: ExceptionTag[T2], t1: T1, t2: [Exception[T2] → T1]): MACRO T1 =
    catch_lift(tag, λ (d: Dummy): t1, t2)

  catch_list_lift(l: list[ExceptionTag[T2]], f1: [Dummy → T1],
    f2: [Exception[T2] → T1]): RECURSIVE
    T1 =
  CASES l OF
    null: f1(FALSE),
    cons(e, r): catch_lift(e, λ (d: Dummy): catch_list_lift(r, f1, f2), f2)
  ENDCASES
  MEASURE l BY <<

  catch(l: list[ExceptionTag[T2]], t1: T1, t2: [Exception[T2] → T1]): MACRO T1 =
    catch_list_lift(l, λ (d: Dummy): t1, t2)

  throw(tag: ExceptionTag[T2], e: Exception[T2]): T1

  throw(tag: ExceptionTag[T2], val: T2): MACRO T1 =
    throw(tag, make_exc(tag, val))

END stdcatch

```



```

stdprog[T: TYPE+]: THEORY
BEGIN

  prog(s: void, t: T): T = t

  error(mssg: string): T

  exit: T

  catch(tag: ExceptionTag[void], t1, t2: T): MACRO T =
    catch_lift[T, void]
      (tag, λ (d: Dummy): t1, λ (e: Exception[void]): t2)

  throw(tag: ExceptionTag[void]): MACRO T =
    throw(tag, make_exc(tag, fail))

  catch(l: list[ExceptionTag[void]], t1, t2: T): MACRO T =
    catch_list_lift[T, void]
      (l, λ (d: Dummy): t1, λ (e: Exception[void]): t2)

  UndefinedMutableVariable: ExceptionTag[void] =
    "UndefinedMutableVariable"

  Mutable: TYPE+

  ref(t: T): Mutable

  new: Mutable

  undef(v: Mutable): bool

  val_lisp(v: Mutable): T

  val(v: Mutable): T =
    IF undef(v)
      THEN throw(UndefinedMutableVariable, make_exc(UndefinedMutableVariable, fail))
    ELSE val_lisp(v)
    ENDIF

  def(v: Mutable, t: T): T = t

  set(v: Mutable, t: T): void = LET nt = def(v, t) IN skip

```

```
CONVERSION val
```

```
Global: TYPE+ = Mutable
```

```
loop_lift(f: [Dummy → void]): T
```

```
loop(s: void): MACRO T = loop_lift( $\lambda$  (d: Dummy): s)
```

```
return(t: T): void = fail
```

```
format(s: string, t: T): string
```

```
END stdprog
```

```
stdglobal[T : TYPE+, t : T]: THEORY
BEGIN

  Global: TYPE+ = Mutable[T]

END stdglobal
```

```

stdpvs[T: TYPE+]: THEORY
BEGIN

  typeof(t: T): string

  str2pvs(s: string): T

  pvs2str_lisp(t: T): string

  pvs2str(t: T): MACRO string =
    catch_lift("cant-translate", λ (d: Dummy): pvs2str_lisp(t),
              λ (e: Exception[void]): "<?>")

  Slisp: TYPE+

  slisp(l: list[T]): Slisp

  {||}(l: list[T]): Slisp = slisp(l)

END stdpvs

```

```

stdstr: THEORY
BEGIN

  NotARealNumber: ExceptionTag[string] = "NotARealNumber"

  NotAnInteger: ExceptionTag[string] = "NotAnInteger"

  charcode( $n: \mathbb{N}$ ): string

  chartable: void

  emptystr: string = ""

  space: string = " "

  newline: string

  tab: string

  doublequote: string = charcode(34)

  singlequote: string = "'"

  backquote: string = "`"

  spaces( $n: \mathbb{N}$ ): string

  upcase( $s: string$ ): string

  downcase( $s: string$ ): string

  capitalize( $s: string$ ): string

  strfind( $s_1, s_2: string$ ):  $\mathbb{Z}$ 

  substr( $s: string, i, j: \mathbb{N}$ ): string

  real2str( $r: \mathbb{R}$ ): string

  bool2str( $b: bool$ ): string = IF  $b$  THEN "TRUE" ELSE "FALSE" ENDIF

  tostr( $r: \mathbb{R}$ ): MACRO string = real2str( $r$ )

```

```

tostr(b: bool): MACRO string = bool2str(b)

str2real(s: string):  $\mathbb{Q}$ 

str2int(s: string):  $\mathbb{Z}$ 

str2bool(s, answer: string): bool =
  downcase(s) = downcase(answer)

number?(s: string): bool

int?(s: string): bool

concat(s1, s2: string): string = s1 ◦ s2;

s1: string + s2: string: MACRO string = concat(s1, s2);

r:  $\mathbb{R}$  + s: string: MACRO string = concat(real2str(r), s);

s: string + r:  $\mathbb{R}$ : MACRO string = concat(s, real2str(r));

b: bool + s: string: MACRO string =
  concat(bool2str(b), s);

s: string + b: bool: MACRO string =
  concat(s, bool2str(b))

pad(n:  $\mathbb{N}$ , s: string): RECURSIVE string =
  IF n = 0 THEN emptystr ELSE concat(s, pad(n - 1, s)) ENDIF
  MEASURE n

strcmp(s1, s2: string, sensitive: bool):  $\mathbb{Z}$ 

strcmp(s1, s2: string): MACRO  $\mathbb{Z}$  = strcmp(s1, s2, TRUE)

strtrim(s1, s2: string): string

strtrim_left(s1, s2: string): string

strtrim_right(s1, s2: string): string

```

`trim(s: string): string`

`trim_left(s: string): string`

`trim_right(s: string): string`

`filename(s: string): string`

`directory(s: string): string`

`END stdstr`

stdio: THEORY

BEGIN

FileNotFound: ExceptionTag[string] = "FileNotFound"

FileAlreadyExists: ExceptionTag[string] = "FileAlreadyExists"

WrongInputStreamMode: ExceptionTag[string] = "WrongInputStreamMode"

WrongOutputStreamMode: ExceptionTag[string] = "WrongOutputStreamMode"

ClosedStream: ExceptionTag[string] = "ClosedStream"

EndOfFile: ExceptionTag[string] = "EndOfFile"

IOExceptionTags: list[ExceptionTag[string]] =
(:NotARealNumber, NotAnInteger, FileNotFound, FileAlreadyExists, WrongInput-
StreamMode,
WrongOutputStreamMode, ClosedStream, EndOfFile:)

assert(*b*: bool, *str*: string): void =
 b OR error[void](concat("[Assertion Failure] ", *str*)) & fail

break: MACRO void = return(skip)

while(*b*: bool, *s*: void): MACRO void =
 loop_lift(λ (*d*: Dummy):
 IF *b* THEN *s* ELSE return(skip) ENDIF)

for(*si*, *b*, *sinc*, *s*: void): MACRO void =
 si &
 loop_lift(λ (*d*: Dummy):
 IF *b* THEN *s* & *sinc* ELSE return(skip) ENDIF)

printstr(*s*: string): void = skip

print(*s*: string): MACRO void = printstr(*s*)

print(*r*: \mathbb{R}): MACRO void =
 printstr(concat(real2str(*r*), emptystr))

print(*b*: bool): MACRO void =


```

    printstr(concat(bool2str(b), emptystr))

println(s: string): MACRO void = printstr(concat(s, newline))

println(r:  $\mathbb{R}$ ): MACRO void =
    printstr(concat(real2str(r), newline))

println(b: bool): MACRO void =
    printstr(concat(bool2str(b), newline))

query_token(mssg, s: string): string

query_word(mssg: string): MACRO string = query_token(mssg, emptystr)

query_line(mssg: string): string

query_real(mssg: string):  $\mathbb{Q}$ 

query_int(mssg: string):  $\mathbb{Z}$ 

query_bool(mssg, answer: string): bool =
    str2bool(query_token(mssg, emptystr), answer)

read_token(s: string): MACRO string = query_token(emptystr, s)

read_word: MACRO string = query_token(emptystr, emptystr)

read_line: MACRO string = query_line(emptystr)

read_real: MACRO  $\mathbb{Q}$  = query_real(emptystr)

read_int: MACRO  $\mathbb{Z}$  = query_int(emptystr)

read_bool(answer: string): MACRO bool = query_bool(emptystr, answer)

Stream: TYPE+

IStream: TYPE+ FROM Stream

OStream: TYPE+ FROM Stream

fclose(f: Stream): void = skip

```

```

fexists(s: string): bool

fopen?(f: Stream): bool

strstream?(f: Stream): bool

filestream?(f: Stream): bool

sdtstream?(f: Stream): bool =
    NOT (filestream?(f) OR strstream?(f))

finput?(f: Stream): bool

foutput?(f: Stream): bool

stdin: IStream

stdout: OStream

stderr: OStream

Mode: TYPE = {input, output, create, append, overwrite, rename, str}

mode2str(m: Mode): string =
    CASES m OF
        input: "input",
        output: "output",
        create: "create",
        append: "append",
        overwrite: "overwrite",
        rename: "rename",
        str: "str"
    ENDCASES

tostr(m: Mode): MACRO string = mode2str(m)

fopenin_lisp(s: string): IStream

fopenout_lisp(s: string, n: ℕ): OStream

sopenin(s: string): IStream

```

sopenout(*s*: string): OStream

```
fopenin(m: Mode, s: string): IStream =  
  IF m = input AND length(s) = 0  
    THEN stdin  
  ELSIF m = input THEN fopenin_lisp(s)  
  ELSIF m = str THEN sopenin(s)  
  ELSE throw(WrongInputStreamMode, make_exc(WrongInputStreamMode, mode2str(m)))  
  ENDIF
```

fopenin(*s*: string): IStream = fopenin_lisp(*s*)

```
fopenout(m: Mode, s: string): OStream =  
  IF m = output AND length(s) = 0  
    THEN stdout  
  ELSIF m = output THEN fopenout_lisp(s, 0)  
  ELSIF m = create THEN fopenout_lisp(s, 1)  
  ELSIF m = append THEN fopenout_lisp(s, 2)  
  ELSIF m = overwrite THEN fopenout_lisp(s, 3)  
  ELSIF m = rename THEN fopenout_lisp(s, 4)  
  ELSIF m = str THEN sopenout(s)  
  ELSE throw(WrongOutputStreamMode, make_exc(WrongOutputStreamMode, mode2str(m)))  
  ENDIF
```

fopenout(*s*: string): OStream = fopenout(output, *s*)

fname_lisp(*f*: Stream): string

fgetstr_lisp(*f*: OStream): string

eof_lisp(*f*: IStream): bool

flength_lisp(*f*: Stream): \mathbb{N}

fgetpos_lisp(*f*: Stream): \mathbb{N}

fsetpos_lisp(*f*: Stream, *n*: \mathbb{N}): void

fprint_lisp(*f*: OStream, *s*: string): void = skip

fname(*f*: Stream): string =

```

IF filestream?(f) THEN fname_lisp(f) ELSE emptystr ENDIF

fgetstr(f: OStream): string =
  IF fopen?(f)
    THEN fgetstr_lisp(f)
  ELSE throw(ClosedStream, make_exc(ClosedStream, fname(f)))
  ENDIF

eof?(f: IStream): bool =
  IF fopen?(f)
    THEN eof_lisp(f)
  ELSE throw(ClosedStream, make_exc(ClosedStream, fname(f)))
  ENDIF

flength(f: Stream): ℕ =
  IF fopen?(f)
    THEN flength_lisp(f)
  ELSE throw(ClosedStream, make_exc(ClosedStream, fname(f)))
  ENDIF

fgetpos(f: Stream): ℕ =
  IF fopen?(f)
    THEN fgetpos_lisp(f)
  ELSE throw(ClosedStream, make_exc(ClosedStream, fname(f)))
  ENDIF

fprintf(f: OStream, s: string): void =
  IF fopen?(f)
    THEN fprintf_lisp(f, s)
  ELSE throw(ClosedStream, make_exc(ClosedStream, fname(f)))
  ENDIF

fprintf(f: OStream, r: ℝ): MACRO void =
  fprintf(f, concat(real2str(r), emptystr))

fprintf(f: OStream, b: bool): MACRO void =
  fprintf(f, concat(bool2str(b), emptystr))

fsetpos(f: Stream, n: ℕ): void =
  IF fopen?(f)
    THEN fsetpos_lisp(f, n)
  ELSE throw(ClosedStream, make_exc(ClosedStream, fname(f)))

```

ENDIF

fprintln(f : OStream, s : string): MACRO void =
fprint(f , concat(s , newline))

fprintln(f : OStream, r : \mathbb{R}): MACRO void =
fprint(f , concat(real2str(r), newline))

fprintln(f : OStream, b : bool): MACRO void =
fprint(f , concat(bool2str(b), newline))

echo(f : OStream, s : string): MACRO void =
printstr(s) & fprint(f , s)

echo(f : OStream, r : \mathbb{R}): MACRO void =
printstr(concat(real2str(r), emptystr)) &
fprint(f , concat(real2str(r), emptystr))

echo(f : OStream, b : bool): MACRO void =
printstr(concat(bool2str(b), emptystr)) &
fprint(f , concat(bool2str(b), emptystr))

echoIn(f : OStream, s : string): MACRO void =
printstr(concat(s , newline)) &
fprint(f , concat(s , newline))

echoIn(f : OStream, r : \mathbb{R}): MACRO void =
printstr(concat(real2str(r), newline)) &
fprint(f , concat(real2str(r), newline))

echoIn(f : OStream, b : bool): MACRO void =
printstr(concat(bool2str(b), newline)) &
fprint(f , concat(bool2str(b), newline))

fread_token_lisp(f : IStream, s : string): string

fread_line_lisp(f : IStream): string

fread_real_lisp(f : IStream): \mathbb{Q}

fread_int_lisp(f : IStream): \mathbb{Z}

```

fcheck(f: IStream): bool =
    (fopen?(f) OR throw(ClosedStream, make_exc(ClosedStream, fname(f))) AND
    (NOT eof?(f) OR
    throw(EndOfFile, make_exc(EndOfFile, fname(f))))

fread_token(f: IStream, s: string): string =
    prog(fcheck(f), fread_token_lisp(f, s))

fread_word(f: IStream): MACRO string = fread_token(f, emptystr)

fread_line(f: IStream): string =
    prog(fcheck(f), fread_line_lisp(f))

fread_real(f: IStream):  $\mathbb{Q}$  =
    prog(fcheck(f), fread_real_lisp(f))

fread_int(f: IStream):  $\mathbb{Z}$  =
    prog(fcheck(f), fread_int_lisp(f))

fread_bool(f: IStream, answer: string): MACRO bool =
    str2bool(fread_token(f, emptystr), answer)

END stdio

```

stdmath: THEORY

BEGIN

MathExceptions: list[ExceptionTag[string]] =
(:NotARealNumber, NotAnInteger:)

Π : $\mathbb{R}_{>0}$

SIN(x : \mathbb{R}): { x : \mathbb{R} | $-1 \leq x$ AND $x \leq 1$ }

COS(y : \mathbb{R}): { x : \mathbb{R} | $-1 \leq x$ AND $x \leq 1$ }

EXP(x : \mathbb{R}): $\mathbb{R}_{>0}$

RANDOM: { y : $\mathbb{R}_{\geq 0}$ | $0 \leq y$ AND $y \leq 1$ }

NRANDOM(n : $\mathbb{N}_{>0}$): { y : \mathbb{N} | $0 \leq y$ AND $y < n$ }

sqrt_lisp(x : $\mathbb{R}_{\geq 0}$): $\mathbb{R}_{\geq 0}$

log_lisp(x : $\mathbb{R}_{>0}$): \mathbb{R}

atan_lisp(x , y : \mathbb{R}): \mathbb{R}

asin_lisp(x : \mathbb{R}): \mathbb{R}

acos_lisp(x : \mathbb{R}): \mathbb{R}

SQRT(x : \mathbb{R}): $\mathbb{R}_{\geq 0}$ =
IF $x < 0$
THEN throw(NotARealNumber,
 make_exc(NotARealNumber, concat(concat("SQRT(", real2str(x)), ")")))
ELSE sqrt_lisp(x)
ENDIF

LOG(x : \mathbb{R}): \mathbb{R} =
IF $x \leq 0$
THEN throw(NotARealNumber,
 make_exc(NotARealNumber, concat(concat("LOG(", real2str(x)), ")")))
ELSE log_lisp(x)
ENDIF

```

TAN( $x: \mathbb{R}$ ):  $\mathbb{R}$  =
  LET  $d = \text{COS}(x)$  IN
    IF  $d = 0$ 
      THEN throw(NotARealNumber,
                 make_exc(NotARealNumber,
                          concat(concat("TAN (", real2str( $x$ )), "0) ")))
    ELSE SIN( $x$ )/COS( $x$ )
    ENDF

```

```

ATAN( $y, x: \mathbb{R}$ ):  $\mathbb{R}$  =
  IF  $x = 0$  AND  $y = 0$ 
    THEN throw(NotARealNumber, make_exc(NotARealNumber, "ATAN (0, 0) "))
  ELSE atan_lisp( $y, x$ )
  ENDF

```

```

ASIN( $x: \mathbb{R}$ ):  $\mathbb{R}$  =
  IF  $x < -1$  OR  $x > 1$ 
    THEN throw(NotARealNumber,
               make_exc(NotARealNumber, concat(concat("ASIN (", real2str( $x$ )), " ) ")))
  ELSE asin_lisp( $x$ )
  ENDF

```

```

ACOS( $x: \mathbb{R}$ ):  $\mathbb{R}$  =
  IF  $x < -1$  OR  $x > 1$ 
    THEN throw(NotARealNumber,
               make_exc(NotARealNumber, concat(concat("ACOS (", real2str( $x$ )), " ) ")))
  ELSE acos_lisp( $x$ )
  ENDF

```

```

BRANDOM: bool = (NRANDOM(2) = 0)

```

```

END stdmath

```



```

stdfmap[T: TYPE+]: THEORY
BEGIN

  fmap(f: IStream, fread: [IStream → string], t: T, st: [[T] → T],
        n: ℕ): RECURSIVE
    T =
  IF fopen?(f)
    THEN IF n = 0 OR eof?(f)
          THEN t
          ELSE LET s = fread(f) IN
                LET nt = st(s, t) IN fmap(f, fread, nt, st, n - 1)
          ENDIF
    ELSE throw(ClosedStream, make_exc(ClosedStream, fname(f)))
    ENDIF
  MEASURE n

  fmap(f: IStream, fread: [IStream → string], t: T, st: [[T] → T]): T =
    LET l = flength(f) IN
      LET nt = fmap(f, fread, t, st, l) IN prog(fclose(f), nt)

  fmap_line(f: IStream, t: T, st: [[T] → T]): T =
    LET l = flength(f) IN
      LET nt = fmap(f, fread_line, t, st, l) IN
        prog(fclose(f), nt)

  printf(s: string, t: T): MACRO void = printstr(format(s, t))

  fprintf(f: OStream, s: string, t: T): MACRO void =
    fprintf(f, format(s, t))

END stdfmap

```

```

stdindent: THEORY
BEGIN

Indent: TYPE+

create_indent(n: ℕ, s: string): Indent

push_indent(i: Indent, n: ℕ): void

pop_indent(i: Indent): void

top_indent(i: Indent): ℕ

get_indent(i: Indent): ℕ

set_indent(i: Indent, n: ℕ): void

get_prefix(i: Indent): string

set_prefix(i: Indent, s: string): void

create_indent(n: ℕ): MACRO Indent = create_indent(n, emptystr)

open_block(i: Indent, n: ℕ): MACRO void = push_indent(i, n)

open_block(i: Indent): MACRO void =
    push_indent(i, get_indent(i))

close_block(i: Indent): MACRO void = pop_indent(i)

indent(i: Indent): string =
    concat(get_prefix(i), spaces(top_indent(i)))

indent(i: Indent, s: string): string = concat(indent(i), s)

prindent(i: Indent, s: string): void = printstr(indent(i, s))

prindentln(i: Indent, s: string): void =
    printstr(concat(indent(i, s), newline))

fprindent(f: OStream, i: Indent, s: string): void =
    fprintf(f, indent(i, s))

```

```

fprindentln(f: OStream, i: Indent, s: string): void =
    fprint(f, concat(indent(i, s), newline))

center(col: ℕ, s: string): string =
    format(concat(concat("~", real2str(col)), ":@<~a~>"), s)

flushleft(col: ℕ, s: string): string =
    format(concat(concat("~", real2str(col)), "a"), s)

flushright(col: ℕ, s: string): string =
    format(concat(concat("~", real2str(col)), "@a"), s)

END stdindent

```

stdtokenizer: THEORY

BEGIN

NoError: $\mathbb{N} = 0$

FileNotFound: $\mathbb{N} = 1$

EndOfTokenizer: $\mathbb{N} = 2$

InvalidToken: $\mathbb{N} = 3$

ExpectingWord: $\mathbb{N} = 4$

ExpectingTestWord: $\mathbb{N} = 5$

ExpectingInt: $\mathbb{N} = 6$

ExpectingTestInt: $\mathbb{N} = 7$

ExpectingReal: $\mathbb{N} = 8$

ExpectingTestReal: $\mathbb{N} = 9$

Tokenizer: TYPE =

[#stream: [$\mathbb{N} \rightarrow \text{string}$],
lines: [$\mathbb{N} \rightarrow \mathbb{N}$],
casesen: boolean,
separ: string,
error: \mathbb{Z} ,
val_int: \mathbb{Z} ,
val_real: \mathbb{R} ,
length: \mathbb{N} ,
pos: upto(length)#]

init_tokenizer(casesen: bool, separ: string): Tokenizer =

(#stream := $\lambda (x: \mathbb{N}): \text{emptystr}$,
lines := $\lambda (x: \mathbb{N}): 0$,
casesen := casesen,
separ := separ,
error := NoError,
val_int := 0,
val_real := 0,

```

length := 0,
pos := 0#)

empty_tokenizer: Tokenizer =
  (#stream := λ (x: ℕ): emptystr,
   lines := λ (x: ℕ): 0,
   casesen := TRUE,
   separ := emptystr,
   error := NoError,
   val_int := 0,
   val_real := 0,
   length := 0,
   pos := 0#)

TokenizerOfLength(l: ℤ): TYPE = {t: Tokenizer | t`length = l}

set_casesen(t: Tokenizer, c: bool): Tokenizer =
  t WITH [ `casesen := c ]

error?(t: Tokenizer): MACRO bool = t`error ≠ NoError

set_error(t: Tokenizer, code: ℤ): TokenizerOfLength(t`length) =
  t WITH [ `error := code ]

last_token(t: Tokenizer): string =
  IF t`pos = 0
  THEN emptystr
  ELSE t`stream(t`pos - 1)
  ENDIF

peek(t: Tokenizer, n: ℕ>0): MACRO string =
  t`stream(t`pos + n - 1)

next_token(t: Tokenizer): MACRO string =
  t`stream(t`pos + 1 - 1)

tostr(t: Tokenizer, i: upto(t`length)): RECURSIVE string =
  IF i = t`length
  THEN emptystr
  ELSIF i = t`pos
  THEN concat(concat(concat(" [ ", t`stream(i)), " ] "), tostr(t, i + 1))
  ELSE concat(concat(t`stream(i), space), tostr(t, i + 1))

```

```

ENDIF
MEASURE t`length - i

add_token(s: string, t: Tokenizer, l: ℕ): MACRO Tokenizer =
  LET n = t`length IN
    t
    WITH [ `stream(n) := s,
           `lines(n) := l,
           `length := n + 1 ]

read_token(t: Tokenizer)(f: IStream): string =
  fread_token(f, t`separ)

line_tokenizer(s: string, tl: [ℕ]): [ℕ] =
  LET (t, l) = tl IN
    LET g = fopenin(str, s) IN
      LET f =
        (λ (mys: string, myt: Tokenizer):
          LET n = myt`length IN
            myt
            WITH [ `stream(n) := mys,
                  `lines(n) := l,
                  `length := n + 1 ])
      IN
        LET nt = fmap(g, read_token(t), t, f, length(s)) IN
          prog(fclose(g), (nt, l + 1))

file2tokenizer(s: string, t: Tokenizer): Tokenizer =
  IF fexists(s)
    THEN LET (nt, l) = fmap_line(fopenin(s), (t, 1), line_tokenizer) IN nt
  ELSE LET filename = s IN t WITH [ `stream(0) := filename, `error := FileNotFound ]
  ENDF

file2tokenizer(s: string): Tokenizer =
  file2tokenizer(s, empty_tokenizer)

str2tokenizer(s: string, t: Tokenizer): Tokenizer =
  LET f = fopenin(str, s) IN
    LET (nt, l) = fmap_line(f, (t, 1), line_tokenizer) IN nt

str2tokenizer(s: string): Tokenizer =
  str2tokenizer(s, empty_tokenizer)

```

eot?(t: Tokenizer): bool = t`pos = t`length

get_line(t: Tokenizer): MACRO $\mathbb{N} = t`lines(t`pos)$

consume(t: Tokenizer, n: $\mathbb{N}_{>0}$): TokenizerOfLength(t`length) =
IF t`error \neq NoError
THEN t
ELSIF eot?(t) OR t`pos + n > t`length THEN t WITH [\backslash error := EndOfTokenizer]
ELSE t WITH [\backslash pos := t`pos + n]
ENDIF

go_next(t: Tokenizer): MACRO TokenizerOfLength(t`length) =
consume(t, 1)

go_back(t: Tokenizer): TokenizerOfLength(t`length) =
IF t`pos = 0
THEN t
ELSE t WITH [\backslash pos := t`pos - 1]
ENDIF

pos_go_next: LEMMA

$\forall (t_1: \text{Tokenizer}):$
LET $t_2 = \text{consume}(t_1, 1)$ IN
NOT $t_2`error \neq \text{NoError}$ IMPLIES $t_2`pos = t_1`pos + 1$

accept_word(t: Tokenizer, test: [string \rightarrow bool]): TokenizerOfLength(t`length) =
IF t`error \neq NoError
THEN t
ELSIF eot?(t) THEN t WITH [\backslash error := EndOfTokenizer]
ELSIF number?(t`stream(t`pos)) THEN t WITH [\backslash error := ExpectingWord]
ELSIF test(t`stream(t`pos)) THEN t WITH [\backslash pos := t`pos + 1]
ELSE t WITH [\backslash error := ExpectingTestWord]
ENDIF

pos_accept_word: LEMMA

$\forall (t_1: \text{Tokenizer}, \text{test}: [\text{string} \rightarrow \text{bool}]):$
LET $t_2 = \text{accept_word}(t_1, \text{test})$ IN
NOT $t_2`error \neq \text{NoError}$ IMPLIES $t_2`pos = t_1`pos + 1$

the_word(s: string)(token: string): bool = str2bool(s, token)

```

accept_word(t: Tokenizer, s: string): MACRO TokenizerOfLength(t`length) =
  accept_word(t, the_word(s))

any_word(token: string): bool = TRUE

accept_word(t: Tokenizer): MACRO TokenizerOfLength(t`length) =
  accept_word(t, any_word)

accept_int(t: Tokenizer, test: [ $\mathbb{Z} \rightarrow \text{bool}$ ]): TokenizerOfLength(t`length) =
  IF t`error  $\neq$  NoError
    THEN t
  ELSIF eot?(t) THEN t WITH [ $\text{error} := \text{EndOfTokenizer}$ ]
  ELSIF NOT int?(t`stream(t`pos)) THEN t WITH [ $\text{error} := \text{ExpectingInt}$ ]
  ELSE LET  $i = \text{str2int}(t`stream(t`pos))$  IN
    IF test( $i$ )
      THEN t WITH [ $\text{pos} := t`pos + 1, \text{val\_int} := i, \text{val\_real} := i$ ]
    ELSE t WITH [ $\text{error} := \text{ExpectingTestInt}$ ]
  ENDIF
ENDIF

pos_accept_int: LEMMA
 $\forall (t_1: \text{Tokenizer}, \text{test}: [\mathbb{Z} \rightarrow \text{bool}]):$ 
  LET  $t_2 = \text{accept\_int}(t_1, \text{test})$  IN
    NOT  $t_2`error \neq \text{NoError}$  IMPLIES  $t_2`pos = t_1`pos + 1$ 

any_int( $i: \mathbb{Z}$ ): bool = TRUE

accept_int(t: Tokenizer): MACRO TokenizerOfLength(t`length) =
  accept_int(t, any_int)

accept_real(t: Tokenizer, test: [ $\mathbb{R} \rightarrow \text{bool}$ ]): TokenizerOfLength(t`length) =
  IF t`error  $\neq$  NoError
    THEN t
  ELSIF eot?(t) THEN t WITH [ $\text{error} := \text{EndOfTokenizer}$ ]
  ELSIF NOT number?(t`stream(t`pos)) THEN t WITH [ $\text{error} := \text{ExpectingReal}$ ]
  ELSE LET  $r = \text{str2real}(t`stream(t`pos))$  IN
    IF test( $r$ )
      THEN t WITH [ $\text{pos} := t`pos + 1, \text{val\_real} := r$ ]
    ELSE t WITH [ $\text{error} := \text{ExpectingTestReal}$ ]
  ENDIF
ENDIF

```



```

pos_accept_real: LEMMA
  ∀ (t1: Tokenizer, test: [ℝ → bool]):
    LET t2 = accept_real(t1, test) IN
      NOT t2`error ≠ NoError IMPLIES t2`pos = t1`pos + 1

any_real(r: ℝ): bool = TRUE

accept_real(t: Tokenizer): MACRO TokenizerOfLength(t`length) =
  accept_real(t, any_real)

Messenger: TYPE = [ℤ → string]

std_mssg(t: Tokenizer)(code: ℤ): string =
  IF code ≤ 0
    THEN emptystr
  ELSIF code = FileNotFound
    THEN concat(concat(concat(concat("File not found: ", doublequote),
                                     t`stream(t`pos + 1 - 1)),
                                     doublequote),
                ".")
  ELSIF code = EndOfTokenizer THEN "Found EOT."
  ELSIF code = InvalidToken
    THEN concat(concat(concat(concat("Invalid Token: ", doublequote),
                                     t`stream(t`pos + 1 - 1)),
                                     doublequote),
                ".")
  ELSIF code = ExpectingWord
    THEN concat(concat(concat(concat("Expecting a word. Found: ", double-
quote),
                                     t`stream(t`pos + 1 - 1)),
                                     doublequote),
                ".")
  ELSIF code = ExpectingTestWord
    THEN concat(concat(concat(concat("Expecting word that satisfies test. Found:
                                     doublequote),
                                     t`stream(t`pos + 1 - 1)),
                                     doublequote),
                ".")
  ELSIF code = ExpectingInt
    THEN concat(concat(concat(concat("Expecting an integer. Found: ", dou-
blequote),
                                     t`stream(t`pos + 1 - 1)),
                                     doublequote),
                ".")

```

```

                                doublequote),
                                ".")
    ELSIF code = ExpectingTestInt
        THEN concat(concat(concat(concat("Expecting integer that satisfies test. Found: ",
                                doublequote),
                                t`stream(t`pos + 1 - 1)),
                                doublequote),
                                ".")
    ELSIF code = ExpectingReal
        THEN concat(concat(concat(concat("Expecting a real. Found: ", double-
quote),
                                t`stream(t`pos + 1 - 1)),
                                doublequote),
                                ".")
    ELSIF code = ExpectingTestReal
        THEN concat(concat(concat(concat("Expecting real that satisfies test. Found: ",
                                doublequote),
                                t`stream(t`pos + 1 - 1)),
                                doublequote),
                                ".")
    ELSE emptystr
    ENDIF

```

```

print_error(t: Tokenizer, m: Messenger): MACRO void =
    IF t`error ≠ NoError
        THEN printstr(concat(concat(concat(concat(concat("Syntax Error. ", "Line "),
                                real2str(t`lines(t`pos))),
                                ": "),
                                m(error(t))),
                                newline))
        & fail
    ELSE skip
    ENDIF

```

```

print_error(t: Tokenizer): MACRO void =
    IF t`error ≠ NoError
        THEN printstr(concat(concat(concat(concat(concat("Syntax Error. ", "Line "),
                                real2str(t`lines(t`pos))),
                                ": "),
                                std_mssg(t)(error(t))),
                                newline))
        & fail
    ELSE skip
    ENDIF

```

```

ELSE skip
ENDIF;

m : Messenger × cs : [string] : Messenger =
  LET (code, s) = cs IN m WITH [ `code := s ]

tokenizer2str(t: Tokenizer): string =
  concat(concat(tostr(t, 0), newline),
    std_mssg(t)(t`error))

CONVERSION tokenizer2str

tostr(t: Tokenizer): MACRO string = tokenizer2str(t)

END stdtokenizer

```

```
stdpvzio: THEORY
BEGIN

    help_pvs_attachment(s: string): void

    help_pvs_theory_attachments(s: string): void

    pvzio_version: string

    set_promptin(s: string): void

    set_promptout(s: string): void

END stdpvzio
```

```
stdsys: THEORY
BEGIN
```

```
Time: TYPE =
  [#second: below(60),
   minute: below(60),
   hour: below(24),
   day: subrange(1, 31),
   month: subrange(1, 12),
   year:  $\mathbb{N}$ ,
   dow: below(7),
   dst: bool,
   tz: { $x: \mathbb{Q} \mid -24 \leq x \text{ AND } x \leq 24$ }#]
```

```
tinybang: Time =
  (#day := 1,
   dow := 0,
   dst := FALSE,
   hour := 0,
   minute := 0,
   month := 1,
   second := 0,
   tz := 0,
   year := 0#)
```

```
days_of_week(dow: below(7)): string =
  IF dow = 0
    THEN "Monday"
  ELSIF dow = 1 THEN "Tuesday"
  ELSIF dow = 2 THEN "Wednesday"
  ELSIF dow = 3 THEN "Thursday"
  ELSIF dow = 4 THEN "Friday"
  ELSIF dow = 5 THEN "Saturday"
  ELSE "Sunday"
ENDIF
```

```
months(month: subrange(1, 12)): string =
  IF month = 1
    THEN "January"
  ELSIF month = 2 THEN "February"
  ELSIF month = 3 THEN "March"
  ELSIF month = 4 THEN "April"
```

```

ELSIF month = 5 THEN "May"
ELSIF month = 6 THEN "June"
ELSIF month = 7 THEN "July"
ELSIF month = 8 THEN "August"
ELSIF month = 9 THEN "September"
ELSIF month = 10 THEN "October"
ELSIF month = 11 THEN "November"
ELSE "December"
ENDIF

```

```

tostr(t: Time): string =
  concat(concat(concat(concat(concat(days_of_week(dow(t)), " "),
                                     months(month(t))),
                                     format(" ~2, '0d ", day(t))),
                                     real2str(year(t))),
          format(", ~2, '0d:~2, '0d:~2, '0d ",
                (hour(t), minute(t), second(t))),
          format(" (GMT~@d) ", -tz(t)))

```

```

get_time: Time

```

```

today: string =
  LET t = get_time IN
    format("~d/~2, '0d/~d", (month(t), day(t), year(t)))

```

```

date: string = LET t = get_time IN tostr(get_time)

```

```

sleep(n: ℕ): void

```

```

get_env(name, default: string): string

```

```

get_env(name: string): MACRO string = get_env(name, emptystr)

```

```

printf(s: string): MACRO void = printstr(format(s, ""))

```

```

fprintf(f: OStream, s: string): MACRO void =
  fprintf(f, format(s, ""))

```

```

Blisp: TYPE+

```

```

blisp(b: bool): Blisp

```

```
{||}(b: bool): Blisp = blisp(b)
```

```
END stdsys
```